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Nguyen Nhu Thang

Dai hoc Ky thuat Le Quy Don

Tran Manh Hoang

Dai hoc Ky thuat Le Quy Don

Ba Cao Nguyen

Dai hoc Ky thuat Le Quy Don

Phuong T. Tran (✉ tranthanhphuong@tdtu.edu.vn)

Ton Duc Thang University <https://orcid.org/0000-0002-1448-8882>

Research

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Posted Date: December 3rd, 2020

DOI: <https://doi.org/10.21203/rs.3.rs-117143/v1>

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RESEARCH

Enhancing the performance of downlink NOMA relaying networks by RF energy harvesting and data buffering at relay

Nguyen Nhu Thang¹, Tran Manh Hoang¹, Ba Cao Nguyen¹ and Phuong T. Tran^{2*}

*Correspondence:

tranthanhphuong@tdtu.edu.vn

²Wireless Communications

Research Group, Faculty of

Electrical and Electronics

Engineering, Ton Duc Thang

University, Ho Chi Minh City,

Vietnam

Full list of author information is
available at the end of the article

Abstract

Recently, non-orthogonal multiple access (NOMA) has been considered as a promising candidate for next-generation mobile communications because it can significantly improve the spectral efficiency of wireless networks. In this paper, we investigate a novel solution to enhance the reliability and the supply stability of a downlink NOMA relaying networks, in which we integrate two techniques: (i) simultaneous wireless information and power transfer, i.e. the relay node can harvest the energy from source signals and use this energy to help forward information from source node to two user nodes; and (ii) data buffer aid at relay node, i.e. the data packets received from the source can be stored in a buffer and then be re-transmitted to the destination nodes only when the channel condition is good. The performance of the proposed system is analyzed rigorously to derive the system outage probability and the average packet delay. Furthermore, a power allocation optimization problem to minimize the outage probability is formulated and solution to this problem is also provided in this paper. Monte Carlo simulations are conducted to verify the analytical results, which confirms that with the data buffer at the relay, the overall outage probability (OOP) has been reduced significantly.

Keywords: NOMA; energy harvesting; successive interference cancellation; power allocation; buffer-aided

Introduction

The non-orthogonal multiple-access (NOMA) is considered as a promising multiuser communication technique for the fifth-generation (5G) mobile network since

¹it can achieve superior spectral efficiency [1]. Unlike the orthogonal-multiple-access¹
²(OMA), NOMA users can share the same radio resources, including time and band-²
³width. The key idea for this advantage is to employ the power domain or codes³
⁴domain, where different users are distinguished by their power levels or different⁴
⁵code[2]. Recent researcher have shown that NOMA can be applied to not only⁵
⁶point-to-point but also relay networks [3]. While the application of NOMA to the⁶
⁷point-to-point networks were well investigated, there are still increasing needs for⁷
⁸the case of cooperative relaying networks [4, 5]. The work in [4] studied the conven-⁸
⁹tional cooperative NOMA system with buffer-aided relaying. Under the assumption⁹
¹⁰that the relay node possesses a buffer. Herein the authors considered an adaptive¹⁰
¹¹transmission scheme in which different working modes are employed in different¹¹
¹²time slots. The authors of [5] proposed a dual-hop cooperative relaying scheme us-¹²
¹³ing NOMA, where two source nodes communicate with each other simultaneously¹³
¹⁴via a common relay on the same frequency band. In this scheme, after receiving¹⁴
¹⁵symbols transmitted in parallel by both sources with different power levels, the¹⁵
¹⁶relay forwards the superposition coded composite signal using NOMA to two des-¹⁶
¹⁷tinations. However, in this work power control for uplink multiple access was not¹⁷
¹⁸considered. 18

¹⁹ In addition, harvesting energy from the ambient environment has become a ¹⁹
²⁰promising solution for energy-constrained electronic devices, which are convention-²⁰
²¹ally supported by limited power sources such as battery [6, 7, 8, 9]. In some special²¹
²²applications, charging the battery is too expensive or even impossible, e.g. sensor²²
²³network works under toxic environment and body area network. Moreover, some²³
²⁴natural energy sources such as solar and wind, and radio frequency (RF) can be²⁴
²⁵also utilized as effective sources for energy harvesting (EH). Compared with other²⁵
²⁶kinds of energy sources, the RF energy harvesting [10], also known as wireless en-²⁶
²⁷ergy transfer, has some advantages. Since the RF energy harvesting is an active²⁷
²⁸energy supply method, it can provide more reliable energy flow to guarantee the²⁸
²⁹quality of service. 29

³⁰ Therefore, utilizing RF-EH technique together with NOMA scheme helps prolong ³⁰
³¹the lifetime and improves the spectral utilization efficiency of the energy-constrained ³¹
³²multi-user wireless relaying networks. The NOMA systems combining with RF en- ³²
³³ergy harvesting is investigated in [11, 12, 13, 14]. A simultaneous wireless informa- ³³

tion and power transfer (SWIPT) of NOMA networks were considered in [11], where¹
base station serves two types of users, namely, relay user and far user. The outage²
performance and spectral efficiency of the NOMA-EH relaying networks with an-³
tenna selection were investigated in [13], where the power splitting (PS) protocol is⁴
applied at the relay to harvest the energy. In [13], the performance of NOMA sys-⁵
tem is also compared with OMA system. The authors of [14] investigated a NOMA⁶
system in which NOMA users near to the source act as the EH relays to assist the⁷
far NOMA users in forwarding the information. In these works, the authors con-⁸
sidered the users reception signals in two different time slots, however, the service⁹
was simultaneously provided to users in the NOMA system. The impact of power¹⁰
allocation in the cooperative NOMA network with SWIPT was investigated in [3].¹¹
In this work, Yang et. al. proposed two types of NOMA power allocation policies,¹²
namely NOMA with fixed power allocation (F-NOMA) and cognitive radio inspired¹³
NOMA (CR-NOMA). The results of these above-mentioned works shown that the¹⁴
performance of NOMA outperform OMA scheme. However, the diversity gain was¹⁵
not improved. 16

An energy buffer-aided EH relay was applied in cooperative communication sys-¹⁷
tem to improve system performance in [15], but in this work the authors didn't¹⁸
consider buffer-aided data. On the other hand, the work in [16] proposed a hybrid¹⁹
NOMA/OMA system with buffer-aided relay selection. As a result, two buffer-aided²⁰
opportunistic relay selection algorithms were proposed. The aim of that work is to²¹
improve the outage performance and sum-rate of the system. 22

The authors in [17] proposed a priority-based max-link relay selection for data-²³
buffer-aided decode-and-forward DF cooperative networks. In this work, the authors²⁴
derived analytical expressions of outage probability and bit error rate to evaluate²⁵
the system performance. In addition to NOMA downlink system investigation, the²⁶
hybrid NOMA/OMA uplink system with the help of a buffer-aided relay was also²⁷
considered in [18]. 28

So far, all previous works related to cooperative communication protocols and²⁹
NOMA technique in literature have proposed the superposition signal coding at³⁰
the source. Meanwhile, the relay node, which has a fixed power, only decodes and³¹
forwards signals to destinations. However, due to the random nature of wireless³²
channel, the amount of energy harvested at the relay is usually very small and vari-³³

able. Hence, the power allocation at the relay can reduce the feedback energy and¹
 guarantee the performance fairness for all users. To the best of our knowledge, com-²
 bining SWIPT with NOMA relaying system where the buffer-aided relay technique³
 is employed at the relay has not been investigated in literature and the derivation⁴
 of the overall outage probability expression of this system has not been carried out,⁵
 either. Motivated by these facts, in this paper, we proposed a different cooperative⁶
 decode-and-forward relaying scheme where a source transmits information packets⁷
 to the relay, while relay broadcasts the modulation superposition signals to two⁸
 users. However, R employs time switching based EH prior to the communication⁹
 with destinations. Based on the channel gain from the relay node to the destination¹⁰
 node, the relay performs fixed and optimal power allocation for two users. Further-¹¹
 more, NOMA technology is investigated for the cooperative transmission in term¹²
 of two scenarios, i.e. with and without buffer aid at the relay node. ¹³

14

The main contributions of this paper are summarized as below: ¹⁵

16

- We proposed a novel downlink NOMA relay system applying SWIPT, where¹⁶
 designated relay node is either equipped with buffer or not. To harvest energy,¹⁷
 the relay uses time switching protocol. The performance of the system is¹⁸
 improved, and additionally the spectral utilization efficiency and the lifetime¹⁹
 of wireless networks will be enhanced. ²⁰
- The optimal power allocation at the relay node to minimum outage probability²¹
 is also considered in this paper. Since the harvested energy is very small, the²²
 reallocation of the harvested energy after converting from the RF signals of²³
 the source is important for saving cost. ²⁴
- Markov chain model and state-transition matrix is used to describe the ran-²⁵
 dom process at the buffer-aided relay. On the other hand, with buffer-aided²⁶
 relay the diversity gain is improved significantly. ²⁷
- The system performance is demonstrated by the outage probability and the²⁸
 sum end-to-end ergodic capacity over Rayleigh fading channel. We derive the²⁹
 closed-form expression to evaluate the rate of symbols and outage performance³⁰
 of NOMA-SWIPT system. ³¹
- The analytical results are validated by simulation. From the derived closed-³²
 form expressions, practical networks can be investigated in the in future. ³³

¹ Our proposed system can be applied in surveillance sensor networks for disaster¹
² detection or Internet of Things (IoT), where installing fixed power lines or frequently²
³ replacing the batteries for a large number of nodes is not convenient. Besides, its³
⁴ advantages such as low energy cost, reducing greenhouse effect, and prolonging⁴
⁵ lifetime are useful for future mobile networks. 5

⁶ The remaining of the paper is organized as follows. Section “System model”⁶
⁷ presents the NOMA-SWIPT system model and channel model. The analysis of⁷
⁸ outage probability with and without buffer aid at the relay node is given in Sec-⁸
⁹ tions “Outage probability without buffer-aided relay” and “Analysis of the out-⁹
¹⁰ age probability with buffer-aided relay”, respectively. The average packet delay is¹⁰
¹¹ demonstrated in Section “Average packet delay of the buffer-aided relay system”.¹¹
¹² Numerical results, which verify our analysis, are presented in Section “Numerical¹²
¹³ results”. Finally, the conclusion is given in Section “Conclusion”. 13

¹⁴ For the convenience, we provide in Table 1 the notations along with their descrip-¹⁴
¹⁵ tions used in this paper. 15

¹⁶ **Table 1 Mathematical notations** 16

Notation	Description	17
Pr	Probability	17
$F_X(x)$	Cumulative distribution function (CDF)	18
$f_X(x)$	Probability density function (PDF)	19
$\mathcal{CN}(\mu, \sigma^2)$	A circularly symmetric complex Gaussian RV x with mean μ and variance σ^2	20
$\mathbb{E}\{\cdot\}$	The statistical expectation operator	20
$\Gamma(\cdot)$	Gamma function [27]	21
$\mathcal{K}_n(\cdot)$	The second kind of Bessel function order n [27]	22
$E_n(z)$	Exponential integral function n [27]	22
$G_{pq}^{mn}(x _{b_s}^{a_r})$	Meijer's G-Function [27, 9.3]	23

²⁴ 24

²⁵ **Methods** 25

²⁶ System model 26

²⁷ The system model of a NOMA downlink relaying network investigated in this paper²⁷
²⁸ is shown in Fig 1. According to this model, a source (S) wants to send its messages²⁸
²⁹ to two destinations (D_1) and (D_2) simultaneously with the help of relay node, (R),²⁹
³⁰ which is capable of energy harvesting. It is assumed that S, D_1 , and D_2 have fixed³⁰
³¹ power supply while relay node have no extra embedded energy supply, hence, R³¹
³² needs to harvest energy from S. In addition, the relay node is assumed to have an³²
³³ unlimited-size information buffer to store the received messages [19]. We assume³³

¹that the direct link between the source and destination is not available due to far¹
²distance or deep shadow fading in the channel. All nodes are equipped with a single²
³antenna and operate in a half-duplex mode^[1]. 3

⁴ The channels between two arbitrary nodes are subject to block and flat Rayleigh⁴
⁵fading. This means that the channel coefficients are constant during each data block⁵
⁶transmission interval T but vary from one block to another. 6

⁷ In the case that the relay uses a buffer for data processing, we assume that it⁷
⁸has perfect channel state information (CSI) of the links $S \rightarrow R$ and $D_1, D_2 \rightarrow R$ ⁸
⁹at the beginning of each time slot by using a short reference signals. Based on this⁹
¹⁰set of information, R can decide whether it is ready to operate in transmitting or¹⁰
¹¹receiving mode [20]. 11

¹² As shown in Fig. 1, the complex channel coefficient of the link between S and R¹²
¹³is denoted by $h_1 \sim \mathcal{CN}(0, \Omega_0)$. The complex channel coefficient between R and D_i ,¹³
¹⁴is $g_i \sim \mathcal{CN}(0, \Omega_i)$, where $i = \{1, 2\}$ and $\Omega_1 = \mathbb{E}\{|g_1|^2\}$, $\Omega_2 = \mathbb{E}\{|g_2|^2\}$. The additive¹⁴
¹⁵white Gaussian noise (AWGN) at R, D_1 and D_2 is respectively denoted by $w_{\mathcal{A}} \sim$ ¹⁵
¹⁶ $\mathcal{CN}(0, N_0)$, where $\mathcal{A} \in \{R, D_1, D_2\}$. Without loss of generality, we assume that the¹⁶
¹⁷users' channel gains are sorted in the descending order as follows: $|g_1|^2 > |g_2|^2$.¹⁷
¹⁸ 18

¹⁹*Data buffer at relay* 19

²⁰ In order to process the data the relay is equipped with a buffer for storing the²⁰
²¹signals received from the source. For time switching (TS) scheme the relay also²¹
²²has an energy storage device to store the harvested energy^[2]. R first harvests the²²
²³energy from the RF signal transmitted by S and then performs signal reception and²³
²⁴transmission using the harvestet and transmit strategy [21]. 24

²⁵ The system operates according to the time-division duplex (TDD) mode where²⁵
²⁶each transmission period is divided into equal time slots of length $\tau(1 - \alpha)$. At each²⁶
²⁷time slot, the relay or source node is selected to transmit data depending on the²⁷
²⁸status of the relay buffers and the available links that can provide the successful²⁸
²⁹transmission or reception of one packet. 29

³⁰^[1]This model can employ two antennas for the relay node and operate in a full-³⁰
³¹duplex mode. 31

³²^[2]The storing of harvested energy in TS scheme is refered as charging-then-³²
³³communicate. In contrast, PS scheme is refered as charging-and-communicate. 33

¹ If S is selected, it will generate a transmission frame of size $2r_0\tau$ bits, intended for¹
²two destination nodes D_1 and D_2 , to send to the relay node, where r_0 is the target²
³transmission rate of the system. Each frame contains two segments, the first one is³
⁴used for transmission symbols to D_1 and the second one is used for transmission⁴
⁵symbols to D_2 . The relay buffer has $L \geq 2$ storage units, each can store $2r_0\tau(1-\alpha)$ ⁵
⁶bits. The relay node decodes the received frame and stores it into the storage device.⁶
⁷Each storage device is split into two parts of the same length, which are used to⁷
⁸store the information symbols intended for D_1 and D_2 , respectively. ⁸

¹⁰*Signal model* ¹⁰

¹¹In each time slot, if the source is selected to transmit with unicast communication,¹¹
¹²it combines two signals x_1 and x_2 into a transmission packet. Then, the received¹²
¹³signal at the relay is given by ¹³

$$\sup{15} \quad y_R = h_1 \sqrt{P_S} x_S + w_R. \quad (1) \sup{15}$$

¹⁷The signal-to-noise ratio (SNR) of the source-to-relay link is given by [19] ¹⁷

$$\sup{19} \quad \gamma_R = \frac{P_S |h_1|^2}{N_0}. \quad (2) \sup{19}$$

²¹ The time switching (TS) architecture for harvesting energy is applied as [22].²¹
²²Refer to [22] for more detailed explanation^[3]. ²²

²³ Herein, T denote the block duration of an entire communication period in which²³
²⁴the information is transmitted from S to D_i . For each period T , the first amount²⁴
²⁵of time, αT , is used for EH at R, while the remaining amount of time, $(1-\alpha)T$,²⁵
²⁶is used for transmitting and receiving the information, where α denotes the EH²⁶
²⁷time fraction in one transmission block and $0 \leq \alpha \leq 1$. Therefore, the amount of²⁷
²⁸harvested energy at the relay for the case of linear model in the i th time slot is²⁸
²⁹given by [24, 25] ²⁹

$$\sup{31} \quad E_h = \alpha T \eta P_S |h_1|^2, \quad (3) \sup{31}$$

³²^[3]The proposed analytical approach can be applied to the power spitting EH model ³²
³³[23] ³³

¹where η denotes the energy conversion efficiency whose value ranges from 0 to 1,¹
²depending on the harvesting electric circuitry.²

³ *Remark 1:* In practice, the energy harvester will output a constant power because³
⁴of the circuit design for EH-RF. The key of such a non-linear model is that it can⁴
⁵capture the joint effect of the non-linear phenomena caused by hardware constraints⁵
⁶including circuit sensitivity limitations and current leakage. The main cause of non-⁶
⁷linear energy harvesting models can be mentioned by the relationship between the⁷
⁸input RF power and the output direct current of energy harvester. The cause to⁸
⁹make the nonlinear function can be explained by circuit devices such as diodes and⁹
¹⁰transistors in the energy harvester structure. Denote P_{th} as the saturation power¹⁰
¹¹threshold of harvester, thus the transmit power of the relay node in proposed model¹¹
¹²of manuscript is given by¹²

$$P_{\text{R}} = \begin{cases} \frac{2\alpha\eta}{1-\alpha} P_{\text{S}} |h_1|^2, & P_{\text{S}} |h_1|^2 \leq P_{\text{th}} \\ \frac{2\alpha\eta}{1-\alpha} P_{\text{th}}, & P_{\text{S}} |h_1|^2 > P_{\text{th}}. \end{cases} .$$

¹³
¹⁴
¹⁵
¹⁶
¹⁷ We assume that all the amount of power harvested is consumed by the relay¹⁷
¹⁸for forwarding signals to all users D_i , the processing power for the transmitting/¹⁸
¹⁹receiving circuitry at the relay is generally negligible compared to the power used¹⁹
²⁰for signal transmission and perhaps venial. So, from (3), the transmission power of²⁰
²¹the relay is given as²¹

$$P_{\text{R}} = \frac{E_h}{(1-\alpha)T/2} = \frac{2\alpha\eta P_{\text{S}} |h_1|^2}{(1-\alpha)}. \quad (4)$$

²²
²³
²⁴
²⁵ In a specific time slot, if R is selected, it transmits a modulation superimposition²⁵
²⁶information symbol $x_{\text{R}} = \sqrt{a_1 P_{\text{R}}} x_1 + \sqrt{(1-a_1) P_{\text{R}}} x_2$ stored in the buffer through²⁶
²⁷multicast communication, where x_1 and x_2 denote the information symbols intended²⁷
²⁸for D_1 and D_2 , respectively. a_1 is the power allocation coefficient for D_1 . At the end²⁸
²⁹of a time slot, the received signal at the destinations is given by²⁹

$$y_{D_i} = \sqrt{P_{\text{R}}} g_i (\sqrt{a_1} x_1 + \sqrt{1-a_1} x_2) + w_{D_i}. \quad (5)$$

³⁰
³¹
³² When $|g_1|^2 > |g_2|^2$, according to the NOMA principle, the relay allocates more³²
³³power for D_2 , in order to balance the fairness of the system performance. Due to³³

the broadcast nature of the wireless environment, we have the instantaneous signal-to-interference-and-noise ratio (SINR) of the $R \rightarrow D_2$ link given by

$$\gamma_{D_2}^{x_2} = \frac{(1 - a_1)P_R|g_2|^2}{a_1P_R|g_2|^2 + N_0}, \quad (6)$$

where the information symbol x_1 is treated as the interference at D_2 . At D_1 , the goal is to decode information symbol x_1 of themselves. By applying the SIC principle^[4], D_1 can remove the detected information symbol x_2 from the set of received signals. From (5), the instantaneous SNR and SINR of the $R \rightarrow D_1$ link is expressed as

$$\gamma_{D_1}^{x_2 \rightarrow x_1} = \frac{(1 - a_1)P_R|g_1|^2}{a_1P_R|g_1|^2 + N_0}, \quad (7)$$

$$\gamma_{D_1}^{x_1} = \frac{a_1P_R|g_1|^2}{N_0}. \quad (8)$$

Outage probability without buffer-aided relay

In this section, we investigate the outage performance of the SWIPT NOMA downlink relaying in two cases, i.e., the overall outage probability and the outage probability for each destination.

Overall outage probability

The overall outage probability (OOP) of the system is defined as the probability that neither the source-to-relay link nor the relay-to-both destinations links is unavailable for transmission to achieve the target predefined transmission rate. For simplicity, we assume that the target transmission rates from the source to the relay and the relay to the destinations are the same and equal to r_0 . Hence, the instantaneous end-to-end capacity is $\frac{1-\alpha}{2} \log_2(1 + \gamma_{e2e}) < r_0$, and the outage event happens, the factor of $\frac{1-\alpha}{2}$ is due to the two consecutive time slots for communication between the source and the destination. Outage probability is equivalent to the probability that output SNR, γ_{e2e} , falls below a certain threshold, $\gamma_{th} = 2^{\frac{2r_0}{1-\alpha}} - 1$.

The following theorem provides the exact closed-form expression of the overall outage probability and the approximation of the outage probability of the SWIPT-NOMA downlink relaying system.

^[4]In this paper, we assume that the system is equipped with ideal successive interference cancellation (SIC) technique [26].

¹**Theorem 1** *The overall outage probability of the system when the relay knows*¹
²*both g_1 and g_2 is given by*²

$$\begin{aligned} \text{OOP} &= 1 - \frac{1}{\Omega_1} \sum_{k=0}^{\infty} \frac{(-1)^k}{k!} \left(\frac{\Psi_{\min}}{\Omega_2} \right)^k \\ &\times \left[\frac{(-1)^k}{(k-1)!} \left(\frac{1}{\Omega_1} \right)^{k-1} \text{Ei} \left(\frac{-\gamma_{\text{th}}}{\Omega_1 P_S} \right) + \frac{\exp \left(\frac{-\gamma_{\text{th}}}{\Omega_1 P_S} \right)}{\left(\frac{\gamma_{\text{th}}}{P_S} \right)^{k-1}} \sum_{j=0}^{k-2} \frac{(-1)^j \left(\frac{\gamma_{\text{th}}}{\Omega_1 P_S} \right)^j}{\prod_{\ell=0}^j (k-1-\ell)} \right], \end{aligned} \quad (9)$$

³where $\text{Ei}(x)$ denotes the exponential integral function [27], $\phi = \frac{2\alpha\eta}{1-\alpha}$, $\Psi_{\min} =$ ³
⁴ $\min \left\{ \frac{\gamma_{\text{th}}}{a_1 \phi P_S}, \frac{\gamma_{\text{th}}}{\phi P_S (1-a_1(1+\gamma_{\text{th}}))} \right\}$, and the condition $a_1 < \frac{1}{1+\gamma_{\text{th}}}$ holds.⁴

⁵At high SNR regime, the approximation of OOP is given by⁵

$$\text{OOP} \approx 1 - \exp \left(-\frac{\gamma_{\text{th}}}{\Omega_1 P_S} \right) - \sqrt{\frac{4\Psi_{\min}}{\Omega_1 \Omega_2}} \mathcal{K}_1 \left(\sqrt{\frac{4\Psi_{\min}}{\Omega_1 \Omega_2}} \right), \quad (10)$$

⁶where $\mathcal{K}_1(\cdot)$ is the first-order modified Bessel function of the second kind.⁶

⁷*Proof* Please refer to Appendix A. When the condition $a_1 < \frac{1}{1+\gamma_{\text{th}}}$ holds, which⁷
⁸means the overall outage probability does not occur, we need to allocate more power⁸
⁹to D_2 . With assumption that SIC information symbol x_2 at the D_1 is perfect, if⁹
¹⁰the relay knows the channel responses g_1 and g_2 , it can adjust the power allocation¹⁰
¹¹coefficient a_1 to balance the outage probability of the relay-to-destinations links.¹¹

¹²It should be noted that if $|g_1|^2 > |g_2|^2$, we have¹²

$$\frac{(1-a_1)P_R|g_1|^2}{a_1 P_R|g_1|^2 + N_0} > \frac{(1-a_1)P_R|g_2|^2}{a_1 P_R|g_2|^2 + N_0}. \quad (11)$$

¹³This remark is very important for analyzing the outage probability expressions in¹³
¹⁴the next part. ■¹⁴

¹⁵*Outage probability at the destination*¹⁵

¹⁶In this section, we derive a closed-form expression of the outage probability at each¹⁶
¹⁷destination. When one destination is in outage, the other can detect its correspond-¹⁷
¹⁸ing information symbol. The system may switch to the conventional OMA system.¹⁸

¹⁹However, in this case, the system performance will be degraded because the total¹⁹
²⁰power of the relay have been fixed division for D_1 and D_2 .²⁰

²¹Theorem 2 provides closed-form expressions of outage probabilities at D_1 and D_2 ,²¹
²²respectively.²²

¹**Theorem 2** The outage probability at D_1 and D_2 are given respectively in (12)¹
²and (13) below ²

$$\begin{aligned} \text{OP}_{D_1} &= 1 - \frac{1}{\Omega_1} \sum_{t=0}^{\infty} \frac{(-1)^t}{t!} \left(\frac{\mathcal{Q}_{\max}}{\Omega_2} \right)^t \\ &\times \left[\frac{(-1)^t}{(t-1)!} \left(\frac{1}{\Omega_1} \right)^{t-1} \text{Ei} \left(\frac{-\xi_1}{\Omega_1 P_S} \right) + \frac{\exp \left(\frac{-\xi_1}{\Omega_1 P_S} \right)}{\left(\frac{\xi_1}{P_S} \right)^{t-1}} \sum_{k=0}^{t-2} \frac{(-1)^k \left(\frac{1}{\Omega_1} \right)^k \left(\frac{\xi_1}{P_S} \right)^k}{\prod_{\ell=0}^k (t-1-\ell)} \right], \end{aligned} \quad (12)$$

$$\begin{aligned} \text{OP}_{D_2} &= 1 - \frac{1}{\Omega_1} \sum_{m=0}^{\infty} \frac{(-1)^m}{m!} \left(\frac{b}{\Omega_2} \right)^m \\ &\left[\frac{(-1)^m}{(m-1)!} \left(\frac{1}{\Omega_1} \right)^{m-1} \text{Ei} \left(\frac{-\xi_2}{\Omega_1 P_S} \right) + \frac{\exp \left(\frac{-\xi_2}{\Omega_1 P_S} \right)}{\left(\frac{\xi_2}{P_S} \right)^{m-1}} \sum_{q=0}^{m-2} \frac{(-1)^q \left(\frac{1}{\Omega_1} \right)^q \left(\frac{\xi_2}{P_S} \right)^q}{\prod_{v=0}^q (m-1-v)} \right]. \end{aligned} \quad (13)$$

¹⁶where $\mathcal{Q}_{\max} = \max \left\{ \frac{\xi_1}{a_1 \phi P_S}, \frac{\xi_1}{\phi P_S (1 - a_1 (1 + \xi_1))} \right\}$, $b = \frac{\xi_2}{\phi P_S (1 - a_1 (1 + \xi_2))}$, $\xi_1 = 2^{\frac{2r_1}{1-\alpha}} - 1$,
¹⁷ $\xi_2 = 2^{\frac{2r_2}{1-\alpha}} - 1$, and the condition $a_1 \leq \frac{1}{1+\xi_i}$, $i \in \{1, 2\}$ holds. r_1 and r_2 are the
¹⁸target transmission rates at D_1 and D_2 , respectively. ¹⁸

²⁰*Proof* To obtain the outage probability expression of D_1 and D_2 , we first analyze
²¹the instantaneous SINR and SNR of D_1 and D_2 . It should be noted that in order
²²to prove this theorem we assume that $|g_1|^2 > |g_2|^2$. Please refer to Appendix B. \blacksquare
²³

²⁴**Optimal power allocation to minimize the outage probability** ²⁴

²⁵In this section, we study the power allocation problem to minimize the outage
²⁶probability of the EH-NOMA system. ²⁶

²⁸**Theorem 3** The optimal power allocation coefficient a_1^* to minimize the outage
²⁹probability in the EH-NOMA system is given by ²⁹

$$a_1^* = \frac{1}{2 + \gamma_{\text{th}}}. \quad (14)$$

³²*Proof* To obtain the minimum OOP and minimum OP_{D_1} in (9) and (12) we formu-
³³late the outage probability minimization problems P1 and P2 as follows: ³³

$$\begin{aligned}
& \text{P1 : } \min_{a_1} \text{ OOP,} \\
& \text{s.t. } 0 \leq a_1 \leq 1, \\
& a_1 < \frac{1}{1 + \gamma_{\text{th}}}.
\end{aligned}
\tag{15}$$

and

$$\begin{aligned}
& \text{P2 : } \min_{a_1} \text{ OP}_{D_1}, \\
& \text{s.t. } 0 \leq a_1 \leq 1, \\
& a_1 < \frac{1}{1 + \xi_1}.
\end{aligned}
\tag{17}$$

The condition $a_1 < \frac{1}{1 + \gamma_{\text{th}}}$ indicates that the outage event does not occur, i.e., the outage probability is less than one. The problems P1 and P2 are equivalent to maximizing $\Psi_{\min} = \min \left\{ \frac{\gamma_{\text{th}}}{a_1 \phi P_S}, \frac{\gamma_{\text{th}}}{\phi P_S (1 - a_1 (1 + \gamma_{\text{th}}))} \right\}$ and $\mathcal{Q}_{\max} = \max \left\{ \frac{\xi_1}{a_1 \phi P_S}, \frac{\xi_1}{\phi P_S (1 - a_1 (1 + \xi_1))} \right\}$. In addition, we assume that $\gamma_{\text{th}} = \xi_1$, which means that the data rate of D_1 is equal to the system data rate.

$$\begin{aligned}
& \text{P1 : } \max \left\{ \min \left(\frac{1}{a_1}, \frac{1}{1 - a_1 (1 + \gamma_{\text{th}})} \right) \right\}, \\
& \text{s.t. } 0 \leq a_1 \leq 1, \\
& a_1 < \frac{1}{1 + \gamma_{\text{th}}}.
\end{aligned}
\tag{19}$$

$$\begin{aligned}
& \text{P2 : } \max \left\{ \max \left(\frac{1}{a_1}, \frac{1}{1 - a_1 (1 + \gamma_{\text{th}})} \right) \right\}, \\
& \text{s.t. } 0 \leq a_1 \leq 1, \\
& a_1 < \frac{1}{1 + \gamma_{\text{th}}}.
\end{aligned}
\tag{21}$$

We consider two cases of the objective functions P1 and P2 under the following conditions:

$$a_1^* = \left(\frac{1}{a_1} \leq \frac{1}{1 - a_1 (1 + \gamma_{\text{th}})} \right) \cup \left(\frac{1}{a_1} \geq \frac{1}{1 - a_1 (1 + \gamma_{\text{th}})} \right).
\tag{22}$$

These problems can be solved similarly as those in [28, 4.1].

¹ After some mathematical manipulations, we obtain the optimal power allocation¹
² coefficient a_1^* as shown in Theorem (3). It should be noticed that the power allo-²
³ cation for D_1 is considered, subject to the condition that the outage probability of³
⁴ D_2 certainly occurs^[5]. ■⁴

⁶Analysis of the outage probability with buffer-aided relay ⁶

⁷In this section, we investigate the outage probability of the system where the buffer⁷
⁸ aid is employed at the relay node. For convenience, we assume that the source node⁸
⁹ always has data to transmit. We also consider the number of transmitted symbols⁹
¹⁰ as the number of transmitted packets. The relay chooses a node to transmit (source¹⁰
¹¹ or relay) in a given time slot. To perform this, the information of the outage states¹¹
¹² of the links $S \rightarrow R$ and $R \rightarrow D$ is required. Therefore, the system uses one bit for¹²
¹³ the feedback information from the destination to the relay. This information helps¹³
¹⁴ R known if the link $R \rightarrow D$ is in outage or not. One bit which feedbacks from the¹⁴
¹⁵ relay to source is used to control the source in the transmit or silent mode. The¹⁵
¹⁶ source transmits the packets to the relay. Then, the relay decodes the packets and¹⁶
¹⁷ stores the decoded packets in its buffer. After that, the relay transmits the packets¹⁷
¹⁸ to the destination node. If the source is selected to transmit but the link $S \rightarrow R$ is¹⁸
¹⁹ in outage, the source remains silent and the outage occurs. Similarly, if the relay¹⁹
²⁰ is selected to transmit but the link $R \rightarrow D$ is in outage, the relay remains silent²⁰
²¹ and the outage occurs. Therefore, the system performance will be improved with²¹
²² the help of a buffer-aided relay. Unlike the case of without a buffer-aided relay,²²
²³ in this case the outage event is defined as the probability that the relay does not²³
²⁴ receive and transmit. In other words, the relay remains silent. To describe the state²⁴
²⁵ transition of the buffer-aided relay, we denote the outage events of $S \rightarrow R$ link and²⁵
²⁶ $R \rightarrow D$ link by \mathcal{O}_{SR} and \mathcal{O}_{RD} , respectively. When the links are not in outage, the²⁶
²⁷ probabilities are: $1 - \mathcal{O}_{SR} = \bar{\mathcal{O}}_{SR}$ and $1 - \mathcal{O}_{RD} = \bar{\mathcal{O}}_{RD}$, respectively. In addition,²⁷
²⁸ the relay decision scheme is described as in Table 2. In Table 2, ‘SR’ denotes the²⁸
²⁹ link from the source to the relay nodes, ‘RD’ refers to the link from the relay to²⁹
³⁰ the destination nodes; ‘ l ’ and ‘ L ’ respectively represents the packets stored in the³⁰

³¹ ^[5] Moreover, if the relay knows the channel gains g_1 and g_2 and the total power factor³¹
³² $a_1 + a_2$ is equal to one, we can allocate power for D_1 and D_2 by adapting to channel³²
³³ gains, i.e. $a_1 = |g_2|^2 / (|g_1|^2 + |g_2|^2)$, to ensure the fairness of the outage performance.³³

Table 2 The Relay Decision Scheme

Case	SR	RD	l	Relay	The outage
A	0	0		Silent	$\mathcal{O}_{\text{SR}}\mathcal{O}_{\text{RD}}$
B	0		$l = 0$	Silent	\mathcal{O}_{SR}
C		0	$l = L$	Silent	\mathcal{O}_{RD}
D	1	0	$l < L$	Receive	$\bar{\mathcal{O}}_{\text{SR}}\mathcal{O}_{\text{RD}}$
E	0	1	$l > 0$	Transmit	$\mathcal{O}_{\text{SR}}\bar{\mathcal{O}}_{\text{RD}}$
F	1	1	$l \geq 2$	Transmit	$\bar{\mathcal{O}}_{\text{SR}}\bar{\mathcal{O}}_{\text{RD}}$
G	1	1	$l \leq 1$	Receive	$\bar{\mathcal{O}}_{\text{SR}}\bar{\mathcal{O}}_{\text{RD}}$

buffer and the buffer size at the relay node. ‘Relay’ denotes the decision of the relay node (silent, receive or transmit), ‘ \mathcal{OP} ’ is the outage probability of the considered system. It is noted that in Table 2, the outage and non-outage links are indicated by ‘0’ and ‘1’, respectively.

To calculate the \mathcal{OP} of the system, from the Table 2, we build the Markov chain. We start at the initial state $l = 0$ (i.e. when the buffer is empty). If the link SR is in outage which means the source does not transmit, then, the buffer will be empty. In other words, the buffer state moves from $l = 0$ to $l = 0$ with probability of \mathcal{O}_{SR} (Case B in Table 2). When the link SR is not in outage, we consider two cases. The first case is when the link RD is in outage (Case D). Consequently, the relay receives the signal, making the buffer state moves from $l = 0$ to $l = 1$ with probability of $(1 - \mathcal{O}_{\text{SR}})\mathcal{O}_{\text{RD}}$. The second case is when the link RD is not in outage (Case G). The relay receives the signal, making the buffer state moves from $l = 0$ to $l = 1$ with probability of $(1 - \mathcal{O}_{\text{SR}})(1 - \mathcal{O}_{\text{RD}})$. Combining these two cases, the buffer state moves from $l = 0$ to $l = 1$ with probability of $1 - \mathcal{O}_{\text{SR}}$. Similarly, we can obtain the probability of moving to the next state. From here, we have the Markov chain showing the state transitions as depicted in Fig. 2. When the buffer is empty ($l = 0$), it stays empty with probability of \mathcal{O}_{SR} (case B) and receives a packet with probability of $1 - \mathcal{O}_{\text{SR}}$ (case D, G). When the buffer has one packet ($l = 1$), it stays in the current state with probability of $\mathcal{O}_{\text{SR}}\mathcal{O}_{\text{RD}}$ if the relay does not receive and transmit (case A). If the relay receives one packet, it moves to the new state ($l = 2$) with probability of $1 - \mathcal{O}_{\text{SR}}$ (case D, G) and is back to the initial state ($l = 0$) with probability of $\mathcal{O}_{\text{SR}}(1 - \mathcal{O}_{\text{RD}})$ (case E). When the buffer has l packets ($2 \leq l \leq L - 1$), it stays in this state with probability of $\mathcal{O}_{\text{SR}}\mathcal{O}_{\text{RD}}$ (case A),

¹receives one packet with probability of $(1 - \mathcal{O}_{\text{SR}})\mathcal{O}_{\text{RD}}$ (case D), and transmits one¹
²packet with probability of $1 - \mathcal{O}_{\text{RD}}$ (case E, F). If the buffer is full, which means²
³that it has L packets, it remains the same state with probability of \mathcal{O}_{RD} (case C)³
⁴and transmits one packet with probability of $1 - \mathcal{O}_{\text{RD}}$ (case E, F).⁴

⁵ From Table 2 and the presented Markov chain, the outage probability of the⁵
⁶system is calculated as⁶

$$\begin{aligned} \mathcal{OP} &= \mathcal{O}_{\text{SR}} \Pr\{l = 0\} + \mathcal{O}_{\text{RD}} \Pr\{l = L\} \\ &+ \mathcal{O}_{\text{SR}}\mathcal{O}_{\text{RD}}(1 - \Pr\{l = 0\} - \Pr\{l = L\}), \end{aligned} \quad (24)$$

¹¹where $\Pr\{l = 0\}$ and $\Pr\{l = L\}$ are the probabilities of the events that the buffer¹¹
¹²is empty and full, respectively. To derive the \mathcal{OP} of the system in (24), we define a¹²
¹³state transition matrix \mathbf{A} with size of $(L + 1) \times (L + 1)$ of the Markov chain, where¹³
¹⁴ \mathbf{A}_{ij} denotes the element of the i th row and j th column of the matrix \mathbf{A} . It should¹⁴
¹⁵be reminded that \mathbf{A}_{ij} refers to the probability of moving from state i at time t to¹⁵
¹⁶state j at time $t + 1$, i.e.,¹⁶

$$\mathbf{A}_{ij} = \Pr\{l_{t+1} = j | l_t = i\}. \quad (25)$$

¹⁹For the case of $L = 5$, matrix \mathbf{A} is expressed as follows¹⁹

$$\mathbf{A} = \begin{pmatrix} \mathcal{O}_{\text{SR}} & \mathcal{O}_{\text{SR}} & 0 & 0 & 0 & 0 \\ \mathcal{O}_{\text{SR}}\mathcal{O}_{\text{RD}} & \mathcal{O}_{\text{SR}}\mathcal{O}_{\text{RD}} & \mathcal{O}_{\text{SR}} & 0 & 0 & 0 \\ 0 & \mathcal{O}_{\text{RD}} & \mathcal{O}_{\text{SR}}\mathcal{O}_{\text{RD}} & \mathcal{O}_{\text{SR}}\mathcal{O}_{\text{RD}} & 0 & 0 \\ 0 & 0 & \mathcal{O}_{\text{RD}} & \mathcal{O}_{\text{SR}}\mathcal{O}_{\text{RD}} & \mathcal{O}_{\text{SR}}\mathcal{O}_{\text{RD}} & 0 \\ 0 & 0 & 0 & \mathcal{O}_{\text{RD}} & \mathcal{O}_{\text{SR}}\mathcal{O}_{\text{RD}} & \mathcal{O}_{\text{SR}}\mathcal{O}_{\text{RD}} \\ 0 & 0 & 0 & 0 & \mathcal{O}_{\text{RD}} & \mathcal{O}_{\text{RD}} \end{pmatrix}. \quad (26)$$

²⁶We should note that matrix \mathbf{A} is not symmetric because the states are not sym-²⁶
²⁷metric and the number of links to other states is not the same, leading to the²⁷
²⁸transition probabilities are not the same. Then, the stationary distribution $\boldsymbol{\pi}$ of the²⁸
²⁹Markov chain is expressed as²⁹

$$\boldsymbol{\pi} = (\mathbf{A} - \mathbf{I} + \mathbf{B})^{-1}\mathbf{b}, \quad (27)$$

³²where \mathbf{I} is an identity matrix, \mathbf{B} is an $(L + 1) \times (L + 1)$ matrix with all elements³²
³³equal to 1, and $\mathbf{b} = (1 \ 1 \ \dots \ 1)^T$.³³

¹**Theorem 4** *With the buffer-aided relaying, the outage probability of the system*¹
²*becomes*²

$$\mathcal{OP} = \sum_{i=1}^{L+1} \pi_i \mathbf{A}_{ii}. \quad (28)$$

⁵To determine the state transit matrix \mathbf{A} , we need to derive \mathcal{O}_{SR} and \mathcal{O}_{RD} . We
⁶assume that the minimum data transmission rate from $\text{S} \rightarrow \text{R}$ is r_0 , then the outage
⁷probability of $\text{S} \rightarrow \text{R}$ link is defined as follows⁷

$$\mathcal{O}_{\text{SR}} = \Pr \left(\frac{1-\alpha}{2} \log_2(1 + \gamma_{\text{R}}) < r_0 \right) = 1 - \exp \left(-\frac{\gamma_{\text{th}}}{\Omega_{\text{SR}} P_{\text{S}}} \right). \quad (29)$$

¹¹According to the SIC principle, if D_1 is able to remove x_2 from its received signal,¹¹
¹²the outage probability of the link from the relay to the destinations nodes is given¹²
¹³by¹³

$$\mathcal{O}_{\text{RD}} = \Pr \left(\frac{1-\alpha}{2} \log_2 \left(1 + \max \{ \gamma_{\text{D}_1}^{x_1}, \gamma_{\text{D}_2}^{x_2} \} \right) < r_0 \right). \quad (30)$$

¹⁶After some manipulations, we have¹⁶

$$\begin{aligned} \mathcal{O}_{\text{RD}} = 1 - & \sqrt{\frac{4\mathcal{A}}{\Omega_{\text{SR}}}} K_1 \left(\sqrt{\frac{4\mathcal{A}}{\Omega_{\text{SR}}}} \right) - \sqrt{\frac{4\mathcal{B}}{\Omega_{\text{SR}}}} K_1 \left(\sqrt{\frac{4\mathcal{B}}{\Omega_{\text{SR}}}} \right) \\ & + \sqrt{\frac{4(\mathcal{A} + \mathcal{B})}{\Omega_{\text{SR}}}} K_1 \left(\sqrt{\frac{4(\mathcal{A} + \mathcal{B})}{\Omega_{\text{SR}}}} \right), \end{aligned} \quad (31)$$

²²where $\mathcal{A} = \frac{\gamma_{\text{th}}}{\Omega_{\text{RD}_1} a_1 \phi P_{\text{S}}}$ and $\mathcal{B} = \frac{\gamma_{\text{th}}}{\Omega_{\text{RD}_2} \phi P_{\text{S}} (1 - a_1 (1 + \gamma_{\text{th}}))}$.²²

²³For the detailed derivations of \mathcal{O}_{RD} , please refer to Appendix C.²³

²⁴Average packet delay of the buffer-aided relay system²⁴

²⁵In this section, the average packet delay of the system is considered. This delay
²⁶includes the average packet delay at the source and the relay.²⁶

²⁷The average packet delay at the source is determined as²⁷

$$\mathcal{D}_{\text{S}} = \frac{1 + \mathcal{OP}}{1 - \mathcal{OP}}. \quad (32)$$

³¹while the average packet delay at the relay is expressed as³¹

$$\mathcal{D}_{\text{R}} = \frac{2}{1 - \mathcal{OP}} \sum_{i=2}^{L+1} \pi_i (i - 1). \quad (33)$$

¹Therefore, the average packet delay of the system is given by

²

$$\mathcal{D} = \mathcal{D}_S + \mathcal{D}_R. \quad (34)$$

⁴

⁵Results and Discussion ⁵

⁶In this section, detailed numerical results are provided to illustrate the impact of ⁶
⁷power allocation on the performance of SWIPT-NOMA system in terms of the OP ⁷
⁸and the EC. For comparison, we also provide the performance of the SWIPT-OMA ⁸
⁹system with the same parameters. Configurations and parameters of the system are ⁹
¹⁰explained as follows. D_1 is closer to the relay nodes than D_2 . Hence, we need to ¹⁰
¹¹allocate more power to D_2 than D_1 to ensure the user fairness. The optimal power ¹¹
¹²allocation coefficient is derived in Theorem 3. The power allocation coefficient for D_1 ¹²
¹³is fixed at $a_1 = 0.3$ and that for D_2 is $1 - a_1$. The energy harvesting fraction $\alpha = 0.3$ ¹³
¹⁴and the energy conversion efficiency $\eta = 0.8$. The system data rate $r_1 = 1$ while ¹⁴
¹⁵ $r_2 = r_0 = 0.5[b/s/Hz]$. The obtained numerical results show that the optimal power ¹⁵
¹⁶allocation can increase the system performance and the NOMA scheme significantly ¹⁶
¹⁷improves the spectrum utilization. ¹⁷

¹⁸ Fig. 4 illustrates the overall outage probability in terms of the average SNR in dB ¹⁸
¹⁹for two cases, i.e., with and without buffer-aided data at the relay node. As observed ¹⁹
²⁰from Fig. 4, the overall outage probability of the system (including the outage events ²⁰
²¹of both D_1 and D_2) in the case of optimal power allocation outperforms the case of ²¹
²²fixed power allocation. From Fig 4, we can see that the benefit of the optimal power ²²
²³allocation in terms of the overall OP compared with the fixed power allocation is ²³
²⁴not significant. This is because we have choose the fixed power allocation $a_1 = 0.3$, ²⁴
²⁵which is the approximate of the optimal power allocation a_1^* (when $\alpha = 0.3$ and ²⁵
²⁶ $r = 0.5$ returns $a_1^* = 0.2709$). In order to achieve the fairness of the overall system ²⁶
²⁷performance, the transmitter needs to allocate power according to the channel gains ²⁷
²⁸of $R \rightarrow D_1$ and $R \rightarrow D_2$. Furthermore, the approximation results calculated from ²⁸
²⁹(10) are very close to the exact results obtained from (9), especially at high SNR ²⁹
³⁰regime. Therefore, we can use (10) to calculate the OOP of the system easily. We ³⁰
³¹can also see in Fig. 4 that the diversity order of the system with buffer aid is equal ³¹
³²to 2. Meanwhile, for the case of without buffer-aided relaying, the diversity order ³²
³³of the system is equal to 1. Hence, employing data buffer at the relay leads to the ³³

¹reduction of OP, but it trades off with the packet delay. We can also see that the ¹
²analytical results agree well with the simulation results. ²

³ Fig. 5 plots the overall OP and the OP of D_1 and D_2 , respectively. The optimal ³
⁴power allocation coefficient as presented in Theorem 3 is used for both cases with ⁴
⁵and without buffer aided data at the relay node. We can see that the outage per- ⁵
⁶formance of D_1 is better than D_2 . This is because the distance from the relay to ⁶
⁷ D_2 is longer than that from the relay to D_1 [29] ^[6] The overall OP is calculated ⁷
⁸as the probability of the events that both D_1 and D_2 cannot decode their symbols ⁸
⁹successfully. The simulation and analytical results are in excellent match, validating ⁹
¹⁰the correctness of the closed-form expressions of (9), (12) and (13). From Fig. 5, we ¹⁰
¹¹can observe that the joint outage events of D_1 and D_2 are less than each individual ¹¹
¹²outage event of D_1 and D_2 . This is suitable in practice where the probability that ¹²
¹³both D_1 and D_2 are in outage is always less than the probability that D_1 or D_2 is ¹³
¹⁴in outage. ¹⁴

¹⁵ The OP_{D_1} and OP_{D_2} are shown in Fig. 6 and Fig. 7, respectively. The power ¹⁵
¹⁶allocation coefficient is fixed at $a_1 = 0.3$ when investigating the outage performance ¹⁶
¹⁷of both OP_{D_1} and OP_{D_2} . Moreover, we also conduct the optimal power allocation ¹⁷
¹⁸for the relay as described in the Theorem 3. Again, we can see that the analytical ¹⁸
¹⁹results are in excellent agreement with the simulation results. From these figures ¹⁹
²⁰we can see that the NOMA system with the optimal power allocation has better ²⁰
²¹outage performance than the OMA system. In the case of fixed power allocation, the ²¹
²²outage performance of D_1 is better than the OMA system, but outage performance ²²
²³of D_2 is worse than the OMA system. However, the NOMA system provides better ²³
²⁴spectral efficiency because two users are served simultaneously. Different from Fig 4, ²⁴
²⁵the gap of the curves plotted in Fig 6 and Fig 7 is more significant for the two cases ²⁵
²⁶with fixed and optimal power allocation. The reason is that the probability of the ²⁶
²⁷event that both D_1 and D_2 are in outage is less than the probability of each event ²⁷
²⁸that D_1 or D_2 is in outage. ²⁸

²⁹ Fig. 8 depicts the effect of the power allocation coefficient on the OP. It should ²⁹
³⁰be noted that we only define the coefficient for D_2 while the coefficient for D_1 is ³⁰
³¹

^[6]The signal power of far-field RF transmission is reduced according to the mu- ³¹
³²tual distance between receiver and transmitter, specifically, $20dB$ per decade of the ³²
³³distance. ³³

¹derived from the condition $a_2 = 1 - a_1$. As shown in Fig. 8, different data rates r_2 ¹
²exhibit different minimum values of the OP. In this figure we can see that when the ²
³rate transmission r_2 is reduced, the power allocation coefficient for D₂ decreases to ³
⁴get better system performance for the fairness of the outage performance of D₁ and ⁴
⁵D₂. ⁵

⁶ The system ergodic capacity is shown in Fig. 9. According to the NOMA theory, ⁶
⁷the ergodic capacity of the system is the summation of the ergodic capacity of ⁷
⁸all users. Let β and $(1 - \beta)$ [Hz] denote the bandwidth assigned for D₁ and the ⁸
⁹remaining bandwidth assigned for D₂, where $(0 \leq \beta \leq 1)$. Using [30, eq. (7.4)], the ⁹
¹⁰sum capacity of the OMA system is given by ¹⁰

$$\begin{aligned}
 C_{OMA} &= \frac{1 - \alpha}{2} \beta \log_2 \left(1 + \min \left\{ \frac{P_S |h_1|^2}{\beta}, \frac{P_R |g_1|^2}{\beta} \right\} \right) \\
 &\quad + \frac{(1 - \alpha)(1 - \beta)}{2} \log_2 \left(1 + \min \left\{ \frac{P_S |h_1|^2}{1 - \beta}, \frac{P_R |g_2|^2}{1 - \beta} \right\} \right). \quad (35)
 \end{aligned}$$

¹⁷ From Fig. 9, we can see that the ergodic capacity of the NOMA system is better ¹⁷
¹⁸than the OMA system. Additionally, the ergodic capacity of D₁ in our proposed ¹⁸
¹⁹system is better than that in the OMA system. However, the ergodic capacity of D₂ ¹⁹
²⁰is not higher than OMA system due to the poor channel condition from the relay to ²⁰
²¹the D₂. Therefore, the advantage of the NOMA system is to improve the capacity ²¹
²²significantly. We also see that the analytical results are very close to the simulation ²²
²³results because we use an approximation of the CDF of X_2 as given in (56). ²³

²⁵ Conclusion ²⁵

²⁶ In this paper, we proposed a NOMA cooperative relaying network with and with- ²⁶
²⁷out data buffer-aided relay. In addition, the relay node harvests the energy from the ²⁷
²⁸source using the time-switching mechanism. We focus on deriving the OP and er- ²⁸
²⁹godic capacity of the system over Rayleigh fading channels. Moreover, we proposed ²⁹
³⁰a power allocation scheme at the relay node which aims to reduce the feedback cost. ³⁰
³¹Numerical results of the OP and capacity showed that the proposed NOMA down- ³¹
³²link relaying system significantly outperformed the OMA system. The data buffer ³²
³³aid employed at the relay helps improve the performance of the system. However, ³³

¹the proposed system trades off with the permissible packet delay. All closed-form¹
²expressions derived in this paper were verified by the Monte-Carlo simulations. ²

³ The impact of fixed and optimal power allocation on the performance of EH-³
⁴NOMA downlink relaying network was also investigated. In this model, all nodes⁴
⁵are equipped with a single antenna. However, it can be developed for multiple⁵
⁶antenna systems. ⁶

⁷ Our proposed relaying network can achieve two goals: (i) the energy efficiency is⁷
⁸improved by the harvested energy from the ambient RF environment. This idea can⁸
⁹be applied to sensor nodes in wireless body area networks for healthcare and other⁹
¹⁰medical applications, (ii) the spectrum utilizing efficiency is superior to that of the¹⁰
¹¹OMA system. It is a promising application which can enhance the performance of¹¹
¹²the 5G networks and the wireless sensor and healthcare networks. ¹²

¹⁴**Appendix A** ¹⁴

¹⁵The goal of this appendix is to provide the overall OP of the SWIPT-NOMA system¹⁵
¹⁶over Rayleigh fading channels. ¹⁶

¹⁷The overall OP of the system can be expressed as ¹⁷

$$\sup{19} \quad \text{OOP} = \underbrace{\Pr(\gamma_R \leq \gamma_{\text{th}})}_{\text{OP}_1} + \underbrace{\Pr(\gamma_R > \gamma_{\text{th}}, \max(\gamma_{D_1}^{x_1}, \gamma_{D_2}^{x_2}, \gamma_{D_1}^{x_2 \rightarrow x_1}) \leq \gamma_{\text{th}})}_{\text{OP}_2}. \quad (36) \sup{19}$$

²²From (2), we obtain the closed-form expression of the first term of (36) as ²²

$$\sup{23} \quad \text{OP}_1 = \Pr\left(|h_1|^2 \leq \frac{\gamma_{\text{th}}}{P_S}\right) = 1 - \exp\left(-\frac{\gamma_{\text{th}}}{\Omega_1 P_S}\right). \quad (37) \sup{23}$$

²⁵To obtain the closed-form expression of the second term of (36), we rewrite the²⁵
²⁶second term of (36) as follows ²⁶

$$\sup{29} \quad \text{OP}_2 = \Pr(\gamma_R > \gamma_{\text{th}}, \gamma_{D_1}^{x_1} \leq \gamma_{\text{th}}, \gamma_{D_1}^{x_2 \rightarrow x_1} < \gamma_{\text{th}}, \gamma_{D_2}^{x_2} < \gamma_{\text{th}}). \quad (38) \sup{29}$$

³¹From (38) and the condition in (11), i.e. $\gamma_{D_1}^{x_2 \rightarrow x_1} > \gamma_{D_2}^{x_2}$, we have ³¹

$$\sup{33} \quad \text{OP}_2 = \Pr(\gamma_R > \gamma_{\text{th}}, \gamma_{D_1}^{x_1} \leq \gamma_{\text{th}}, \gamma_{D_1}^{x_2 \rightarrow x_1} < \gamma_{\text{th}}). \quad (39) \sup{33}$$

¹ By substituting (2), (7), and (8) into (39) and denoting $|h_1|^2 = X$, $|g_1|^2 = Y$ and ² $|g_2|^2 = Z$, we obtain: ²

$$\text{OP}_2 = \Pr \left(X > \frac{\gamma_{\text{th}}}{P_S}, XY \leq \frac{\gamma_{\text{th}}}{a_1 \phi P_S}, XY < \frac{\gamma_{\text{th}}}{\phi P_S (a_2 - a_1 \gamma_{\text{th}})} \right). \quad (40)$$

³ As can be seen from (40), the outage always occurs if $a_1 \geq \frac{1}{1+\gamma_{\text{th}}}$. Thus, allocating ⁴more power to the D₂ is required so that $1 - a_1(1 + \gamma_{\text{th}}) > 0$ always holds. The ⁵condition $a_1 < \frac{1}{1+\gamma_{\text{th}}}$ is used throughout this paper. For simplicity, we can rewrite ⁶(40) as ⁶

$$\text{OP}_2 = \Pr \left(X > \frac{\gamma_{\text{th}}}{P_S}, XY \leq \Psi_{\min} \right), \quad (41)$$

⁷ where $\Psi_{\min} = \min \left\{ \frac{\gamma_{\text{th}}}{a_1 \phi P_S}, \frac{\gamma_{\text{th}}}{\phi P_S (1 - a_1 (1 + \gamma_{\text{th}}))} \right\}$. ⁷

⁸ Based on the conditional probability [31] and the assumption that the channel ⁹gains have exponential distributions, we have ⁹

$$\text{OP}_2 = \int_{\frac{\gamma_{\text{th}}}{P_S}}^{\infty} F_Y \left(\frac{\Psi_{\min}}{x} \right) f_X(x) dx = \int_{\frac{\gamma_{\text{th}}}{P_S}}^{\infty} \left[1 - \exp \left(-\frac{\Psi_{\min}}{\Omega_2 x} \right) \right] f_X(x) dx, \quad (42)$$

¹⁰ where $F_Y(y) = 1 - \exp \left(-\frac{y}{\Omega_2} \right)$ and $f_X(x) = \frac{1}{\Omega_1} \exp \left(-\frac{x}{\Omega_1} \right)$ are the CDF of X and ¹¹the PDF of Y , respectively. ¹¹

¹² Substituting PDF of X into (42) yields ¹²

$$\text{OP}_2 = \exp \left(-\frac{\gamma_{\text{th}}}{\Omega_1 P_S} \right) - \frac{1}{\Omega_1} \int_{\frac{\gamma_{\text{th}}}{P_S}}^{\infty} \exp \left(-\frac{\Psi_{\min}}{\Omega_2 x} - \frac{x}{\Omega_1} \right) dx. \quad (43)$$

¹³ By using the Taylor series expansions of the exponential function and after some ¹⁴manipulations on (43) using [27, 3.351.4], we obtain the second term of (36) as ¹⁵given in (44) below. ¹⁵

$$\text{OP}_2 = \exp \left(-\frac{\gamma_{\text{th}}}{\Omega_1 P_S} \right) - \frac{1}{\Omega_1} \sum_{k=0}^{\infty} \frac{(-1)^k}{k!} \left(\frac{\Psi_{\min}}{\Omega_2} \right)^k$$

$$\left[\frac{(-1)^k}{(k-1)!} \left(\frac{1}{\Omega_1} \right)^{k-1} \text{Ei} \left(\frac{-\gamma_{\text{th}}}{\Omega_1 P_S} \right) + \frac{\exp \left(\frac{-\gamma_{\text{th}}}{\Omega_1 P_S} \right)}{\left(\frac{\gamma_{\text{th}}}{P_S} \right)^{k-1}} \sum_{j=0}^{k-2} \frac{(-1)^j \left(\frac{\gamma_{\text{th}}}{\Omega_x P_S} \right)^j}{\prod_{\ell=0}^j (k-1-\ell)} \right]. \quad (44)$$

¹ We obtain the OP expression of the system by combining (37) and (44). When¹

²the transmission power is high, we have $\gamma_{\text{th}} \ll P_S$. Then, (43) can be approximated²

³as³

⁴⁴

⁵⁵

$$\text{OP}_2 = 1 - \frac{1}{\Omega_1} \int_0^\infty \exp\left(-\frac{\Psi_{\min}}{\Omega_2 x} - \frac{x}{\Omega_1}\right) dx. \quad (45)$$

⁷⁷

⁸ From (45), by using [27, 3.324.1], i.e. $\int_0^\infty e^{-\frac{\beta}{4x} - \gamma x} = \sqrt{\frac{\beta}{\gamma}} K_1(\sqrt{\beta\gamma})$, and after⁸

⁹some mathematical manipulations, we have the approximation of (10) at high SNR⁹

¹⁰regime. The proof of Theorem 1 is completed. ¹⁰

¹¹¹¹

¹²Appendix B ¹²

¹³Due to the imperfect detection at the relay node, it may forward wrong decoded¹³

¹⁴signals to D₁ and D₂ and cannot apply SIC technique on symbol x_2 at the D₁. Hence¹⁴

¹⁵similar to [32], for any modulation scheme, the dual-hop of the links S → R → D₁¹⁵

¹⁶or S → R → D₂ can be modeled as an equivalent one-hop channel whose output¹⁶

¹⁷SINR \mathcal{X}_i , $i \in \{1, 2\}$ at high SNR regime can be tightly approximated. ¹⁷

¹⁸ Let denote \mathcal{X}_1 and \mathcal{X}_2 the SINRs obtained at D₁ and D₂, respectively [5]. ¹⁸

¹⁹¹⁹

$$\mathcal{X}_1 = \min(\gamma_R, \gamma_{D_1}^{x_1}, \gamma_{D_1}^{x_2 \rightarrow x_1}), \quad (46)$$

²¹²¹

²²²²

$$\mathcal{X}_2 = \min(\gamma_R, \gamma_{D_2}^{x_2}). \quad (47)$$

²⁴²⁴

²⁵ To find the OP of D₁, from (46), we have the OP expression of D₁ as ²⁵

²⁶²⁶

$$\begin{aligned} \text{OP}_{D_1} &= 1 - \Pr(\gamma_R > \xi_1, \gamma_{D_1}^{x_1} > \xi_1, \gamma_{D_1}^{x_2 \rightarrow x_1} > \xi_1) \\ &= 1 - \Pr\left(X > \frac{\xi_1}{P_S}, XY \geq \mathcal{Q}_{\max}\right), \end{aligned} \quad (48)$$

²⁹²⁹

³⁰ where $\mathcal{Q}_{\max} = \max\left\{\frac{\xi_1}{a_1 \phi P_S}, \frac{\xi_1}{\phi P_S(1-a_1(1+\xi_1))}\right\}$. ³⁰

³¹ By using the conditional probability [31], we can rewrite (48) as ³¹

³²³²

$$\text{OP}_{D_1} = 1 - \int_{\frac{\xi_1}{P_S}}^\infty \left[1 - F_Y\left(\frac{\mathcal{Q}_{\max}}{x}\right)\right] f_X(x) dx. \quad (49)$$

³³³³

¹ Since the CDF and PDF of X and Y are exponential distribution functions, we¹
² have²

$$\text{OP}_{D_1} = 1 - \frac{1}{\Omega_x} \int_{\frac{\xi_1}{P_S}}^{\infty} \exp\left(-\frac{Q_{\max}}{\Omega_2 x}\right) \exp\left(-\frac{x}{\Omega_1}\right) dx. \quad (50)$$

³
⁴ By using the Taylor series expansions of the exponential function and after some⁴
⁵ manipulations on (50) using [27, 3.351.4], we have the expression of the OP of D_1 ⁵
⁶ as presented in (12).⁶

⁷ Next, we calculate the OP expression at D_2 . With the given SINR at the D_2 and⁷
⁸ the notation $\mathcal{X}_2 = \min(\gamma_R, \gamma_{D_2}^{x_2})$, we have⁸

$$\text{OP}_{D_2} = \Pr(\min(\gamma_R, \gamma_{D_2}^{x_2}) \leq \xi_2) = 1 - \Pr(\gamma_R > \xi_2, \gamma_{D_2}^{x_2} > \xi_2). \quad (51)$$

⁹ Substituting (2) and (6) into (51) yields⁹

$$\begin{aligned} \text{OP}_{D_2} &= 1 - \Pr\left(\frac{P_S |h_1|^2}{N_0} > \xi_2, \frac{a_2 P_R |g_2|^2}{a_1 P_R |g_2|^2 + N_0} > \xi_2\right) \\ &= 1 - \Pr\left(X > \frac{\xi_2}{P_S}, XZ > \frac{\xi_2}{\phi P_S (a_2 - a_1 \xi_2)}\right) \end{aligned} \quad (52)$$

¹⁰ Then, by applying similar calculations in Appendix A we can obtain the OP of¹⁰
¹¹ D_2 as¹¹

$$\begin{aligned} \text{OP}_{D_2} &= 1 - \int_{\frac{\xi_2}{P_S}}^{\infty} \left[1 - F_Z\left(\frac{\xi_2}{x \phi P_S (a_2 - a_1 \xi_2)}\right)\right] f_X(x) dx \\ &= 1 - \frac{1}{\Omega_1} \int_{\frac{\xi_2}{P_S}}^{\infty} \exp\left(-\frac{b}{x \Omega_2}\right) \exp\left(-\frac{x}{\Omega_2}\right) dx. \end{aligned} \quad (53)$$

¹² By using the Taylor series expansions of the exponential function $\exp\left(-\frac{b}{x \Omega_2}\right) =$ ¹²
¹³ $\sum_{t=0}^{\infty} \frac{(-1)^t}{t!} \left(\frac{b}{x \Omega_2}\right)^t$, after some manipulations of (53) using [27, 3.351.4], we obtain¹³
¹⁴ the closed-form expression of the OP of D_2 as given in (13). The proof of Theorem¹⁴
¹⁵ 2 is completed.¹⁵

1 Appendix C

2 From (6) and (8), the expression of \mathcal{O}_{RD} is expressed as

$$\begin{aligned}
 3 \quad \mathcal{O}_{\text{RD}} &= \Pr(\max(\gamma_{\text{D}_1}^{x_1}, \gamma_{\text{D}_2}^{x_2}) < \gamma_{\text{th}}) \\
 4 \quad &= \Pr\left(\frac{a_1 P_{\text{R}} |g_1|^2}{N_0} < \gamma_{\text{th}}, \frac{(1-a_1) P_{\text{R}} |g_2|^2}{a_1 P_{\text{R}} |g_2|^2 + 1} < \gamma_{\text{th}}\right) \\
 5 \quad &= \Pr\left(XY < \frac{\gamma_{\text{th}}}{a_1 \phi P_{\text{S}}}, XZ < \frac{\gamma_{\text{th}}}{\phi P_{\text{S}}(1-a_1(1+\gamma_{\text{th}}))}\right). \quad (54)
 \end{aligned}$$

8 Based on the definition of the conditional probability, we have

$$\begin{aligned}
 9 \quad \mathcal{O}_{\text{RD}} &= \int_0^{\infty} \Pr\left(Y < \frac{\mathcal{A}}{x}, Z < \frac{\mathcal{B}}{x}\right) f_X(x) dx \\
 10 \quad &= \int_0^{\infty} \int_0^{\frac{\mathcal{B}}{x}} \Pr\left(Y < \frac{\mathcal{A}}{x}\right) f_Z(z) f_X(x) dx dz \\
 11 \quad &= \int_0^{\infty} \left[1 - e^{-\frac{\mathcal{A}}{x}} - e^{-\frac{\mathcal{B}}{x}} + e^{-\frac{\mathcal{A}}{x} - \frac{\mathcal{B}}{x}}\right] f_X(x) dx, \quad (55)
 \end{aligned}$$

12 where $\mathcal{A} = \frac{\gamma_{\text{th}}}{a_1 \phi P_{\text{S}}}$, $\mathcal{B} = \frac{\gamma_{\text{th}}}{\phi P_{\text{S}}(1-a_1(1+\gamma_{\text{th}}))}$. After some manipulations, we get (31),
 13 completing the proof of Theorem 4.

18 Appendix D

19 This appendix aims to provide the CDF of the instantaneous SNR of the information
 20 symbol x_1 . The instantaneous end-to-end SNR of symbol x_1 is $X_1 = \min(\gamma^{\text{R}}, \gamma_{x_1}^{\text{D}_1})$.

21 Thus, the CDF of X_1 is given by

$$\begin{aligned}
 22 \quad F_{X_1}(\xi_1) &= \Pr(\min(\gamma^{\text{R}}, \gamma_{x_1}^{\text{D}_1}) \leq \xi_1) \\
 23 \quad &= 1 - \Pr\left(P_{\text{S}} |h_1|^2 > \xi_1, a_1 \phi P_{\text{S}} |h_1|^2 |g_1|^2 > \xi_1\right) \\
 24 \quad &= 1 - \frac{1}{\Omega_1} \int_{\frac{\xi_1}{P_{\text{S}}}}^{\infty} \exp\left(-\frac{\xi_1}{\Omega_2 a_1 \phi P_{\text{S}} x} - \frac{x}{\Omega_1}\right) dx \\
 25 \quad &\approx 1 - \sqrt{\frac{4\xi_1}{\Omega_1 \Omega_2 a_1 \phi P_{\text{S}}}} K_1 \left(\sqrt{\frac{4\xi_1}{\Omega_1 \Omega_2 a_1 \phi P_{\text{S}}}}\right). \quad (56)
 \end{aligned}$$

31 Acknowledgements

32 Not applicable.

33 Funding

There is no funding support for this paper.

1	Abbreviations	1
2	5G Fifth Generation	2
3	AWGN Additive White Gaussian Noise	3
4	CDF Cumulative Distribution Function	4
5	CR-NOMA Cognitive Radio Non-Orthogonal Multiple Access	5
6	CSI Channel State Information	6
7	DF Decode-and-Forward	7
8	EH Energy Harvesting	8
9	F-NOMA Fixed-power-allocation Non-Orthogonal Multiple Access	9
10	IoT Internet of Things	10
11	NOMA Non-Orthogonal Multiple Access	11
12	OMA Orthogonal Multiple Access	12
13	OOP Overall Outage Probability	13
14	OP Outage Probability	14
15	PDF Probability Density Function	15
16	PS Power Splitting	16
17	RF Radio Frequency	17
18	RV Random Variable	18
19	SINR Signal-to-Interference-and-Noise Ratio	19
20	SNR Signal-to-Noise Ratio	20
21	SWIPT Simultaneous Wireless Information and Power Transfer	21
22	TDD Time Division Duplex	22
23	TS Time Switching	23
24	Availability of data and materials	24
25	All necessary data and materials have been presented in this paper. There are no extra data and materials to share.	25
26	Ethics approval and consent to participate	26
27	Not applicable.	27
28	Competing interests	28
29	The authors declare that they have no competing interests.	29
30	Consent for publication	30
31	Not applicable.	31
32	Authors' contributions	32
33	N.T.T. created the main ideas and formulates the problem; T.M.H. executed performance evaluation by rigorous	33
34	analysis and wrote the original draft; B.C.N. wrote simulation programs; P.T.T. worked as an advisor of the project	34
35	to discuss and advise the main ideas, provided recommendations on performance evaluation, and edited the final	35
36	manuscript.	36
37	Authors' information	37
38	N. N. Thang received the B.S. degree in Le Quy Don Technical University, in 1991, the M.S. degree and the Ph.D.	38
39	degree in Le Quy Don Technical University, Hanoi, Vietnam, in 1997 and 2017, respectively. His major interests are	39
40	antenna, wireless communication and cooperative communication. E-mail: nguyennhuthang@tcu.edu.vn.	40
41	T. M. Hoang was born in November 07, 1977. He received the B.S. degree in Communication Command at	41
42	Telecommunications University, Ministry of Defense, Vietnam, in 2002, and the B.Eng. degree in Electrical	42
43	Engineering from Le Quy Don Technical University, Ha Noi, Vietnam, in 2006. He obtained the M.Eng. degree in	43
44	Electronics Engineering from Posts and Telecommunications, Institute of Technology, (VNPT), Vietnam, in 2013.	44
45	He is currently pursuing the Ph.D degree at Le Quy Don Technical University, Hanoi, Vietnam. His research	45
46	interests include energy harvesting, Non-Orthogonal Multiple Access, and signal processing for wireless cooperative	46
47	communications. E-mail: tranmanhhoang@tcu.edu.vn.	47
48	B. C. Nguyen was born in Nghean, Vietnam. He received the B.S. in 2006 in Telecommunication University and	48
49	M.S. in 2011 in Posts and Telecommunications, Institute of Technology, (VNPT), Vietnam. He is currently pursuing	49

¹the Ph.D degree at Le Quy Don Technical University, Hanoi, Vietnam. His research interests include energy 1
²harvesting, full-duplex, and cooperative communication. E-mail: bacao.sqtt@gmail.com. 2
P. T. Tran was born at Ho Chi Minh City, Vietnam, in 1979. He received B.Eng. and M.Eng degrees in electrical 3
³engineering from Ho Chi Minh University of Technology, Ho Chi Minh City, Vietnam in 2002 and 2005, respectively. 3
⁴In 2007, he became a Vietnam Education Foundation Fellow at Purdue University, U.S.A., where he received his 4
⁴M.S. degree in mathematics and Ph.D. degree in electrical and computer engineering in 2013. In 2013, he joined the 4
⁵Faculty of Electrical and Electronics Engineering of Ton Duc Thang University, Vietnam and served as the Vice 5
⁵Dean of Faculty since October 2014. He is currently an IEEE Senior Member. His major interests are in the area of 5
⁶wireless communications and network information theory. Email: tranhanhphuung@tdtu.edu.vn. 6

Author details 7

⁸¹Nha Trang Telecommunications University, Nha Trang, Vietnam. ²Wireless Communications Research Group, 8
⁹Faculty of Electrical and Electronics Engineering, Ton Duc Thang University., Ho Chi Minh City, Vietnam. 9

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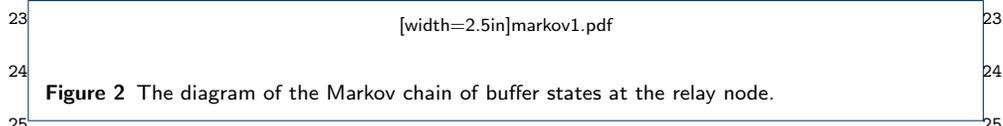
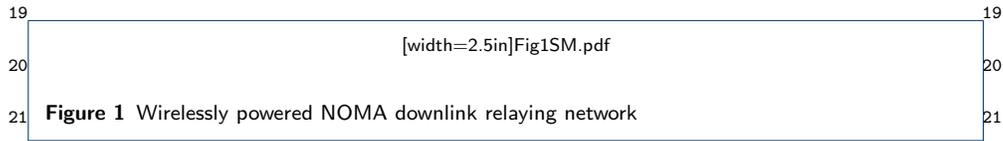
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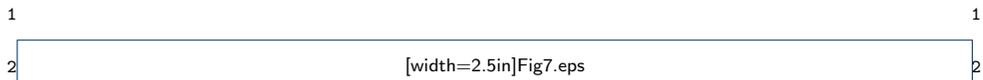


Figure 3 The average packet delay versus SNR according to theoretical analysis.

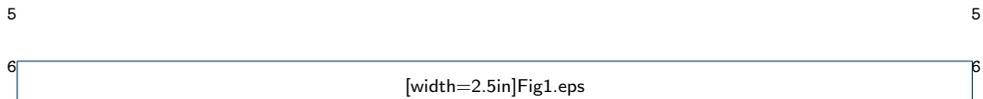


Figure 4 Overall outage probability versus average SNRs for optimal and fixed power allocation.

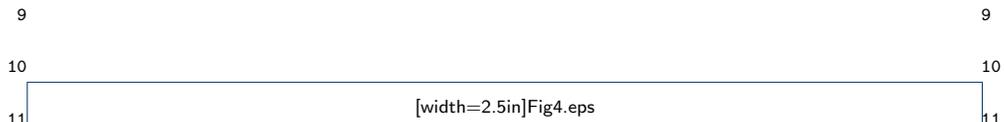


Figure 5 Outage probability versus the SNR with optimal power allocation for the cases of with buffer and without buffer aided relaying.

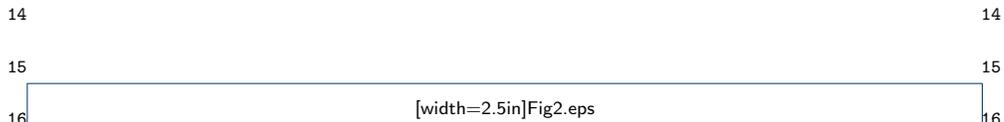


Figure 6 Outage probability of D_1 versus the transmission power of the source for optimal and fixed power allocation.

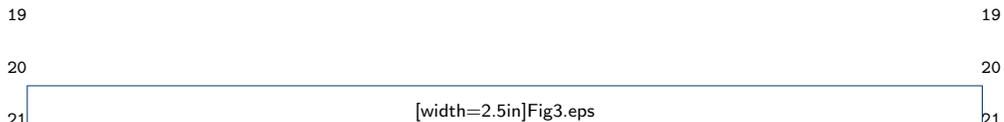


Figure 7 Outage probability of D_2 versus its SNR for the cases of optimal power allocation and fixed power allocation.

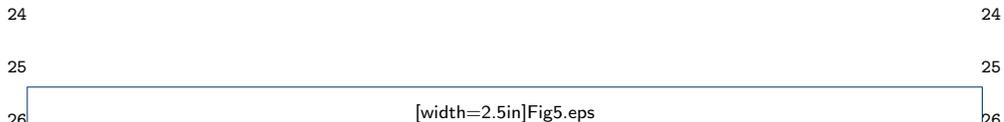


Figure 8 The effect of power allocation coefficient on the OP for different data rates, $E_b/N_0 = 10\text{dB}$.

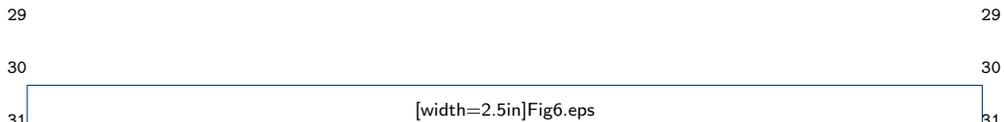


Figure 9 Average capacity of the system versus its SNR.

Figures

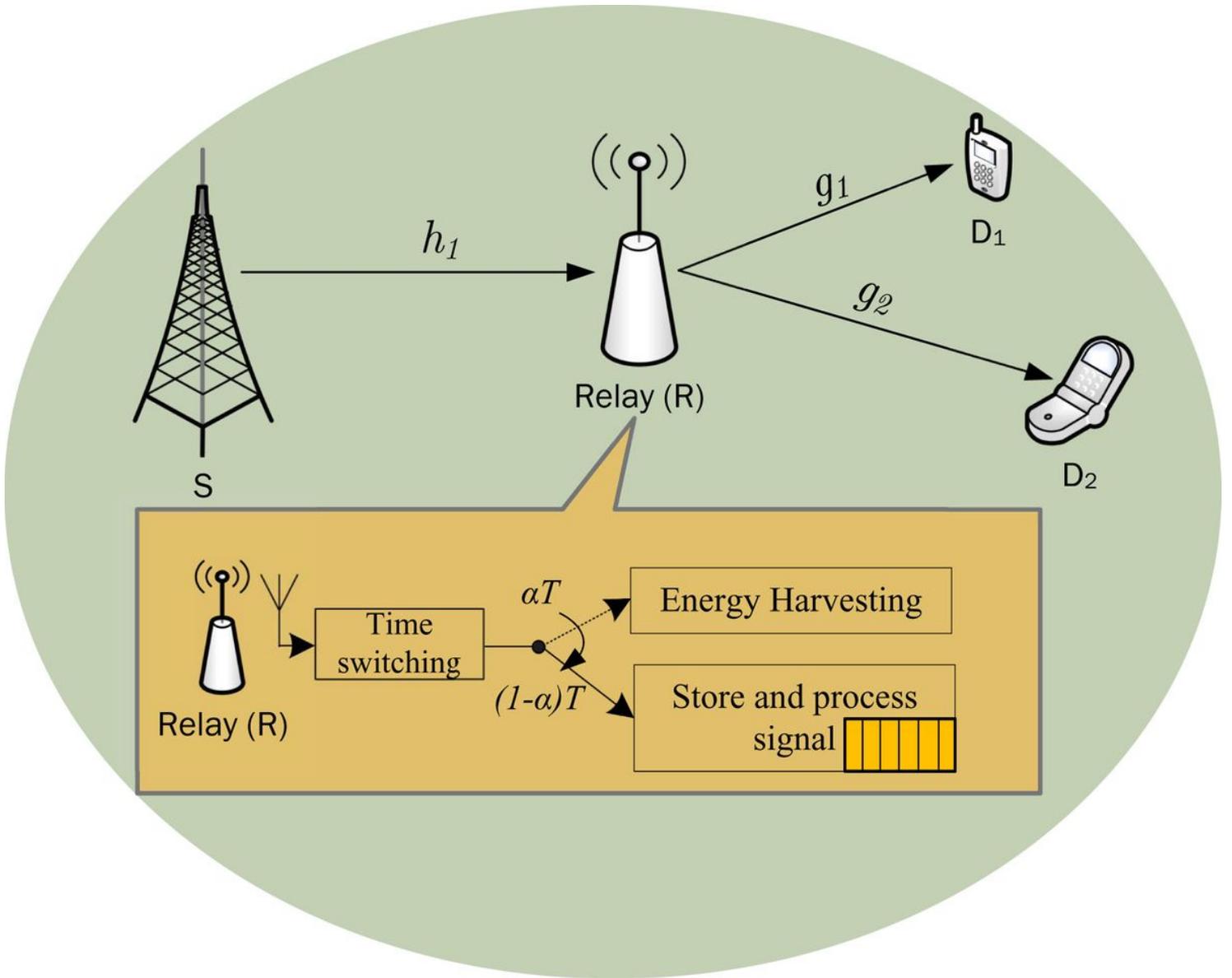


Figure 1

Wirelessly powered NOMA downlink relaying network

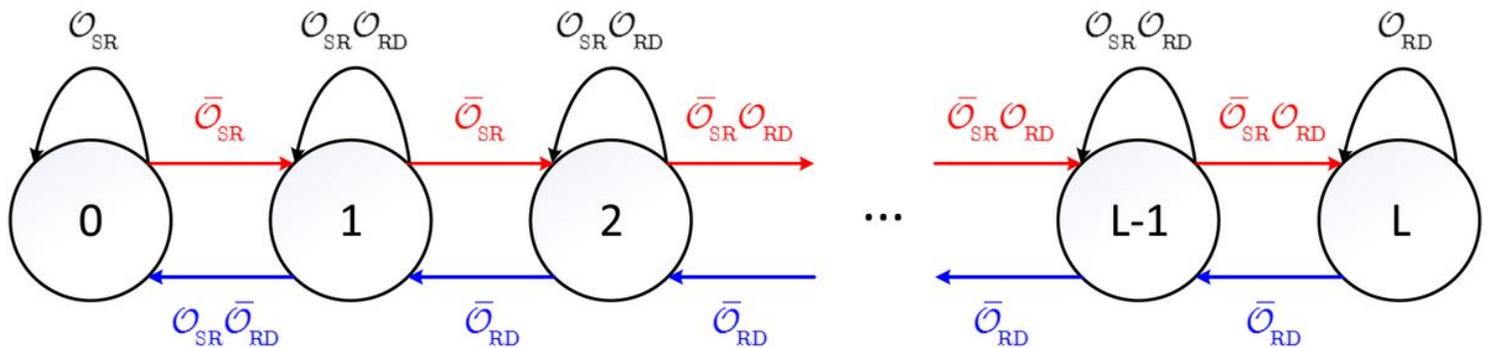


Figure 2

The diagram of the Markov chain of buffer states at the relay node.

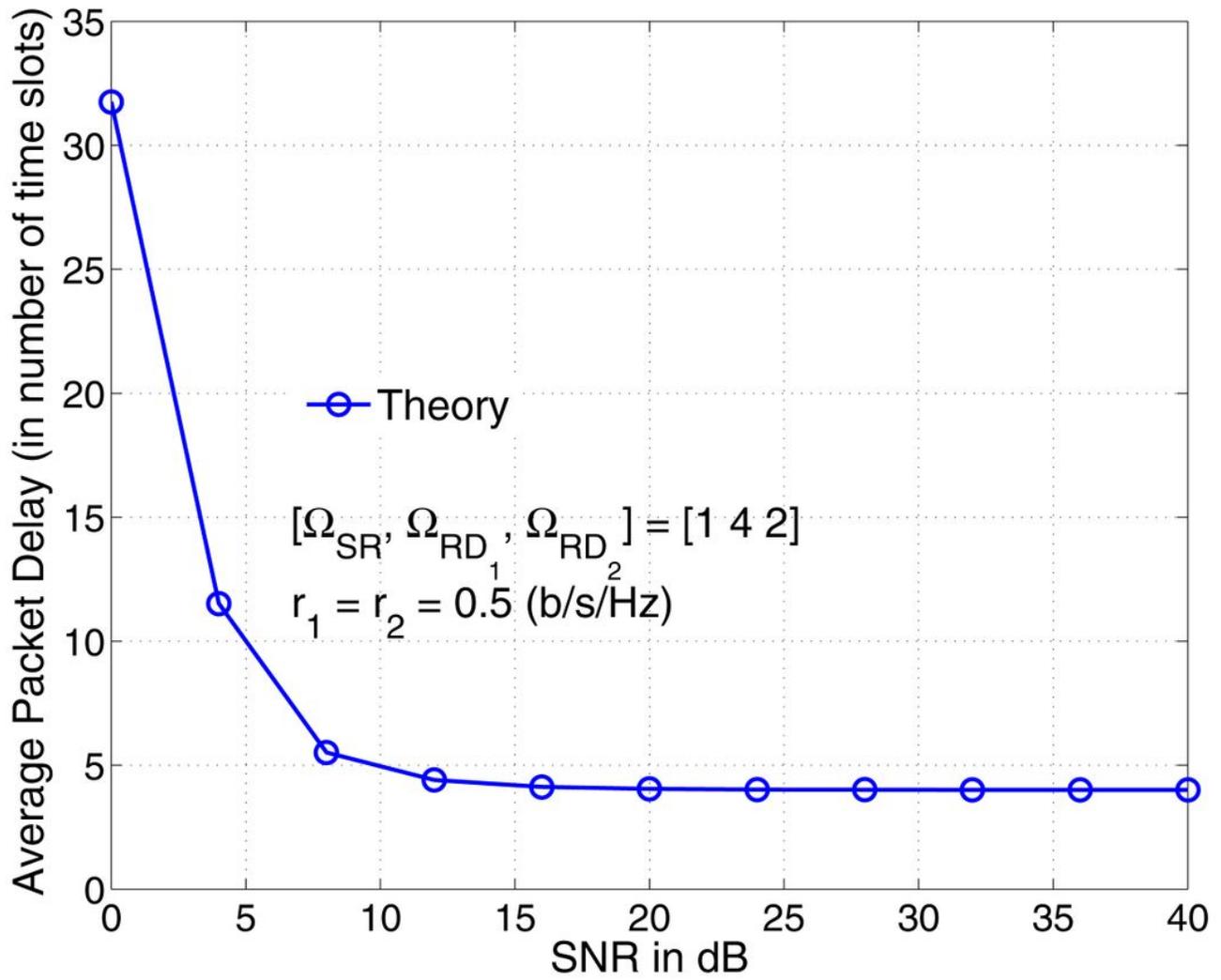


Figure 3

The average packet delay versus SNR according to theoretical analysis.

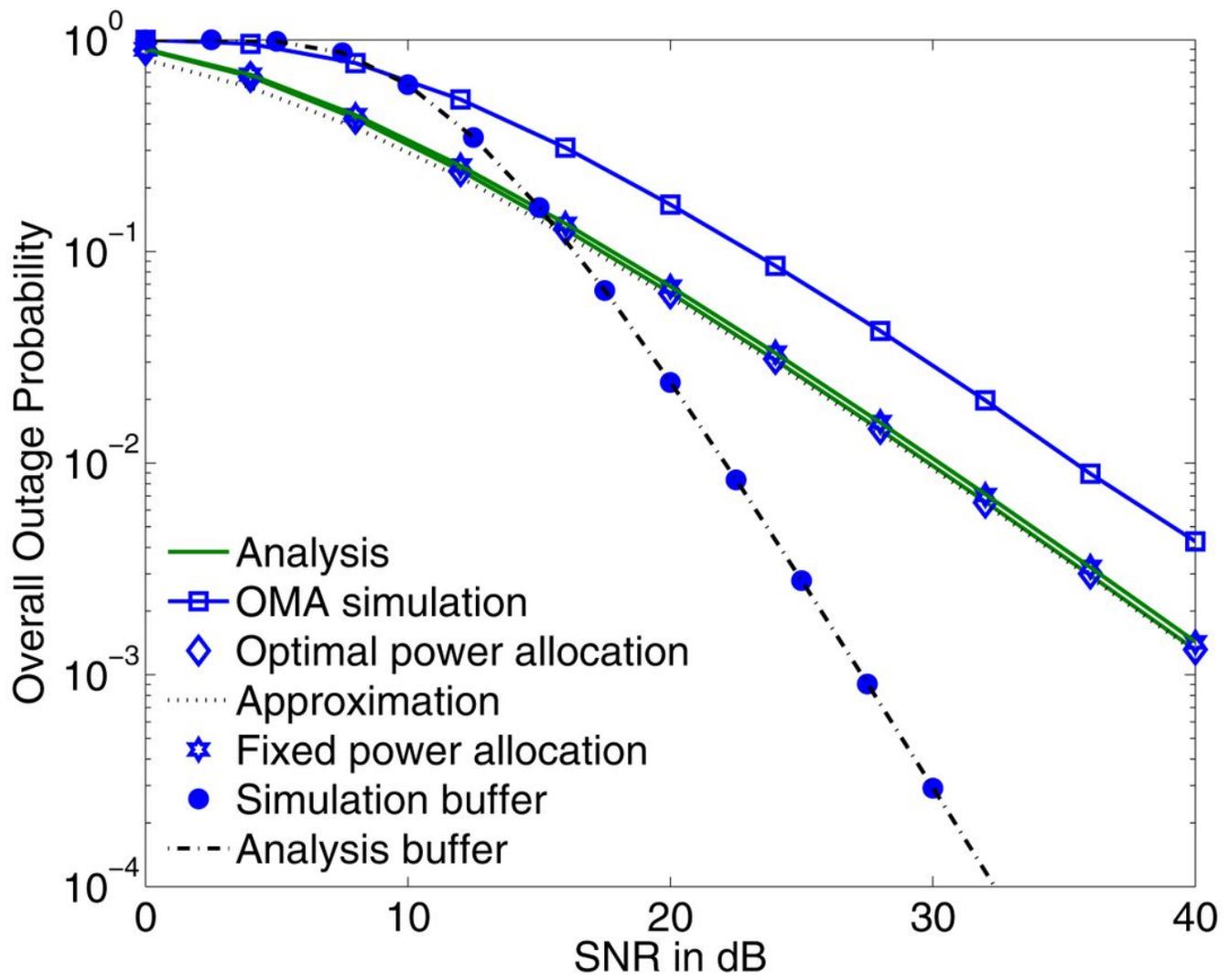


Figure 4

Overall outage probability versus average SNRs for optimal and fixed power allocation.

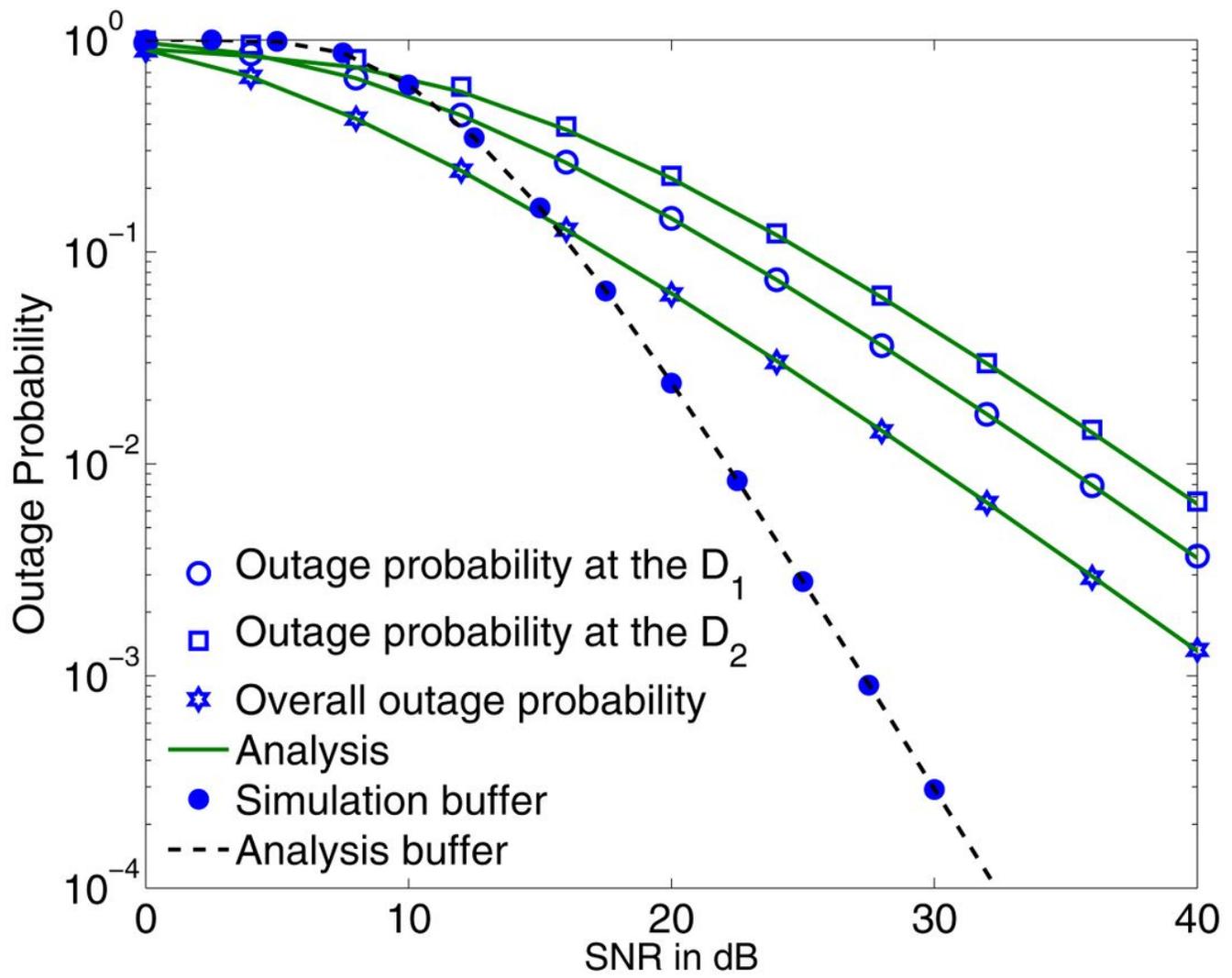


Figure 5

Outage probability versus the SNR with optimal power allocation for the cases of with buffer and without buffer aided relaying.

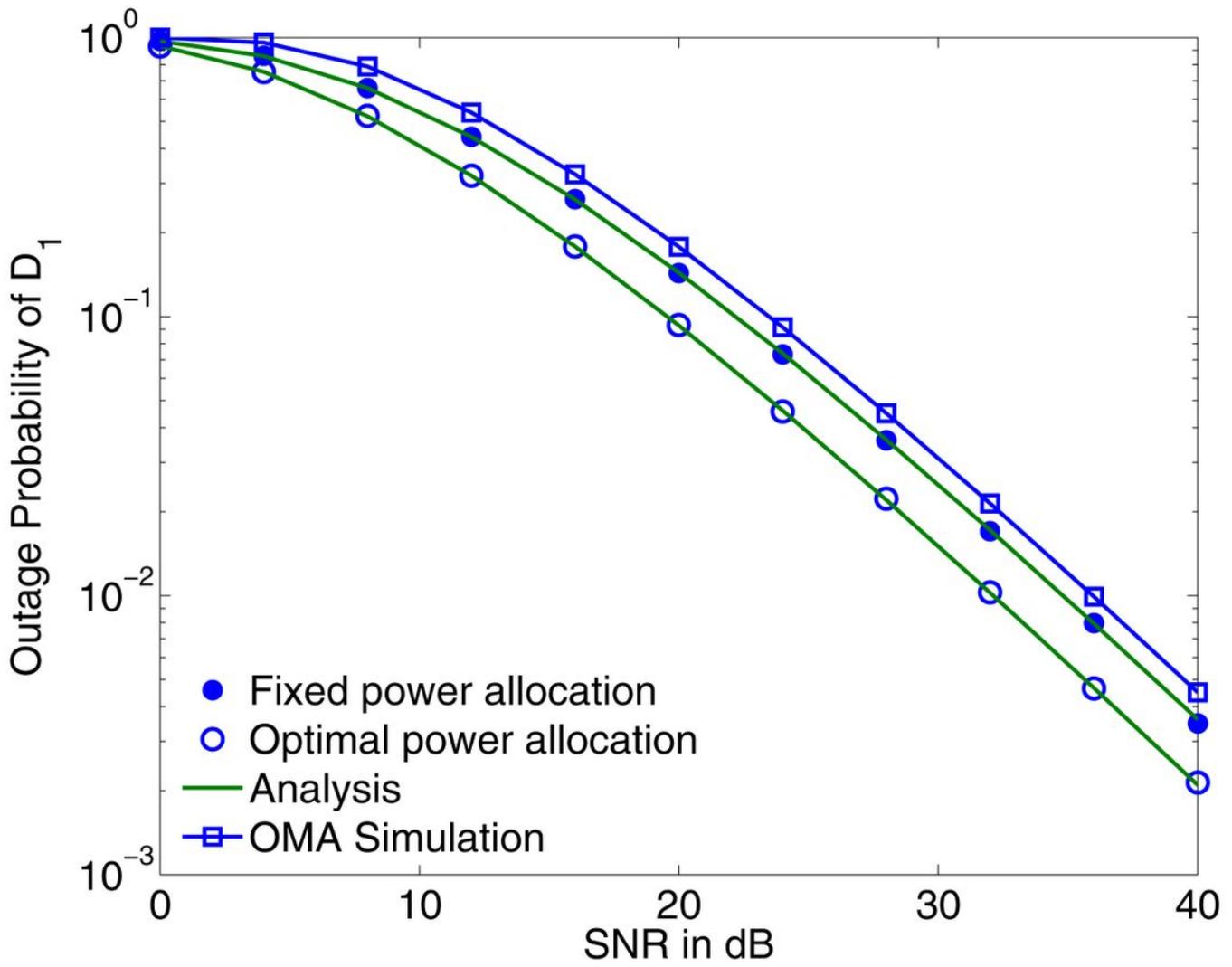


Figure 6

Outage probability of D_1 versus the transmission power of the source for optimal and fixed power allocation.

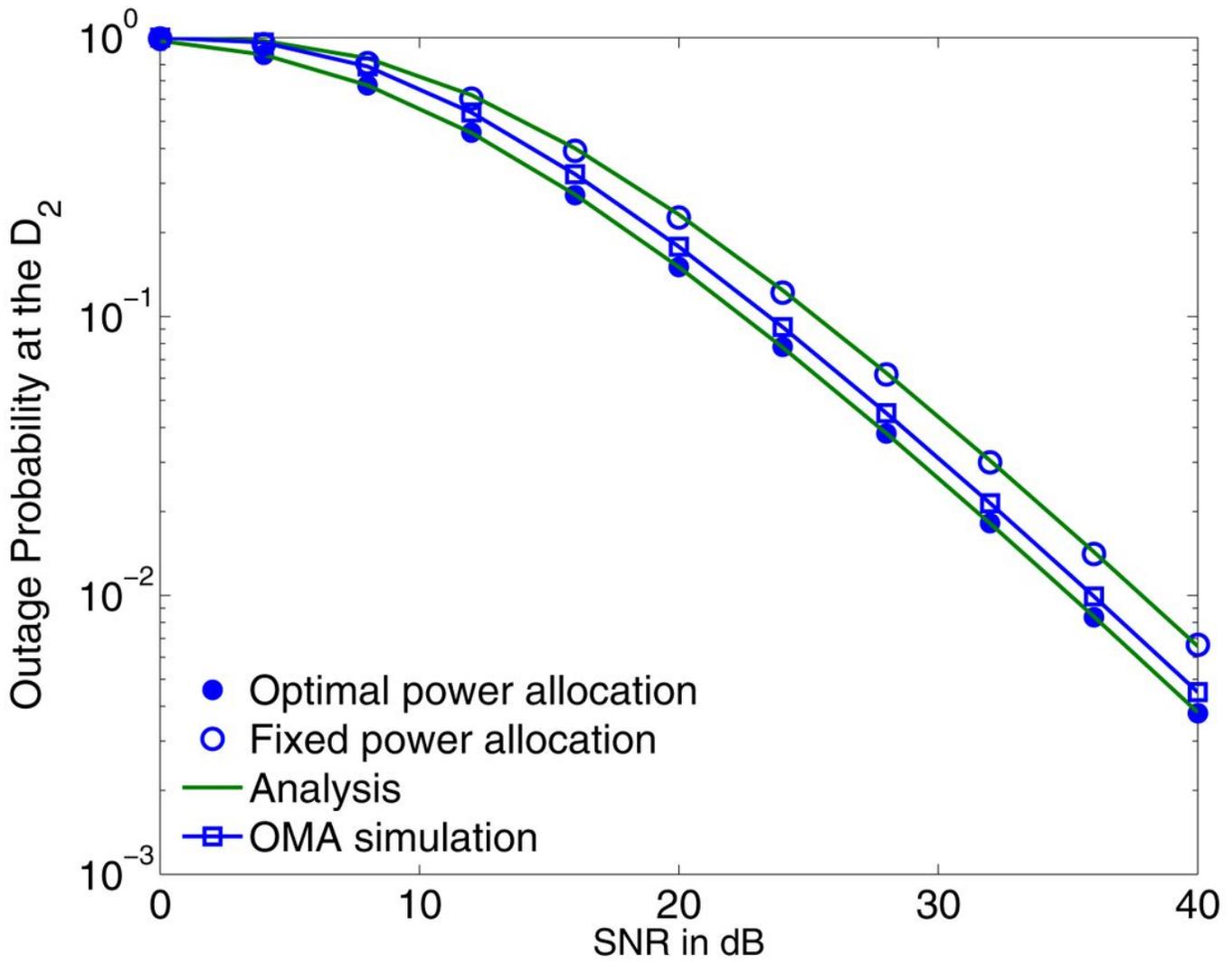


Figure 7

Outage probability of D2 versus its SNR for the cases of optimal power allocation and fixed power allocation.

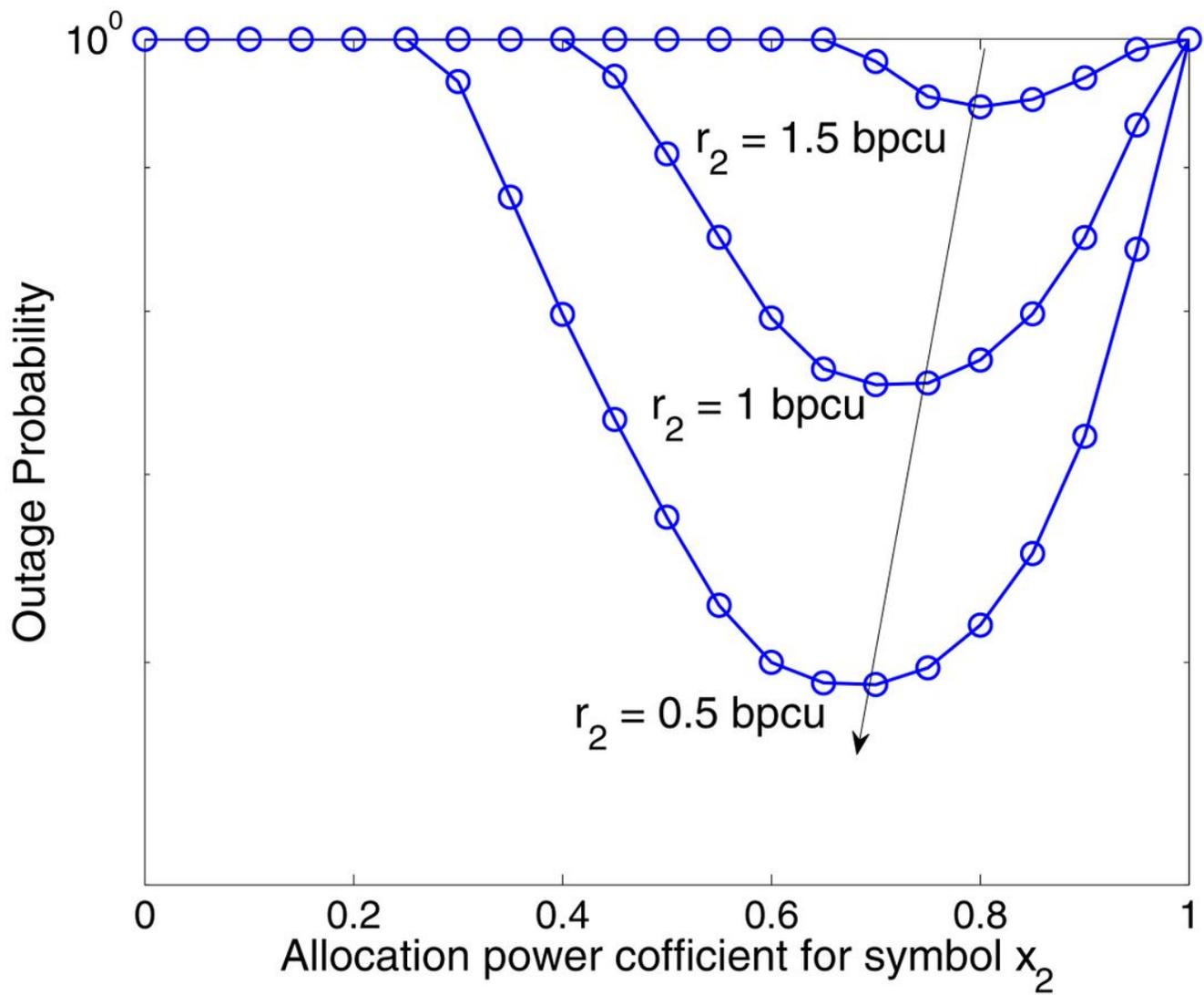


Figure 8

The effect of power allocation coefficient on the OP for different data rates, $E_b/N_0 = 10$ dB.

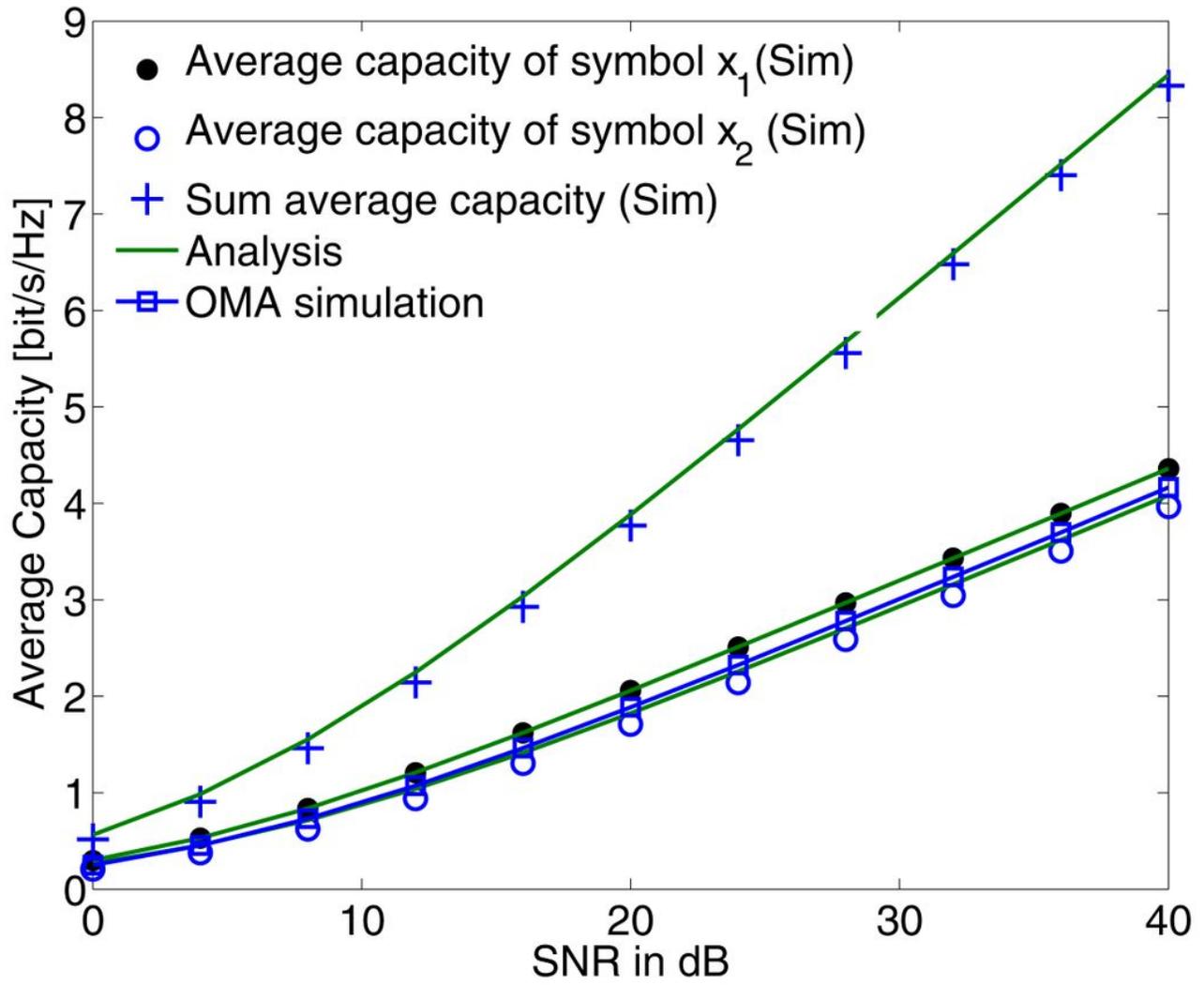


Figure 9

Average capacity of the system versus its SNR.