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A ranking-based differential evolution algorithm for hybrid flow shop sustainable scheduling

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Abstract

With the increasing of environmental pressure, the sustainable production for hybrid flow shops (HFS) has attracted more attention due to its broad industrial applications. In implementing the sustainable production of HFS, the selection of parallel machines for various jobs is a vital step. In light of this, a multi-objective mathematical model for minimization of makespan and energy consumption of HFSP was formulated. The sustainability of parallel machines were evaluated and ranked according to fuzzy TOPSIS method. To solve the multi-objective model of HFSP, an improved differential evolution algorithm was presented to assign the job with the ranked parallel machines, which can narrow the search scope and accelerate the convergence speed. Finally, a case study was presented to evaluate the effectiveness of the proposed ranking-based algorithm and to prove the feasibility of the model. The results showed that the proposed improved algorithm outperforms NSGA-II and PSO in searching for non-dominated solutions, which can effectively solve the sustainable production of HFSP.

Keywords: hybrid flow shop scheduling, sustainable production; improved differential evolution algorithm

Introduction

Hybrid flow shop scheduling problem (HFSP) indicated a combination of the classic flow-shop scheduling problem and the parallel machine scheduling problem, which was characterized by adjusting of jobs processing sequence and allocating the parallel equipment reasonably (Li et al. 2018; Tian et al. 2018; Kong et al. 2020; Xu et al. 2013). With the pressure of environmental problems, sustainable production for HFSP have aroused extensive concerns nowadays. Sustainable production for HFSP can be evaluated with three criterions, i.e., environmental impact, technical performance and

28 production cost. Most previous research for HFSP reported that the arrangement of parallel machines
29 for various jobs quite affect the sustainability in HFSP.

30 According to the characteristic of parallel machines, the classic HFSP can be divided into three
31 types, i.e., the HFSP with identical parallel machines (HFSP-IPM), the HFSP with uniform machines
32 (HFSP-UM), and the HFSP with unrelated parallel machines (HFSP-UPM). For HFSP-IPM, a job at a
33 certain stage have the same speed, which referring identical processing time on every machine; For
34 HFSP-UM, the speed of parallel machines is various for each job at a stage, and the processing time is
35 inversely proportional to the speed. HFSP-UPM is the most complex and closed to real industrial
36 manufacturing scheduling problems, which has a wide engineering application (Meng et al. 2019). For
37 HFSP-UPM, the processing time of a job on any parallel machines are independent and the production
38 efficiency quite depends on the matching degree between the job and the machines (Kong et al. 2020).
39 Hence, to implement sustainable production, reasonable evaluation of sustainability for the parallel
40 machines become an urgent demand for scheduling.

41 The sustainability of parallel machines can be evaluated from three aspects, i.e., the environmental
42 impact, the technical performance and the production cost. Energy consumption of the parallel
43 machines was supposed to the main source of environmental impact, which has been under extensive
44 research. Drake et al. (Drake et al. 2006) showed that there were significant amounts of energy
45 associated with machine start-up and idling. In light of this, Mouzon (Mouzon et al. 2007) proposed a
46 greedy randomized adaptive search procedure integrated with GA to solve the minimization of total
47 energy consumption of HFSP. The result showed that a significant amount of energy can be saved when
48 non-bottleneck are turned on during a long idle time. Liu (Liu et al. 2019) presented an ultra-low idle
49 state of machines by turning off some auxiliary parts in idle state. To calculate the energy consumption
50 of machine tools with respect to various cutting parameters, the specific energy consumption (SEC)
51 was generally recognized as an essential indicator. Besides, subsidiary materials, such as cutting fluid
52 are also have important influence on environmental impact. For production cost of parallel machines,
53 the machining cost, the human labor cost, etc., were generally investigated for optimizing the job
54 scheduling of the workshop (Behnamian et al. 2014). For technical performance, accuracy and
55 reliability in terms of the basic requirements of production have been focused (Dimitrov et al. 2014). In
56 addition, axes configurations, in particular at a high speed spindle rotations and high feedrates, has a
57 significant variation of the dynamic properties for the machines (Brecher et al. 2016). To sum up, a

58 systematic sustainable model of machines with respect to the above factors has to be taken into
59 account.

60 For the optimization of HFSP, the genetic algorithm (Han et al. 2018; Xu et al. 2017), grey wolf
61 algorithm (Ni et al. 2020), backtracking search algorithm (Lu et al. 2019), ant colony optimization
62 (Kang et al. 2014; Feng et al. 2019), particle swarm optimization (Fang et al. 2020), differential
63 evolution algorithm (Tian et al. 2016; Miao et al. 2015) were developed. It is noted that the selection of
64 machines for various jobs in HFSP was random for those algorithms in the iterations, which result in a
65 lower efficiency. To overcome this problem, local search method with a forward decoding method
66 considering idle time (Wang et al. 2012; Tian et al. 2019), energy-saving capability (Ding et al. 2016),
67 the total weighted tardiness (Ding et al. 2016), etc., were proposed. Nevertheless, the methods mostly
68 neglected the sustainability of unrelated parallel machines, which is not applicable to sustainable
69 production.

70 To this end, this work presented a sustainable model of HFSP for minimization of makespan and
71 energy consumption. In addition, an improved differential evolution algorithm integrated with the
72 ranking sustainability of parallel machines was developed to support the HFSP sustainable scheduling.
73 The innovations of the approach are summarized below:

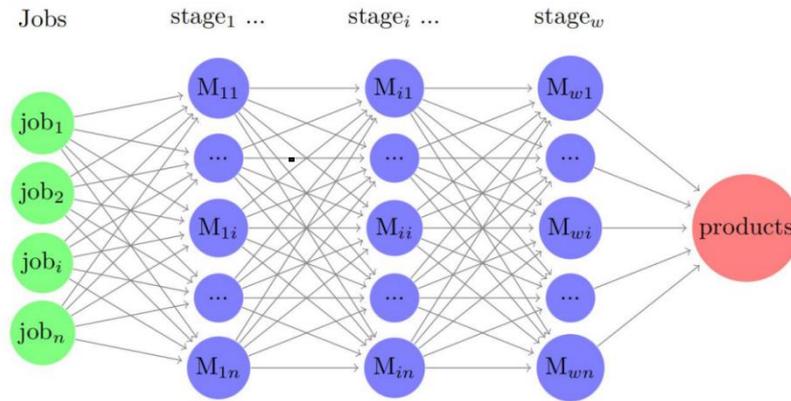
- 74 ● A multi-objective mathematical model for minimization of makespan and energy consumption of
75 HFSP was formulated.
- 76 ● The sustainability of parallel machines were evaluated and ranked according to fuzzy TOPSIS
77 method.
- 78 ● To solve the multi-objective model of HFSP, an improved differential evolution algorithm was
79 presented to assign the job with the ranked parallel machines, which can narrow the search scope
80 and accelerate the convergence speed.

81 The rest of this paper is organized as: Section II describes a multi-objective sustainable model of
82 HFSP. Section III elaborates a ranking-based differential evolution algorithm considering the
83 sustainability of machines. Section IV presents the solutions to several cases. Finally, a conclusion and
84 some future research issues was drawn in section V.

85 **Problem formulation**

86 The hybrid flow shop scheduling problem (Fig.1) is commonly described as follows: there are a

87 set of n jobs has to be processed at w stages in series. Each job i consists of a pre-determined sequence
 88 of operations. Each operation requires one machine selected from a set of available parallel machines,
 89 which is denoted as M_{ij} . The formal mathematical definition of the problem will be described in detail
 90 in the following sections.



91
 92 Fig.1 Illustration of HFSP

93 Hybrid flow shop scheduling has been extensively examined and the main objective has been to
 94 improve production efficiency. However, limited attention has been paid to the consideration of energy
 95 consumption with the advent of green manufacturing. In order to reduce resource and energy
 96 consumption and achieve sustainable production, this paper set the assignment of machines and the
 97 sequence of operations on all the machines as variables in HFSP to minimize the makespan T and
 98 energy consumption E to realize sustainable production.

99 Hypotheses considered in this paper are summarized as follows:

- 100 (1) Jobs are independent, and have equal priority.
 101 (2) After a job is processed on a machine, it is transported to the next machine immediately.
 102 (3) All jobs and machines are available at time zero, turning off the idle machines is not allowed and
 103 machine failure is not considered.
 104 (4) The machine cannot be turned off completely until it has finished all operations assigned to it.
 105 (5) The order of operations for each job is predefined and cannot be modified.
 106 (6) Pre-emption is not allowed, that is, no task can be interrupted before the completion of its current
 107 operation.

108 The notations used throughout the study are listed in Table 1.

109 Table 1 notations of all the parameters for HFSP

$S_{t_{ijk}}$	start time for the i th job of the j th process at machine k
$F_{t_{ijk}}$	completion time for the i th job of the j th process at machine k

i	the job number
j	the operation number
k	the machine number
n	the number of total jobs
w	the number of total stages
m	the number of total machines
k^*	the selected machine number for next stage
X_{ijk}	it is equal to 1 if the machine is selected; otherwise, it is equal to 0
M	machine set
obj	objective function
T	makespan
E	total energy consumption
Pt_{ijk}	processing time for the i th job of the j th process at machine k
Tt_{ijk}	transportation time for the i th job of the j th stage at machine k to $j+1$ th stage
TE_{ijk}	unit transportation energy consumption for the i th job of the j th stage at machine k to $j+1$ th stage
d_{jk}	The distance from the j stage at machine k to the machine tool in the next stage
V	transportation speed
PE_{ijk}	unit processing energy consumption for the i th job of the j th process at machine k
WE_{ijk}	waiting energy consumption per unit time for the i th job of the j th process at machine k
Wt_{ijk}	waiting time for the i th job of the j th process at machine k
PE_{total}	total processing energy consumption
WE_{total}	total waiting energy consumption
TE_{total}	total transportation energy consumption

110 In this work, the process of production can be divided into three stages: (1) processing stage; (2)
111 waiting stage;(3) transportation stage. Each stage corresponding to the consumption of time and energy.
112 The processing time Pt_{ijk} , processing energy consumption per unit time PE_{ijk} , waiting energy
113 consumption per unit time WE_{ijk} of each operation on m machine, transportation distance of machine
114 tools in different stages d_{jk} and transportation energy consumption per unit time TE_{ijk} are deterministic
115 and known in advance. The waiting time Wt_{ijk} was shown in Eq. (1), which was associated with the
116 schedule scenario. It was determined by the completion time Ft_{ijk} of the work piece and the start time of
117 the next work piece $St_{ij(k+1)}$ on the same machine. Clearly, the magnitude of transportation time Tt_{ijk} is
118 determined by the distance between two consecutive machines, which was formulated in Eq. (2). The
119 transport speed V is assumed to be fixed for the convenience of calculation. The makespan means the
120 maximum completed time of all the jobs. It can be described in detail as shown in Eq. (3).

$$Wt_{ijk} = St_{ij(k+1)} - Ft_{ijk} \quad (1)$$

$$Tt_{ijk} = d_{jk} / V \quad (2)$$

$$T = \max\{Ft_{ijk}\} \quad (3)$$

121 The total energy consumption was composed of the energy consumption for the processing
122 stage PE_{total} , the waiting stage WE_{total} and the transportation stage TE_{total} . The processing energy is
123 determined by processing time and the processing power per unit time of machine tools. The waiting

124 energy of machine tools is the energy consumed by machine tools when they are not machining, that is,
 125 waiting to process next jobs. The transportation energy consumption TE_{total} is determined by the
 126 transportation process of the work-pieces between different stages, which can be calculated as Eq. (5).
 127 Therefore, the second objective of this problem is simplified below.

$$E = PE_{total} + WE_{total} \quad (4)$$

$$TE_{total} = \sum_{i=1}^n \sum_{j=1}^w \sum_{k=1}^m Tt_{ijk} \cdot TE_{ijk} \cdot X_{ijk} \quad (5)$$

$$PE_{total} = \sum_{k=1}^m \sum_{i=1}^n \sum_{j=1}^w Pt_{ijk} \cdot PE_{ijk} \cdot X_{ijk} \quad (6)$$

$$WE_{total} = \sum_{k=1}^m \sum_{i=1}^n \sum_{j=1}^w (St_{ij(k+1)} - Ft_{ijk}) \cdot Wt_{ijk} \cdot X_{ijk} \quad (7)$$

128 Mathematically, an integer linear programming model of the HFSP was formulated as the
 129 following, which will be used throughout the paper.

$$\min obj_1 = \max_{\forall i,j,k} \{Ft_{ijk}\} \quad (8)$$

$$\min obj_2 = \sum_{k=1}^m \sum_{i=1}^n \sum_{j=1}^w (Pt_{ijk} \cdot PE_{ijk} + (St_{ij(k+1)} - Ft_{ijk}) \cdot Wt_{ijk}) \cdot X_{ijk} \quad (9)$$

Subject to:

$$\sum_{k \in m} X_{ijk} = 1, \quad i \in \{1, 2, \dots, n\}; j \in \{1, 2, \dots, w\}; k \in \{1, 2, \dots, m\} \quad (10)$$

$$St_{ij} + \sum_{k \in m} Pt_{ijk} \cdot X_{ijk} \leq St_{i(j+1)}, \quad i \in \{1, 2, \dots, n\}; j \in \{1, 2, \dots, w\}; k \in \{1, 2, \dots, m\} \quad (11)$$

$$Ft_{i(j+1)k^*} - Ft_{ijk} \geq Pt_{i(j+1)k^*} \quad \forall i \in n, j \in w, k \in m \quad (12)$$

$$St_{i(j+1)k^*} - Ft_{ijk} \geq 0 \quad \forall i \in n, j \in w, k \in m \quad (13)$$

$$Ft_{ijk} - Ft_{(i-1)jk} \geq Pt_{ijk} \quad \forall i \in n, j \in w, k \in m \quad (14)$$

$$X_{ijk} = \begin{cases} 1, & \text{if machine } k \text{ is selected for operation} \\ 0, & \text{otherwise} \end{cases} \quad \forall i \in n, j \in w, k \in m \quad (15)$$

$$St_{ij(k+1)} - Ft_{ijk} \geq 0 \quad \forall i \in n, j \in w, k \in m \quad (16)$$

$$St_{ijk}, Ft_{ijk} \leq m \cdot X_{ijk} \quad \forall i \in n, j \in w, k \in m \quad (17)$$

$$0 \leq Ft_{ijk} \leq Tt_{ijk} \quad \forall i \in n, j \in w, k \in m \quad (18)$$

$$T_{make} \geq Ft_{ijk} \quad \forall i \in n, j \in w, k \in m \quad (19)$$

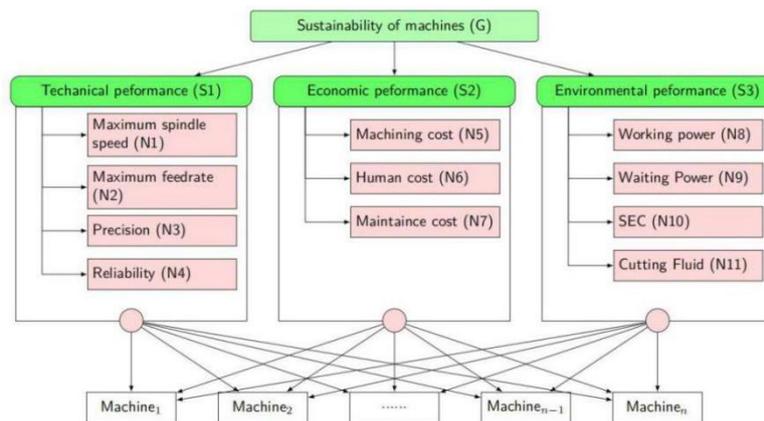
$$E_{cons} \leq E_{max}, T_{make} \leq T_{max} \quad (20)$$

130 where, Eq.(8) and Eq.(9) is the objective function. Eq.(10) ensures that each operation is assigned
 131 to only one machine from its candidate machine set. Eq.(11) means that the operations belonging to the

132 same job satisfy the precedence. constraints Eq.(12) and Eq.(13) ensures that work pieces are processed
 133 according to sequence constraints. Eq.(14) governs that at one time a machine can execute one
 134 operation and it becomes available for other operations only if the previous operation is completed.
 135 Eq.(15) defines the assignment of jobs and the sequence of machines. Eq.(16) represents the
 136 precedence relationship among various operations of a job. Eq.(17) ensures that the starting time of
 137 each operation in the unselected process plan is zero. Besides, the starting time of each operation
 138 should be non-negative. Eq.(18) specifies that the completion time of each operation transported from
 139 one machine to another machine is greater than the completion time of the corresponding operation,
 140 which should be greater than zero. Eq.(19) implies that the makespan is equal to or greater than the
 141 completion time of a schedule that includes the completion time of all jobs. Eq.(20) guarantees that the
 142 production processed meet the constrains of makespan and energy consumption.

143 Ranking for sustainability of parallel machines

144 In HFSP, the selection of machine tool from parallel machines is an important decision-making
 145 process. Generally, the selection processes is randomly assigned and then evaluated in the following
 146 iteration processes of HFSP, which is time-consuming. To overcome this problem, a heuristic method is
 147 proposed to rank and select the sustainability of parallel machines based on the fuzzy TOPSIS method
 148 (Tian et al. 2019) .



149
 150 Fig. 2 The Framework of sustainable evaluation for machine tools

151 Sustainable evaluation system of machines

152 To rank the sustainability of parallel machines, the hierarchical structure of this research decision
 153 problem is shown in Fig. 2. The sustainability of machines was evaluated by three aspects described as
 154 following.

- 155 ● Technical performance: the technical performance refers to the maximum machining quality and
 156 efficiency that the machines can achieve, which were indicated by the maximum spindle speed
 157 (N1), feedrate (N2), the accuracy (N3) and reliability (N4).
- 158 ● Economic performance: the economic performance of machines is indicated by the cost to
 159 processing the corresponding jobs. Generally, the cost can be calculated by the summary of the
 160 machining cost (N5), human cost (N6) and maintenance cost (N7).
- 161 ● Environmental performance: the most environmental impact was owing to the energy
 162 consumption, which can be characterized by the power of machines (working power, N8 and
 163 waiting power, N9) and specific energy consumption (SEC, N10). Besides, the cutting fluid
 164 (N11) is also an important factors for the environment impact.

165 Table 2 Membership function of linguistic scale

Linguistic evaluation of parallel machines	Linguistic evaluation of the weight of criteria	Scale of fuzzy number
Very low (VL)	Of little importance (VL)	(1,1,3)
Low (L)	Moderately important (MI)	(1,2,5)
Good (G)	Important (I)	(3,5,7)
High (H)	Very important (VI)	(5,7,8)
Excellent (Ex)	Absolutely important (AI)	(7,9,9)

166 Based on the construction of the hierarchy, the priority weights of each criteria and the attributes
 167 of parallel machines can be calculated with the fuzzy TOPSIS method. Here, the linguistic variables are
 168 defined and the corresponding membership function by triangular fuzzy number listed in Table 2. The
 169 fuzzy TOPSIS method of calculating priority weights and ranking the parallel machines is described
 170 shown in Fig.3.

Step1: data preparation

- **Isotropy decision matrix**
Construct the fuzzy performance matrix and choose the appropriate linguistic variables for the parallel machines with respect to criteria.

$$\tilde{A} = \begin{matrix} & \begin{matrix} N1 & N2 & \dots & N11 \end{matrix} \\ \begin{matrix} M_{i1} \\ M_{i2} \\ \vdots \\ M_{in} \end{matrix} & \begin{bmatrix} \tilde{a}_{11} & \tilde{a}_{12} & \dots & \tilde{a}_{1n} \\ \tilde{a}_{21} & \tilde{a}_{22} & \dots & \tilde{a}_{2n} \\ \vdots & \vdots & \ddots & \vdots \\ \tilde{a}_{m1} & \tilde{a}_{m2} & \dots & \tilde{a}_{mn} \end{bmatrix} \end{matrix}$$

For the elements in \tilde{A} , it can be obtained by the rating of experts.

- **Normalization of the fuzzy-decision matrix**
The normalized fuzzy-decision matrix denoted by \tilde{R}

$$\tilde{R} = \left[\tilde{r}_{ij} \right]_{m \times n}, i = 1, 2, \dots, m; j = 1, 2, \dots, n$$

$$\tilde{r}_{ij} = \left(\frac{l_{ij}}{u_j^+}, \frac{m_{ij}}{u_j^+}, \frac{u_{ij}}{u_j^+} \right), u_j^+ = \max_i \{u_{ij} | i = 1, 2, \dots, n\}$$

In this process, it is noted that the best level u_j^+ is set as one. The weighted fuzzy normalized decision matrix is shown as following matrix

$$\tilde{V} = \left[\tilde{v}_{ij} \right]_{m \times n}, i = 1, 2, \dots, m; j = 1, 2, \dots, n$$

where $\tilde{v}_{ij} = \tilde{r}_{ij} \otimes \tilde{w}_j$

Step2: data processing

- **Determination of the fuzzy positive-ideal solution (FPIS) and fuzzy negative-ideal solution (FNIS).**
we can define the FPIS A^+ (aspiration levels) and FNIS A^- (the worst levels) as following:

$$A^+ = \left(\tilde{v}_1^+, \dots, \tilde{v}_j^+, \dots, \tilde{v}_n^+ \right)$$

$$A^- = \left(\tilde{v}_1^-, \dots, \tilde{v}_j^-, \dots, \tilde{v}_n^- \right)$$

where $\tilde{v}_j^+ = (1, 1, 1) \otimes \tilde{w}_j = (lw_j, mw_j, uw_j)$
and $\tilde{v}_j^- = (0, 0, 0), j = 1, 2, \dots, n.$

- **Calculate the distance of each parallel machine from FPIS and FNIS**

The distances (d_i^+ and d_i^-) of each parallel machines from A^+ and A^- can be currently calculated by the following equation.

$$d_i^+ = \sum_{j=1}^n d \left(\tilde{v}_{ij}, \tilde{v}_j^+ \right), i = 1, 2, \dots, m; j = 1, 2, \dots, n$$

$$d_i^- = \sum_{j=1}^n d \left(\tilde{v}_{ij}, \tilde{v}_j^- \right), i = 1, 2, \dots, m; j = 1, 2, \dots, n$$

Step3: Machines ranking

- **Determine the fuzzy gap degree**
- **Machines sustainability ranking based on the fuzzy gap degree.**

$$CC_i = \frac{d_i^+}{d_i^+ + d_i^-}, i = 1, 2, \dots, m$$

{M2 M4 M3 ... M1}—Ascending order

171

172

Fig. 3 Procedure of ranking the parallel machines based on fuzzy TOPSIS method

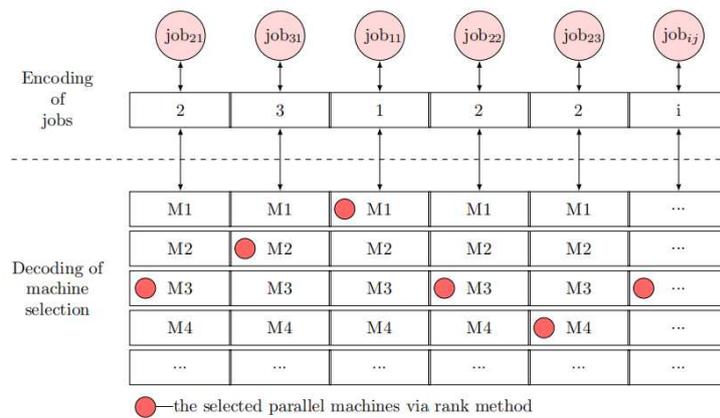
173

174 An improved differential evolution algorithm

175 Differential Evolution (DE) is a stochastic direct search and global optimization algorithm include
 176 genetic operation, i.e., differential mutation, crossover and selection. It is a parallel search evolution
 177 strategy that is fairly fast and reasonably robust, which make differential evolution the versatile tool
 178 today. Inspired by No Free Lunch Theorem, the ranking of machines' sustainability introduced in above
 179 section was integrated with the DE by the heuristic decoding rule. In this work, a ranking-based
 180 differential evolution algorithm (RBDE) was proposed to solve the hybrid flow-shop problems based
 181 on the heuristic decoding rule. In this approach, after the genetic operation, active decoding method
 182 was carried out and integrated into iterations of the algorithm. The selection pressure to the proper
 183 direction could be obtained so as to enhance its performance by accelerating the convergence speed in
 184 the iteration of DE. The procedures are described in detail as below.

185 **Ranking-based Decoding**

186 Decoding method is the key factor to decode the iterative chromosome sequence into a reasonable
 187 production scheduling scheme, which has a great impact on the efficiency of the solution. The heuristic
 188 rule proposed in this paper mainly refers to the directional selection of machine tools according to the
 189 sustainability of the parallel machines. The parallel machines with the highest sustainability are more
 190 suitable for the jobs. In this work, the jobs is encoded by the number and the sequence as shown in
 191 Fig.4. In the decoding processes, the jobs would be assigned to the machines with the highest
 192 sustainability via the above ranking method.



193
194 **Fig.4 heuristic decoding rule**

195
196 **Calculation procedure**

197 Based on the above decoding method, the ranking-based integration with the heuristic decoding
 198 rule was developed for HFSP sustainable scheduling. The flowchart of the RBDE was illustrated in
 199 Fig.5. It can realize directional selection of the most sustainbale machines to accelerate the
 200 convergence speed. The scheduling process was presented as follows:

201 **Initializing:** The population P_0 were generated with uniform distribution in search space, and the
 202 control parameters F and Cr was provided based on the experience. .

203 **Mutation:** NP mutants were generated with the following mutation strategy in Eq.(21).

204
$$Mu(NP) = x_{r1} | P_0 + F \cdot (x_{r2} - x_{r3}) | P_0 \quad (21)$$

205 **Crossover:** Population P_0 and its mutant intermediates NP was implemented as Eq. (22).

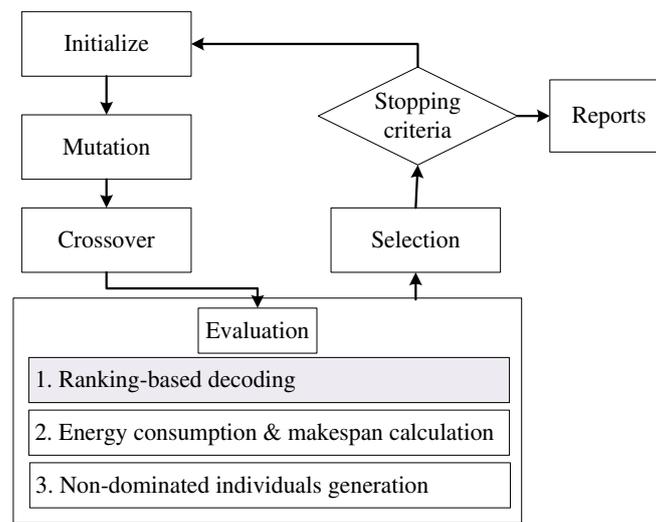
206
$$C_r(P_0 + 1) = \begin{cases} x_{gr}(NP), & \text{if } rand(0,1) \leq cp \\ x_{gr}(P_0), & \text{otherwise} \end{cases} \quad (22)$$

207 where, $rand(0,1)$ is the random numbers which are uniform distribution between 0 and 1. $x_{gr}()$
 208 represents the selection of gene of the r th chromosome for the population.

209 **Evacuation & Selection:** Based on the ranking-based decoding method, the parallel machines were
 210 selected for various jobs. Then, the energy consumption and makespan could be calculated with the
 211 developed HFSP model. In this step, only non-dominated individuals were carried on to the next
 212 generations.

213 **Stopping criteria:** The iteration repeated until it achieved the maximum number of generations.

214 **Report:** The best individuals' outcome as the optimal solution.



215

216 Fig.5 Flowchart of the RBDE

217

218 Table 3 The parameters in HFSP

processes	machines	job1	job2	job3	job4	processing power(kW)	waiting power(kW)
		P_T (min)	P_T (min)	P_T (min)	P_T (min)		
process 1	M1	6	4.2	6.1	5.8	20	15
	M2	2.3	2.6	4.1	3.4	30	25
	M3	4.2	3.2	5.6	4.3	25	20
	M4	8.1	4.2	7.4	6.9	10	5
	M5	1.5	2.6	3.6	2.8	35	30
process 2	M6	2.3	2.4	2.2	2.9	30	25
	M7	3.5	3.6	3.2	2.6	23	18
	M8	4.3	4.3	4.8	3.6	18	13
	M9	5.2	4.5	4.8	4.8	12	7

	M10	6.1	5.2	5.6	5.1	8	5
	M11	8.1	4.8	4.5	4.8	10	5
process 3	M12	7.2	3.3	3.4	3.3	18	13
	M13	3.1	2.5	1.5	1.9	35	30
	M14	5.3	3.4	2.1	2.1	23	18
	M15	9	5.3	5.3	5.6	8	5
	M16	2.6	1.5	2.2	2.8	25	20
process 4	M17	3.6	2.5	4.3	3.5	20	15
	M18	3.8	1.9	3.4	3.5	23	18
	M19	5.1	2.6	5.6	4.6	15	10
	M20	6.2	3.9	6.3	5.3	10	5
	M21	4.5	4.8	7.5	7.3	10	5
process 5	M22	3.6	3.4	6.6	6.8	13	8
	M23	2.7	3.5	5.8	5.6	15	10
	M24	1.3	2.6	3.7	3.3	20	15
	M25	1.5	2.6	4.8	4.6	18	13

219 Note: P_T means the processing time for each jobs with the corresponding machines

220 Case study

221 A case was presented to testify the feasibility and effectiveness of the proposed methodology for
 222 minimizing makespan and energy consumption of HFSP. In this case, four jobs were processed on 25
 223 machines and each job requires 5 steps operations. There were 5 parallel machines in each process. The
 224 related data including job number, processing time, process power and wait power were provided in
 225 Table 3.

226 In addition, the sustainability of each machine were evaluated and ranked with the fuzzy TOPSIS
 227 method listed in Table 4. The higher rank would be selected in the decoding processing of the RBDE
 228 algorithm. Also, the proposed method was implemented in Matlab 2014 and runs on an Intel Core i5
 229 CPU (2.53Ghz/8.00G RAM) PC with a Windows 10 operation system. All the parameters and their
 230 selected values are summarized in Table 5.

231

Table 4 Ranking value of the machines

processes	Machines	(value/rank)			
		job1	job2	job3	job4
process1	M1	0.59/	0.40/5	0.55/4	0.55/5
	M2	0.30/	0.31/1	0.48/3	0.41/2
	M3	0.50/	0.36/4	0.60/5	0.48/4
	M4	0.58/	0.31/2	0.48/2	0.48/3
	M5	0.19/	0.32/3	0.44/1	0.35/1
process2	M6	0.27/	0.28/1	0.28/1	0.38/3
	M7	0.41/	0.40/3	0.38/2	0.28/1

	M8	0.46/	0.43/5	0.48/5	0.36/2
	M9	0.49/	0.41/4	0.43/4	0.43/5
	M10	0.47/	0.38/2	0.41/3	0.38/4
process3	M11	0.70/	0.35/4	0.35/4	0.40/4
	M12	0.75/	0.27/1	0.31/3	0.33/3
	M13	0.79/	0.29/2	0.15/1	0.23/1
	M14	0.61/	0.31/3	0.20/2	0.23/2
	M15	0.71/	0.36/5	0.39/5	0.44/5
process4	M16	0.25/	0.13/1	0.22/1	0.33/1
	M17	0.30/	0.23/3	0.42/3	0.39/2
	M18	0.39/	0.17/2	0.34/2	0.41/3
	M19	0.46/	0.23/4	0.50/5	0.47/4
	M20	0.51/	0.32/5	0.50/4	0.48/5
process5	M21	0.33/	0.38/5	0.60/4	0.56/3
	M22	0.31/	0.32/3	0.65/5	0.65/5
	M23	0.22/	0.34/4	0.60/3	0.58/4
	M24	0.09/	0.25/2	0.44/1	0.36/1
	M25	0.08/	0.25/1	0.54/2	0.51/2

232

233

Table 5 Parameters of the RBDE

Parameters	constraint
initial population size	100
range of crossover probability	[0.1, 1]
range of mutation probability	[0.01, 0.2]
scale factor	0.5
stopping condition	Maximum number of iterations 200 or convergence condition $\ \Delta f\ \leq 10^{-3}$

234

235 Effectiveness of RBDE

236 To test the effectiveness of RBDE, we expanded the instance size by increasing the times of
 237 the data of workpiece/process/machines in Table 2. For the simplicity of presentation, an instance
 238 with n jobs, s processes and m machines were denoted as an $n*s*m$. For each instance, RBDE
 239 algorithm was run for 50 times independently and took the average value of the solutions. The running
 240 time and two extreme results of the Pareto front set were listed in Table 6.

241

Table 6 Optimization results of RBDE

No.	instances	extreme point1		extreme point2		running time (min)
		Makespan (min)	energy consumption (kW.h)	Makespan (min)	energy consumption (kW.h)	
1	4*5*25	21	2404	22.9	1949	34.252
2	8*5*25	26.5	4067	33.8	3089	61.710
3	12*5*25	30.8	4900	38.2	4253	90.262
4	16*5*25	36.6	6051	40.3	5587	117.240
5	20*5*25	41.8	7113	50.2	6791	143.815

6	25*5*25	47.6	9315	56.7	8374	178.785
7	4*10*50	38.4	8249	39.9	6821	61.978
8	8*10*50	46.9	1.15e+04	53.1	1.059e+04	114.406
9	12*10*5	51.2	1.630e+04	59.0	1.3042e+04	168.970
10	16*10*5	58.9	1.733e+04	63	1.64e+04	220.235
11	20*10*5	64.7	1.914e+04	67	1.873e+04	271.624
12	25*10*5	71.7	2.223e+04	74.9	2.185e+04	339.596

242 To visualize the performance of RBDE, we selected the pareto front solutions of No.3 and No.9
 243 instance. It can be seen from Fig. 6 that the results are distributed evenly and widely, which denoted the
 244 effectiveness of the proposed methods.

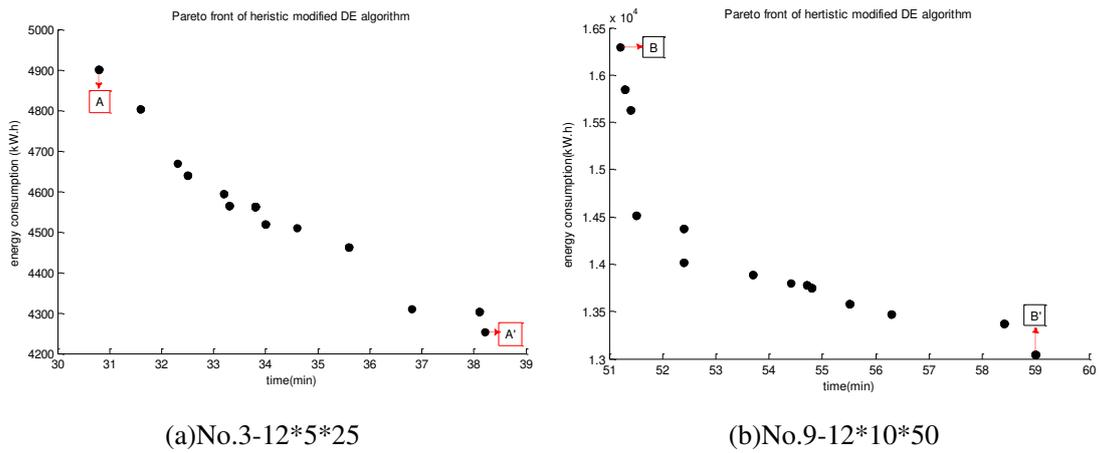
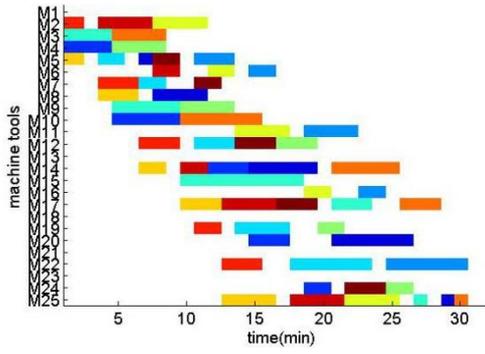


Fig.6 Pareto front solutions of RBDE

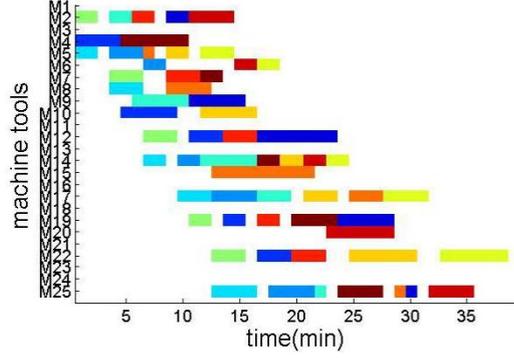
245
 246 To further assess the versatility of the proposed algorithm, the extreme solution points A, A', B, B'
 247 were selected from the obtained Pareto front solutions in Fig.6. Gantt charts of the four test instances
 248 are showed in Fig.7.

249 The optimal solutions corresponding to the four instance are {8,10,11,4,2,7,12,6,5,3,1,9},
 250 {4,6,2,3,10,5,12,8,9,1,11,7}, {10,2,12,11,3,6,4,5,7,8,1,9} and {7,1,11,9,6,5,10,3,2,12,4,8} respectively.
 251 Different ranking value of parallel machine tools for each process and jobs can be obtained according
 252 to section 2, as shown in Table 6. It can be seen that the utilization rates of M1, M8, M9, M10, M13,
 253 M11, M18, M21 and M23 in parallel machines is lower from Gantt chart (a) and (b). Besides, it also
 254 indicated that the data of parallel machine tool in (c), (d) is double in (a), (b), which is consistent with
 255 the ranking of parallel machines.

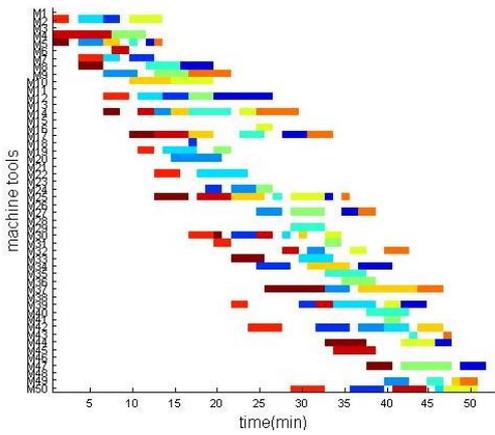
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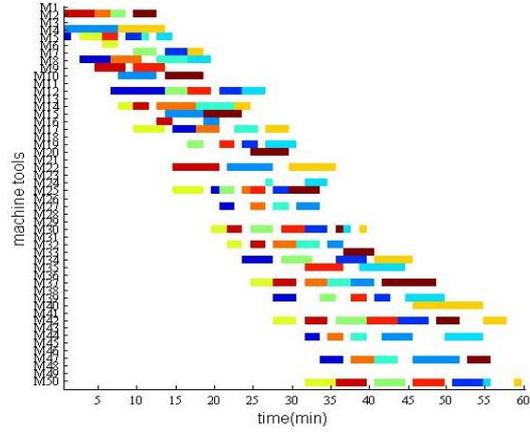
12*5*25- A (a)



12*5*25- A' (b)



12*10*50-B (c)



12*10*50-B' (d)

257

Fig.7 Gantt chart for the four instances

258

Table 7 Number of non-dominated solutions found by each algorithm

NO	instance	RBDE			HMPSO		HMNSGA-II			
		m	min	avg	max	min	avg	max	min	avg
1	4*5*25	6	4	5.2	4	3	3.1	5	3	4.6
2	8*5*25	18	10	14.5	15	9	11.3	17	10	13.2
3	12*5*25	21	12	14.2	16	11	12.5	17	12	13.4
4	16*5*25	28	10	21.7	13	8	10.3	23	10	18.5
5	20*5*25	22	15	19.4	18	12	15.6	20	13	18.9
6	25*5*25	28	13	20.1	18	12	17.4	24	12	18.3
7	4*10*50	5	3	4.5	3	2	2.6	4	3	3.2
8	8*10*50	12	9	11.3	11	5	7.9	13	7	9.8
9	12*10*50	24	13	22.4	18	11	16.3	21	14	18.1
10	16*10*50	23	15	19.5	17	10	12.6	20	12	16.4
11	20*10*50	17	10	12.3	14	8	10.5	15	9	11.2
12	25*10*50	15	7	9.6	12	6	8.3	13	6	8.9

259

260 **Comparisons with other algorithms**

261 In order to test the performance of the proposed algorithms, the heuristic rule was implanted in the
 262 NSGA-II and PSO named as HMNSGA-II and HMPSO. The ‘Max.’, ‘Avg.’, and ‘Min.’ column
 263 represent the maximum, average, and minimum number of non-dominated solutions, respectively. It
 264 can be clearly found from Table 7 that the overall AVG, MAX and MIN yielded by RBDE were better
 265 than those generated by HMNSGA-II and HMPSO algorithms in the same computation time. The
 266 reason can be explained that the proposed algorithm can take full of the non-dominated solutions to
 267 generate excellent offspring and the heuristic rule for iteration operators to disturb old individuals.

268 Besides, in this paper, we utilize the five performance metrics to evaluate the improved methods,
 269 i.e, convergence metric , uniformity performance , diversity metric , hyper-volume and running
 270 time.

271 (1) Convergence metric γ , it means the average value of the minimum distance between each
 272 reference point in the set p and the reference set p^* .

$$273 \quad \gamma = \frac{\sum_{x \in p} \min_{y \in p^*} dis(x, y)}{|p|} \quad (23)$$

274 Where, p is the solution set obtained by the algorithm $dis(x,y)$ represents the Euclidean distance
 275 between point y in reference set p^* and point x in reference set p .

276 (2) Uniformity performance: Spacing Metric sp , it was used to measure the standard deviation of
 277 the minimum distance from each solution to others.

$$278 \quad sp = \sqrt{\frac{1}{|p|} \sum_{i=1}^{|p|} (\bar{d} - d_i)^2} \quad (24)$$

279 where, d_i is the minimum distance from t point i on the pareto frontier of the algorithm to the other
 280 points, \bar{d} is the average value of all distance d_i .

281 (3) Diversity metric Δ , measures the extent of spread achieved among the obtained solutions.

$$282 \quad \Delta = \frac{df + dl + \sum_{i=1}^{N-1} |d_i - \bar{d}|}{df + dl + (N-1)\bar{d}} \quad (25)$$

283 where, df and dl are the Euclidean distances between the extreme solutions and the boundary solutions

284 of the obtained non-dominated set. Assuming that there are N solutions on the best non-dominated
 285 front.

286 (4) Hyper-volume, HV : The volume of the region in the target space enclosed by the
 287 non-dominant solution set and the reference point obtained by the algorithm.

$$288 \quad HV = \delta \left(\bigcup_{i=1}^{|\mathcal{S}|} v_i \right) \quad (26)$$

289 where, δ is the Lebesgue measure, which was used to measure volume. $|\mathcal{S}|$ represents the number of
 290 non-dominated solution sets. v_i means the Super volume composed of the reference point and the i th
 291 solution in the solution set.

292 (5) Running time describes the execution time of an algorithm to reflect the efficiency.

293 For all three algorithms, the population size is set as 100, the number of iterations is 200. For each
 294 instance, all the tested algorithms are run 20 times and the performance metrics averaged are collected,
 295 as shown in Table 8.

296 Table 8 optimization results of algorithms

instance	RBDE					HMPSO					HMNSGA-II					
	γ	Sp	Δ	HV	t_{max}	γ	Sp	Δ	HV	t_{max}	γ	Sp	Δ	HV	t_{max}	
1	4*5*25	69.004	141.443	0.601	13.500	18.246	69.004	141.443	0.601	13.500	66.274	69.004	141.443	0.601	13.500	34.252
2	8*5*25	191.375	112.814	0.748	444.200	30.698	191.407	134.001	1.409	230.300	45.226	198.719	166.435	0.826	239.300	61.710
3	12*5*25	123.606	55.893	0.691	331.800	44.282	153.508	82.550	0.750	150.100	75.452	228.204	100.232	0.730	197.000	90.262
4	16*5*25	162.305	48.048	0.856	353.300	57.873	249.635	64.752	0.946	309.100	69.200	363.703	101.415	0.894	309.100	117.240
5	20*5*25	84.329	55.441	0.855	3879.00	72.763	154.146	86.655	0.967	1068.000	82.845	276.667	161.591	0.935	1506.400	143.815
6	25*5*25	80.378	85.157	0.612	966.500	91.404	140.715	195.662	0.938	333.800	102.018	267.857	98.147	0.916	949.300	178.785
7	4*10*50	35.001	37.789	0.072	37.000	30.860	35.001	37.789	0.072	37.000	112.785	35.001	37.789	0.072	37.000	61.978
8	8*10*50	94.415	163.862	0.853	135.000	56.585	94.415	163.862	0.853	135.000	95.206	94.415	163.862	0.853	135.000	114.406
9	12*10*50	103.750	173.268	0.672	312.000	82.263	113.755	249.986	0.749	120.000	93.207	128.334	181.530	0.729	247.000	168.970
10	16*10*50	112.312	158.39	0.715	321.781	106.638	198.527	205.687	0.741	125.416	120.049	267.815	176.326	0.738	273.246	220.235
11	20*10*50	323.002	123.574	0.733	347.000	132.872	453.750	196.003	0.799	136.000	148.580	761.250	180.193	0.742	146.000	271.624
12	25*10*50	150.000	94.439	0.524	439.500	166.060	236.000	125.623	0.792	376.813	186.421	350.473	82.815	0.734	409.000	339.596

297
 298 Table 8 shows the statistical performance metrics obtained by each of the three algorithms. It is
 299 revealed that RBDE outperforms HMNSAG-II and HMPSO for almost all the metrics, particularly for

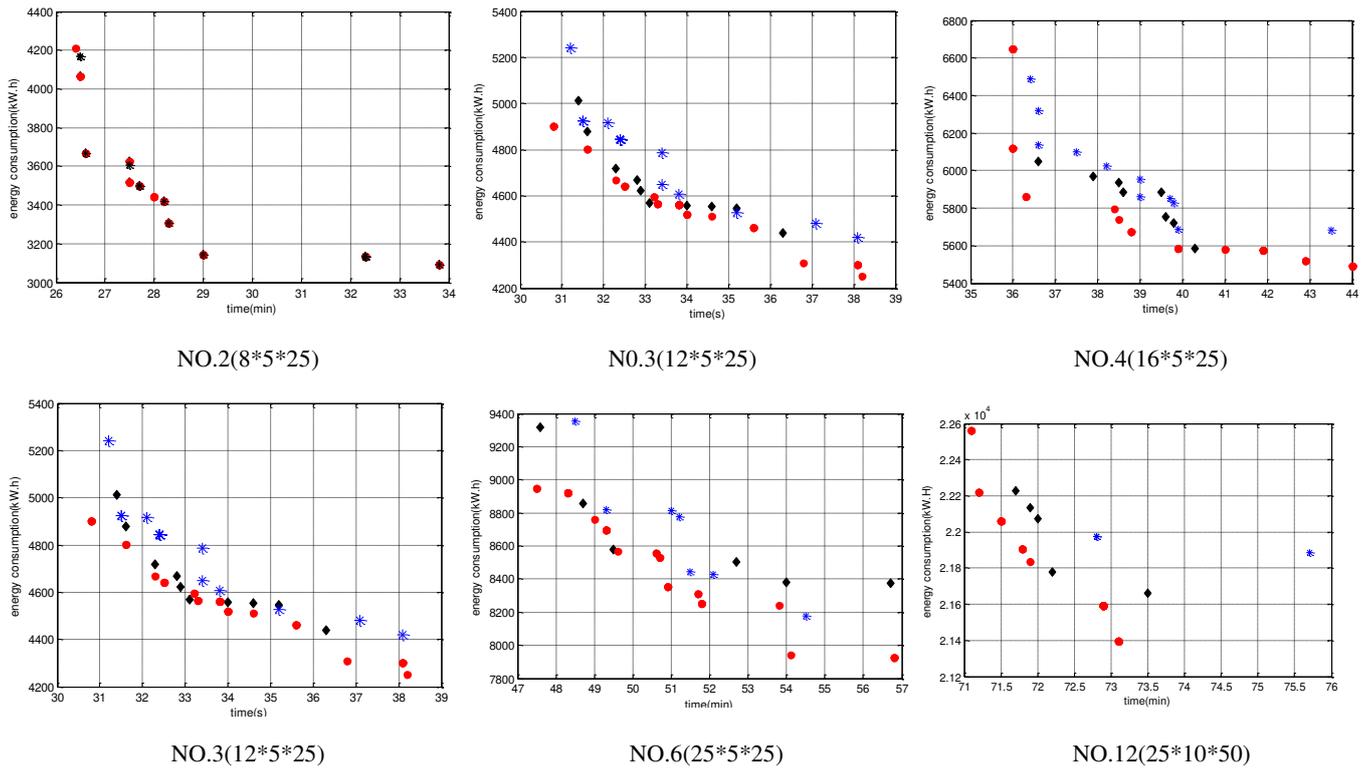
300 γ , sp and running time, which means that RBDE can achieves the most proximity to the reference front.
301 In terms of convergence γ , RBDE and HMPSO can quickly converge to a better value than
302 HMNSAG-II. However, HMPSO is easy to fall into local convergence, and the universality and
303 diversity of RBDE and NSGA-II (Δ and HV) is significantly superior to the HMPSO. From the
304 uniformity performance perspective (i.e, sp), it can be observed that DE and PSO outperforms
305 NSGA-II when the problem size is small. However, such advantage will no longer exist if the problem
306 size is getting larger. In large scale problems, RBDE and HMNSGA-II can find more non-dominated
307 solutions, and it performs especially better in solving those problems with less number of stages and
308 machines. In addition, RBDE is the fastest in solving multi-objective HFSP problems, which means the
309 proposed algorithms have a better searching ability than the others for the problem studied in this work.

310 To further visualize the performance of these algorithms, the pareto fronts are selected to provide a
311 graphical representation for different scare of workpiece instance, i.e, No.2, No.3, No.4 and different
312 scare instance No.3, No.6, No.12, which are selected as typical examples of the small, medium and large
313 scale instances, respectively. Results are demonstrated in Fig. 8.

314 Fig.8 shows zthe optimal pareto frontier as the function of iteration number. As shown, the curves
315 of all the three schemes converge toward the optimal pareto frontier. These figures give an intuitive
316 illustration for some results derived from the performance metrics analysis. As shown in the figure, the
317 RBDE algorithm performs the best among the three algorithms in solution quality and diversity.
318 Besides, it is noted that the RBDE algorithm provides solutions that are relatively closer to the Pareto
319 fronts, whereas the RBDE and HMNSGA-II algorithms provide solutions that are more diversified.
320 This result well explains the above- mentioned performance anomaly between the three algorithms.

321 Overall, the proposed RBDE algorithm is capable of providing better solutions than HMNSGA-II
322 and HMPSO in terms of quality and distribution. The heuristic rule for modified DE algorithm has a
323 positive effect on the behavior of the proposed algorithm.

324



325 Fig.8 Scatter of RBDE, HMNSGA-II and HMPSO algorithms(●-RBDE, ◆-HMNSGA-II *
 326 -HMPSO)

327 **Conclusions**

328 This work investigates a multi-objective hybrid flow shop sustainable scheduling model based on
 329 a modified differential evolution algorithm. The main contributions of the proposed algorithm are
 330 concluded as follows:

- 331 (1) A sustainable scheduling mathematical model considering makespan and energy consumption
- 332 model was established for the hybrid flow shop.
- 333 (2)The sustainability of parallel machines were defined by a hierarchy criteria, which was ranked
- 334 by fuzzy TOPSIS method.
- 335 (3)A ranking-based DE was developed to produce feasible scheduling sequences, in which
- 336 realized active selection of parallel machines with high sustainability. The effectiveness of the proposed
- 337 methods was verified.

338 Although the efficiency of the proposed method has been verified, this paper has some limitations.
 339 1) It does not to use actual jobs-shop data to validate this method to provide the best sustainable
 340 scheduling and 2) It just focuses on the energy-consumption and makespan problem. The future
 341 research should make use of multiobjective optimization tools to achieve the comprehensive

342 optimization by considering large numbers of variables and constraints, e.g., set-up time of machines,
343 limited resources, and requirements of customer order.

344

345 **Author contributions** L.W. and L.K. conceptualised the mathematical model. X.L. made the
346 program. L.W. and F.L. analysed the data and made figures. L.W. wrote the paper. All the authors read
347 and contributed .to the submitted version of the manuscript. L.W. and F.L. acquired the funding and
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349

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354

355 **Ethical Approval**

356 Not applicable.

357

358 **Consent to Participate**

359 All authors agree to participate in the editing of the paper.

360

361 **Consent to Publish**

362 All authors agree to publish this manuscript in this journal.

363

364 **Competing Interests**

365 The authors declare that they have no conflict of interest.

366

367 **Availability of data and materials**

368 Not applicable.

369

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