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Configuration models as an urn problem: the generalized hypergeometric ensemble of random graphs

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ABSTRACT

A fundamental issue of network data science is the ability to discern observed features that can be expected at random from those beyond such expectations. Configuration models play a crucial role there, allowing us to compare observations against degree-corrected null-models. Nonetheless, existing formulations have limited large-scale data analysis applications either because they require expensive Monte-Carlo simulations or lack the required flexibility to model real-world systems. With the generalized hypergeometric ensemble, we address both problems. To achieve this, we map the configuration model to an urn problem, where edges are represented as balls in an appropriately constructed urn. Doing so, we obtain a random graph model reproducing and extending the properties of standard configuration models, with the critical advantage of a closed-form probability distribution.

Introduction

Essential features of complex systems are inferred by studying the deviations of empirical observations from suitable stochastic models. Network models, in particular, have become the state of the art for complex systems analysis, where systems' constituents are represented as vertices, and their interactions are viewed as edges and modelled by means of edge probabilities in a random graph. Such models are built to preserve certain properties of the observed network. How well the model describes other properties of the network distinguishes between randomly expected patterns and “interesting” ones.

The simplest of random graph models, known as the $G(n, p)$, generates edges between a given number of vertices with a fixed probability $p \in (0, 1]$ ¹. In this model, the properties of vertices, such as their degrees, are homogeneously distributed, i.e., they have all the same expected value. However, most empirical graphs show heterogeneous, heavy tailed degree distributions²⁻⁵. Hence, random graph models able to incorporate *specified degree sequences* are of primal importance.

The most common family of models that fix degree sequences is known as the *configuration model* of random graphs. The comparison of an empirical graph with the corresponding configuration model allows quantifying which properties of the original graph can be ascribed to the degree sequence. The properties not explained by the degree sequence highlight the unique structure of the studied empirical graph. In particular, configuration models are invaluable for the macroscopic and the mesoscopic analysis of graphs. Its most notable applications, among others, are graph partitioning through modularity maximisation^{6,7} and quantifying degree correlations⁸.

The concept of a configuration model originates in the early works of Béla Bollobàs, who introduced the term ‘configuration’ to refer to the arrangement of edges in a model⁹. Since then, different configuration models preserving degrees have been proposed. In its standard formulation, the configuration model refers to the uniform sampling of a graph from the space of all graphs with a given degree sequence¹⁰. Such formulation thus fixes exactly the degree sequence and the number of edges in the model. Later, Molloy and Reed^{11,12} popularized the model, leading to the common reference in literature of *Molloy-Reed configuration model*.

Although being the most common version of the configuration model, that of Molloy and Reed is not the only possible one¹³. In particular, different types of configuration models arise if one allows for the existence of multi-edges or not, and of self-loops or not. Similarly, it is possible to distinguish between *vertex-labelled configuration models* and *stub-labelled configuration models*, as thoroughly discussed by¹³. Vertex-labelled graphs are those graphs for which there is a bijection between graph and adjacency matrix. This means that, in the case of a multi-edge graph, swapping the end-points of two multi-edges between the same pair of vertices leads back to the same graph. Stub-labelled graphs, on the other hand, are graphs for which swapping the end-points of two multi-edges between the same pair of vertices *does not lead back to the same graph*. For this reason, multiple distinct stub-labelled graphs may correspond to the same adjacency matrix. In this article, we focus

only on vertex-labelled graphs, i.e., graphs that are completely identified by their adjacency matrix.

The standard configuration model has a crucial drawback: it can only be realized through a Monte-Carlo process, in the form of a repeated rewiring procedure. Each vertex is assigned a number of half-edges, or stubs, corresponding to their degree. A random realisation of the model is obtained by wiring pairs of stubs together uniformly at random. This is a computationally expensive procedure, which does not allow to explore the whole probability space of the model. This problem is exacerbated in the case of larger graphs and graphs with highly heterogeneous degree distributions.

This limitation is a consequence of the hard constraints defining these models. Specifically, it stems from the choice of fixing exactly both the degree sequence of the model and the total number of edges, characteristics common in most variants of the standard configuration model. The only suitable approach to realize such models is that of Markov Chain Monte Carlo sampling (MCMC)¹³. The most popular algorithms rely on double edge swaps, also known as degree-preserving rewirings^{14,15}, checkerboard swaps^{16–18}, tetrads¹⁹, and alternating rectangles²⁰, which randomize a given graph by simultaneously rewiring a pair of edges. This procedure had been already introduced in²¹ in 1891, and had been rediscovered many times later on^{22–24}.

A solution to the lack of tractability and the poor scalability of MCMC based configuration models is relaxing the hard constraints characterising these models such that MCMC is not needed anymore. For this reason, there exist variants of the standard configuration model that relax the hard constraints of preserving exactly degree sequence and number of edges. Such models are often referred to as *soft-configuration models*. The soft-constrained graph model most closely related to the standard configuration model is the so-called *Chung-Lu model*^{4,25}. This model assumes that each edge (i, j) is drawn *independently* from a Poisson distribution with parameters $p_{ij} = \mathbf{k}_i^{\text{out}} \mathbf{k}_j^{\text{in}} / m$, where $\mathbf{k}_i^{\text{out}}$ is the out-degree of vertex i . The Chung-Lu model can be seen as generalising the concept of the $G(n, p)$ model¹, such that not only the expected number of edges of the original graph is preserved, but also its degree distribution. In these terms, the standard configuration model instead can be seen as the degree-preserving counterpart of the $G(n, m)$ model¹⁰, which generates all graphs with n vertices and m edges *uniformly*.

The Chung-Lu model is generally simpler to work with than the standard configuration model. However, it has the disadvantage of specifying not the actual values, but only the *expected* number of edges, and the *expected* degree sequence. This, among other things, means that it cannot specify the exact degree distribution¹⁰. Because of the essential role played by degree distributions, this issue has negatively affected the adoption of this model, in spite of its other advantages. Nevertheless, such an approach has found important applications, e.g., in the degree-corrected stochastic block model²⁶.

There exist further soft-configuration models formulations that fall into the family of exponential random graph models^{10,13,27}. However, we will not discuss these here because they rely on MCMC as well²⁸, and the mixing times of these chains are known to be very poor^{29,30}.

In this article, we propose an analytically tractable model for random graphs with *given expected degree sequences* and a *fixed number of edges*. The model positions itself in between the standard configuration model, which fixes exactly the degree sequences, and the Chung-Lu model that preserves the number of edges in expectation and assumes that all edges are independent of each other. Thanks to this, it closely matches the properties of the standard configuration model, with the advantages given by an analytical formulation. The formulation of our model relies on mapping the process of drawing edges to a multivariate urn problem. Such model has the incomparable advantage of possibly incorporating *patterns that go beyond degree sequences*. In its simplest case, our model corresponds to a configuration model for directed or undirected multi-edge graphs, that fixes the values of vertex degrees in expectation instead of their exact values. In its general form, the model incorporates a parameter for each pair of vertices, which we call *edge propensities*. These parameters control the relative likelihood of drawing an individual edge between the respective pair of vertices, as opposed to any other pair. This is achieved by biasing the combinatorial edge drawing process. Any graph patterns can be modelled within this formulations, as long as they can be reduced to dyadic relations between pairs of vertices.

In the next section, we (i) provide formal definitions of the generalized hypergeometric ensemble of random graphs for the directed and undirected cases, and (ii) investigate the properties of our model and how they compare against existing configuration models.

Results

Let us consider a multi-graph $\mathcal{G} = (V, E)$, where V is a set of n vertices, and $E \subseteq V \times V$ is a multi-set of m (directed or undirected) multi-edges. We indicate with \mathbf{A} the adjacency matrix of the graph where entries $A_{ij} \in \mathbb{N}_0$ capture the multiplicity of an edge $(i, j) \in V \times V$ in the multi-set E . In the case of undirected graphs, the adjacency matrix is symmetric, i.e., $\mathbf{A} = \mathbf{A}^T$, and the elements on its diagonal equal twice the multiplicity of the corresponding self-loops. Because we only deal with multi-graphs, in the rest of the article we will refer to multi-graphs simply as graphs, and to multi-edges simply as edges.

Hypergeometric configuration model

The concept underlying our random graph model is the same as that of the standard configuration model, which is to randomly shuffle the multi-edges of a graph \mathcal{G} while preserving vertex degrees. The standard configuration model generates multi-edges

one after another by sampling uniformly at random a vertex with an available out-stub (outwards half-edge) and a vertex with an available in-stub (inwards half-edge) until all stubs are consumed. Figure 1 illustrates one step of this process. The resulting random graphs all have exactly the same degree sequence as the original graph \mathcal{G} . All pairs of available in- and out-stubs are equiprobable to be picked, so are the corresponding individual multi-edges. Therefore, the probability of observing a multi-edge between a given pair of vertices positively relates to the number of possible stub-pairings of the two vertices, which in turn is defined by the corresponding degrees of these. This probability depends only on the degrees of the two vertices and on the total number of multi-edges in the graph.

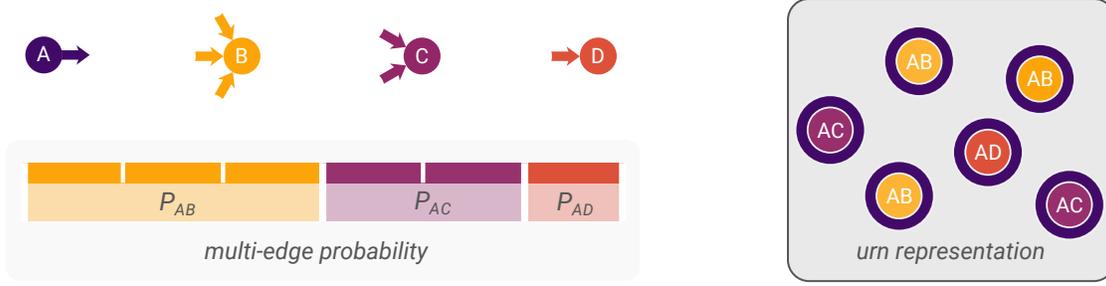


Figure 1. The configuration model represented (**upper left**) as a conventional edge rewiring process and (**right**) as an urn problem. In the former case, once the out-stub (A, \cdot) has been sampled for rewiring, then one in-stub is sampled uniformly at random from those available, to draw a new multi-edge. If we represent each possible combination of an out-stub and an in-stub as a ball, we arrive at the urn problem without replacement. For the shown vertices, the odds of observing a multi-edge (A, B) are three times higher than of observing a multi-edge between (A, D) and 1.5 times higher than of observing a multi-edge between (A, C) in both model representations.

The need to consequently sample two vertices at each step makes it cumbersome to formulate analytically the procedure described above. To overcome this challenge, we take an *edge-centric approach* of sampling m multi-edges from a certain larger multi-set, which we define below in Definition 1. As a consequence of this change of perspective, the model will preserve the *expected* degree sequences instead of the exact ones, resulting in a soft-configuration model. To this end, we introduce the definition of the *hypergeometric (soft-) configuration model*.

For each pair of vertices $i, j \in V$, we define the number Ξ_{ij} of stub combinations that exist between vertices i and j , which can be conveniently represented in matrix form:

Definition 1 (Combinatorial matrix). We define combinatorial matrix $\Xi \in \mathbb{N}^n \times \mathbb{N}^n$ for graph \mathcal{G} as the matrix whose elements Ξ_{ij} are defined as

$$\Xi_{ij} = \mathbf{k}_i^{\text{out}}(\mathcal{G}) \mathbf{k}_j^{\text{in}}(\mathcal{G}) \quad \text{for } i, j \in V, \quad (1)$$

where $\mathbf{k}^{\text{out}}(\mathcal{G})$ and $\mathbf{k}^{\text{in}}(\mathcal{G})$ are the out-degree and in-degree sequences of the graph \mathcal{G} .

Hypergeometric configuration model for directed graphs

Definition 2 (Directed hypergeometric configuration model). Let $\hat{\mathbf{k}}^{\text{in}}, \hat{\mathbf{k}}^{\text{out}} \in \mathbb{N}^n$ be in- and out-degree sequences and \hat{V} a set of n vertices. The hypergeometric configuration model (HypE) X generated by $(\hat{V}, \hat{\mathbf{k}}^{\text{in}}, \hat{\mathbf{k}}^{\text{out}})$ is the n^2 -dimensional random vector X defined on the probability space (S, P) with sample space

$$S = \{\mathcal{G}(V, E) \mid |E| = m\}, \quad m = \sum_{i \in V} \hat{\mathbf{k}}_i^{\text{in}} = \sum_{i \in V} \hat{\mathbf{k}}_i^{\text{out}}, \quad (2)$$

with some probability measure P , such that individual multi-edges are equiprobable and the expected degree sequences of a realisation of X are fixed:

$$\mathbb{E}_P [\mathbf{k}^{\text{in}}(X)] = \hat{\mathbf{k}}^{\text{in}}, \quad \mathbb{E}_P [\mathbf{k}^{\text{out}}(X)] = \hat{\mathbf{k}}^{\text{out}}. \quad (3)$$

This set-up allows to map the model to an *urn problem* and thus to arrive at a closed-form probability distribution function for it. Before doing so, we introduce the concept of *induced random model*.

Definition 3 (Graph-induced random model). We say that the graph $\mathcal{G}(V, E)$ induces the random model X , if the quantities $(\hat{V}, \hat{\mathbf{k}}^{\text{in}}, \hat{\mathbf{k}}^{\text{out}})$ generating X are computed from \mathcal{G} . I.e., $\hat{V} = V$, $\hat{\mathbf{k}}^{\text{in}} = \mathbf{k}^{\text{in}}(\mathcal{G})$, and $\hat{\mathbf{k}}^{\text{out}} = \mathbf{k}^{\text{out}}(\mathcal{G})$.

Under this assumption, we can formulate the following theorem for the distribution of the hypergeometric configuration model X . To keep the notation simple, we will not distinguish between the $n \times n$ adjacency matrix \mathbf{A} and the vector of length n^2 obtained by stacking it by row or column. Similarly, we do the same for all other related $n \times n$ matrices.

Theorem 1 (Distribution of the hypergeometric configuration model). *Let $\mathcal{G}(V, E)$ be a directed graph with $n = |V|$ vertices and $m = |E|$ multi-edges. Let $\mathbf{k}^{\text{in}}(\mathcal{G}) \in \mathbb{N}^n$ and $\mathbf{k}^{\text{out}}(\mathcal{G}) \in \mathbb{N}^n$ be the vectors representing its in-degree and out-degree sequences. Let X be the hypergeometric configuration model induced by \mathcal{G} defined as in Definition 2. If the probability measure P depends only on the degree sequences in \mathcal{G} and the total number of multi-edges $m = |E|$ and all multi-edges are equiprobable, then X follows the multivariate hypergeometric distribution as in Eq. (4).*

Let $\mathbf{A} \in \mathbb{N}_0^n \times \mathbb{N}_0^n$ be an adjacency matrix and $\Xi \in \mathbb{N}^n \times \mathbb{N}^n$ be the combinatorial matrix induced by \mathcal{G} . Then the hypergeometric configuration model X is distributed as follows:

$$\Pr(X = \mathcal{G}) = \frac{\prod_{i,j \in V} \binom{\Xi_{ij}}{A_{ij}}}{\binom{M}{m}}, \quad (4)$$

where $M = \sum_{i,j \in V} \Xi_{ij}$ is the total number of stub combinations between all vertices.

Proof. We want to sample m multi-edges connecting any of the in- and out-stub pairs such that all such multi-edges are equiprobable. The total number of stubs combinations between any two vertices i, j is given by $\Xi_{ij} = \mathbf{k}_i^{\text{out}} \mathbf{k}_j^{\text{in}}$ (cf. Lemma 4 in the SI for a proof). We can hence define the random graph model as follows. We sample m multi-edges without replacement from the multi-set of size $\sum_{i,j \in V} \Xi_{ij}$ that combines all the possible stub pairs combinations Ξ_{ij} between all pairs $i, j \in V$. We sample without replacement because we need to mimic the process of wiring stubs. Once a stub pair has been used, it cannot be sampled again. We can view this model as an *urn problem* where the edges to be sampled are represented by balls in an urn. By representing the multi-edges connecting each different pair of vertices (i, j) as balls of a unique colour, we obtain an urn with a total of $M = \sum_{i,j \in V} \Xi_{ij}$ balls of $n^2 = |V \times V|$ different colours. With this, the sampling of a graph according to our model, corresponds to drawing exactly m balls from this urn. Each adjacency matrix \mathbf{A} with $\sum_{i,j \in V} A_{ij} = m$ corresponds to one particular realisation drawn from this model. The probability of drawing exactly $\mathbf{A} = \{A_{ij}\}_{i,j \in V}$ edges between each pair of vertices is given by the multivariate hypergeometric distribution. \square

From Theorem 1 we derive the following results, whose proofs follow directly from properties of the hypergeometric distribution.

Corollary 1.1. *For each pair of vertices $i, j \in V$, the probability that X has exactly A_{ij} edges between i and j is given by the marginal distributions of the multivariate hypergeometric distribution, i.e.*

$$\Pr(X_{ij} = A_{ij}) = \frac{\binom{\Xi_{ij}}{A_{ij}} \binom{M - \Xi_{ij}}{m - A_{ij}}}{\binom{M}{m}}. \quad (5)$$

Corollary 1.2. *The expected in- and out-degree sequences of realisations of the directed hypergeometric configuration model X correspond to the respective degree sequences of the graph \mathcal{G} inducing X .*

Proof. For each pair of vertices i, j we can calculate the expected number of multi-edges $\mathbb{E}[X_{ij}]$ as

$$\mathbb{E}[X_{ij}] = m \frac{\Xi_{ij}}{M} \quad (6)$$

Moreover, summing the rows and columns of matrix $\mathbb{E}[X_{ij}]$ and assuming directed graphs with self-loops we can calculate the expected in- or out-degrees of all vertices as

$$\begin{aligned} \mathbb{E}[k_j^{\text{in}}(X)] &= \sum_{i \in V} \mathbb{E}[X_{ij}] = m \frac{\sum_{i \in V} \hat{k}_i^{\text{out}} \hat{k}_j^{\text{in}}}{M} = \hat{k}_j^{\text{in}}, \\ \mathbb{E}[k_i^{\text{out}}(X)] &= \sum_{j \in V} \mathbb{E}[X_{ij}] = m \frac{\sum_{j \in V} \hat{k}_i^{\text{out}} \hat{k}_j^{\text{in}}}{M} = \hat{k}_i^{\text{out}}. \end{aligned} \quad (7)$$

Equation (7) confirms that the *expected* in- and out-degree sequence of realisations drawn from X corresponds to the degree sequence of the given graph \mathcal{G} . \square

Hypergeometric configuration model for undirected graphs

So far, we have discussed the hypergeometric configuration model for directed graphs. Specifying the undirected case, on the other hand, requires some more efforts. The reason for this is that, under the assumptions described in the previous section, the undirected version of the hypergeometric configuration model is the degenerate case of its directed counterpart, where the direction of the multi-edges is ignored. In particular, this implies that the random vector corresponding to the undirected model has half the dimensions of the directed one, because any undirected multi-edge between two vertices i and j can either be generated as a directed multi-edge (i, j) or as a directed multi-edge (j, i) . We provide a formal proof of this in Theorem 6, in the SI.

Definition 4 (Undirected hypergeometric configuration model). Let $\hat{\mathbf{k}} \in \mathbb{N}^n$ be a degree sequence and \hat{V} a set of n vertices. The hypergeometric configuration model X generated by $(\hat{V}, \hat{\mathbf{k}})$ is the $(n^2 + n)/2$ -dimensional random vector X defined on the probability space (S, P) , for the sample space

$$S = \{\mathcal{G}(V, E) \mid |E| = m\}, \quad 2m = \sum_{i \in V} \hat{k}_i, \quad (8)$$

with some probability measure P , such that individual multi-edges are equiprobable and the expected degree sequence of a realisation of X is fixed:

$$\mathbb{E}_P[\mathbf{k}(X)] = \hat{\mathbf{k}}. \quad (9)$$

At first sight, it would appear that the undirected hypergeometric configuration model can simply be obtained by restricting the directed model to $n(n+1)/2$ components corresponding to the upper-triangle and the diagonal of the adjacency matrix. That is, we would sample m multi-edges among pairs $i \leq j \in V$ from the multi-set of stub combinations Ξ_{ij} as defined in Eq. (1). However, the resulting model does not satisfy Definition 4, because its expected degree sequence does not equal the degree sequence that induced it. To show this, we follow the same reasoning adopted in the proof of Corollary 1.2. The expected degree $\mathbb{E}[k_i(X)]$ of a vertex i is equivalent to

$$\mathbb{E}[k_i(X)] = \sum_{j \in V} \mathbb{E}[X_{ij}] = m \frac{\sum_{i \in V} \hat{k}_i \hat{k}_j}{\sum_{i \leq j \in V} \hat{k}_i \hat{k}_j} \neq \hat{k}_i.$$

This approach is wrong because the total number of undirected stub combinations is larger than in the directed case. The reason for this is the symmetry in the process of wiring two stubs. Let \hat{k}_i be the degree of vertex i and \hat{k}_j the degree of vertex j . To form a multi-edge (i, j) , each one of the \hat{k}_i stubs of i can be connected to all \hat{k}_j stubs of j , and vice versa, each of the \hat{k}_j stubs of j can be connected to all \hat{k}_i stubs of i . Hence, the total number of combinations of stubs between vertices i and j equals to $\hat{k}_i \hat{k}_j + \hat{k}_j \hat{k}_i = 2\hat{k}_i \hat{k}_j$.

We can now formulate the equivalent of Theorem 1 for the undirected case. The distribution underlying the undirected hypergeometric configuration model can then be computed analogously to Theorem 1.

Theorem 2 (Distribution of undirected hypergeometric configuration model). Let $\mathcal{G}(V, E)$ be an undirected graph with $n = |V|$ vertices and $m = |E|$ multi-edges. Let $\mathbf{k} \in \mathbb{N}^n$ be the vector representing its degree sequence. Let X be the undirected hypergeometric configuration model induced by \mathcal{G} defined as in Definition 4. If the probability distribution underlying X depends only on the degree sequence of \mathcal{G} and the total number of multi-edges m , and all multi-edges are equiprobable, then X follows the multivariate hypergeometric distribution given in Eq. (10).

Let $\mathbf{A} \in \mathbb{N}_0^n \times \mathbb{N}_0^n$ be the symmetric adjacency matrix corresponding to an undirected graph, and $\Xi_{ij} = \mathbf{k}_i \mathbf{k}_j$ be the combinatorial matrix induced by \mathcal{G} . Then the undirected hypergeometric configuration model X is distributed as follows:

$$\Pr(X = \mathcal{G}) = \frac{\prod_{i < j \in V} \binom{2\Xi_{ij}}{A_{ij}} \prod_{l \in V} \binom{\Xi_{ll}}{A_{ll}/2}}{\binom{M}{m}}, \quad (10)$$

where $M = \sum_{i < j \in V} 2\Xi_{ij} + \sum_{l \in V} \Xi_{ll} = \sum_{i, j \in V} \Xi_{ij}$ is the total number of undirected stub combinations between all pair of vertices.

Proof. The proof follows the same reasoning of the proof of Theorem 1, accounting for the fact that the total number of stubs combinations is now given by

$$\begin{cases} 2\Xi_{ij} & \text{if } i \neq j, \\ \Xi_{ii} & \text{if } i = j. \end{cases} \quad (11)$$

□

The undirected version of Corollaries 1.1 and 1.2 are obtained by choosing the suitable probability distribution, given by Eq. (10). For completeness, we report them in the SI.

Comparison with other configuration models

In the sections above, we have provided a parsimonious formulation of a soft-configuration model. Specifically, we have done so in terms of an *hypergeometric ensemble*. The ensemble provides a random graph model formulation for directed and undirected graphs alike, in which (i) the expected in- and out-degree sequences are fixed, and (ii) multi-edges between these vertices with fixed expected degrees are formed uniformly at random. More precisely, the probability for a particular pair of vertices to be connected by an edge is only influenced by combinatorial effects, and thus only depends on the degrees of the vertices and the total number of edges.

We can now discuss the differences of this hypergeometric ensemble with the standard hard-constrained model and Chung-Lu configuration model variants. In particular, we compare here (i) the sample spaces of the three models, and (ii) the distributions characterising the models. In the following, we take the standard model as the reference point.

Sample spaces of the configuration models. For a given degree sequence \mathbf{k} , the standard configuration model samples *uniformly at random* from the space $S_{\mathbf{k}}$ of all graphs with given degree sequence \mathbf{k} . It can be easily specified for both directed or undirected graphs. The Chung-Lu model is instead defined on the space S of *all* graphs. Differently, the hypergeometric ensemble is defined on the space S_m consisting of all graphs with a given number of edges $m = \sum \mathbf{k}$. Because $S_{\mathbf{k}} \subseteq S_m \subseteq S$, the hypergeometric ensemble is closer to the standard configuration model than the Chung-Lu model. The sample space of the Chung-Lu model is much larger than that of the hypergeometric ensemble and, furthermore, this increases the variance of the distribution of degrees, as we outline in the next paragraph.

Distributions of the three models. The standard configuration model assumes all graphs with a given degree sequence to be equiprobable. Differently from this, both the hypergeometric ensemble and Chung-Lu model define a non-uniform probability distribution over a larger set of graphs. However, there is a major difference between the two models. While the hypergeometric ensemble –similarly to the standard configuration model– does not assume the independence of edge probabilities, in the Chung-Lu model edges are independent. This results in a large variance for both the distribution of the number of edges in the graph and the distribution of the degree of a vertex.

In the case of the Chung-Lu model, the probability of observing a directed graph \mathcal{G} follows the Poisson distribution

$$\Pr(X_{CL} = \mathcal{G}) := \prod_{ij} (A_{ij}!)^{-1} \left(\frac{\mathbf{k}_i^{\text{out}} \mathbf{k}_j^{\text{in}}}{m} \right)^{A_{ij}} e^{-\frac{\mathbf{k}_i^{\text{out}} \mathbf{k}_j^{\text{in}}}{m}}. \quad (12)$$

This means that the probability distribution of the out-degree $\mathbf{k}_i^{\text{out}}$ follows a Poisson distribution too, with parameter

$$\lambda := \sum_j \frac{\mathbf{k}_i^{\text{out}} \mathbf{k}_j^{\text{in}}}{m} = \mathbf{k}_i^{\text{out}}. \quad (13)$$

The same holds for in-degrees. This result follows from basic probability theory as the degree is simply the sum of n independent Poisson random variables¹⁰. Similarly, the total number of edges m in the graph follows a Poisson distribution with parameter $\lambda = m$. Because the variance of a Poisson distribution with parameter λ is λ , we see that the variance of the number of edges in model grows linearly with the size of the graph.

In the case of the hypergeometric ensemble, the probability to observe a directed graph \mathcal{G} is given by the hypergeometric distribution in Eq. (4). Hence, the distribution of the out-degree $\mathbf{k}_i^{\text{out}}$ follows a hypergeometric distribution as well, as noted in the following corollary to Theorem 1.

Corollary 1.3 (Distribution of degrees). *Let \mathcal{G} be a graph and $\mathbf{k}_i^{\text{out}}$ the out-degree of vertex i . Under the hypergeometric ensemble, the probability that $\mathbf{k}_i^{\text{out}} = k$ follows the hypergeometric distribution*

$$\Pr(k) := \frac{\binom{m \cdot \mathbf{k}_i^{\text{out}}}{k} \binom{M - m \cdot \mathbf{k}_i^{\text{out}}}{m - k}}{\binom{M}{m}}. \quad (14)$$

This result follows from the properties of the hypergeometric distribution, in particular from the fact that generating a out-degree k for vertex i corresponds to sampling exactly k out-stubs for vertex i from the pool of all out-stubs available $\sum_j \Xi_{ij} = \sum_j \mathbf{k}_i^{\text{out}} \mathbf{k}_j^{\text{in}} = m \cdot \mathbf{k}_i^{\text{out}}$. From the previous corollary, we recover the value for the variance of the distribution of the degree

<i>Model</i>	CM	HypE	CL
<i>Distribution</i>	uniform	hypergeometric	Poisson
<i>Variance of m</i>	—	—	m
<i>Variance of degrees</i>	—	$\mathbf{k}_i^{\text{out}} \frac{(m - \mathbf{k}_i^{\text{out}})}{(m+1)}$	$\mathbf{k}_i^{\text{out}}$

Table 1. Properties of the three configuration models CM, HypE, and CL, arranged by increasing variance in the corresponding distributions.

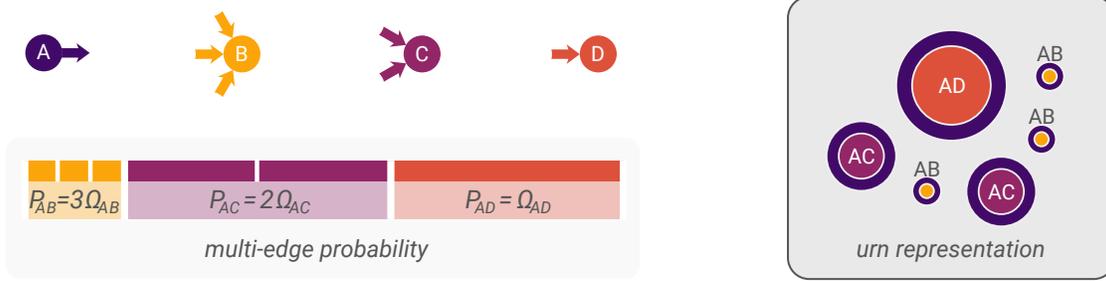


Figure 2. The effect of *edge propensities* on the configuration model. Differently from the standard configuration model, here the stubs are not sampled uniformly at random (cf. Fig. 1). Given an out-stub, each in-stub is characterized by a propensity Ω_{ij} of being chosen. As a result, the probability of wiring the out-stub (A, \cdot) to the vertex D is larger than that of B due to a very large edge propensity Ω_{AD} , even though vertex B has three times more in-stubs than vertex D .

of a vertex which is equal to $\mathbf{k}_i^{\text{out}}(m - \mathbf{k}_i^{\text{out}})(m + 1)^{-1}$. Because this variance is always smaller than $\mathbf{k}_i^{\text{out}}$, the hypergeometric ensemble (HypE) more closely reproduces the standard configuration model (CM) compared to the Chung-Lu model (CL). Table 1 summarizes the differences between the three models.

Numerical comparison. To further highlight the properties of the hypergeometric ensemble and to show how it relates with the other two configuration models, we perform a simple numerical investigation. First, we choose a degree sequence as a starting point. For this task, we use an empirical degree sequence, obtained from the top 285 vertices of the Gentoo collaboration network, ordered by degree, (cf.³¹). The network shows a heavy-tailed degree distribution, with a total of 1366 edges. Having fixed the degree sequence, we generate 50000 realisations from each model and report the distribution of basic graph properties. In Fig. 1 provided in the SI, we also show the simulated distributions obtained for the number of edges m , the mean squared error (MSE) of the degree sequences against the true one, the degree centralisation, and the degree assortativity.

A qualitative comparison shows that, in general, the hypergeometric ensemble fares considerably more closely to the standard configuration model than the Chung-Lu model. This is a consequence of the properties highlighted in the previous paragraph. Such qualitative evaluation can be confirmed quantitatively by comparing the distributions of the different properties. The realisations from the hypergeometric ensembles have in fact (i) a significantly smaller MSE ($p < 1e - 16$), i.e., the degree sequences of the generated graphs deviates less from the original degree sequence, and (ii) a significantly smaller variance in the degree centralisation (bootstrapped p-value $p < 1e - 12$). All three models provide similar degree assortativity distributions ($p > .1$).

Generalized hypergeometric ensemble of graphs

Configuration models sample edges uniformly at random. However, such uniform sampling of edges is generally not enough to describe empirical systems. In empirical graphs, edge probabilities do not only depend on the degree of vertices but also on other characteristics. Examples of such characteristics are vertex labels that lead to observable group structures, distances between vertices, similarities, etc. Below we discuss how such influences on the sampling probabilities of edges can be encoded in a random graph model by means of a dyadic property we call *edge propensity*.

First, we introduce the concept of edge propensity. Given two dyads (i, j) and $(k, l) \in V \times V$ where $\Xi_{ij} = \Xi_{kl}$, according to the hypergeometric configuration model the probabilities of sampling one multi-edge between (i, j) and (k, l) are equal, leading to an odds-ratio of 1 between the two pairs of vertices. Instead, we generalize the model in a way that fixing arbitrary odds-ratios is possible. That is, we specify edge propensities such that the ratio between them is the odds-ratio between the two corresponding vertex pairs of sampling a multi-edge, all else being equal. We use edge propensities to bias the sampling probability of each multi-edge, as illustrated in Fig. 2. This way, the probability of the number of multi-edges between a pair of vertices depends on both the degrees of the vertices and their edge propensity.

We encode edge propensities in a matrix $\mathbf{\Omega}$ defined as follows.

Definition 5 (Propensity matrix). Let $\mathbf{\Omega} = (\Omega_{ij})_{i,j \in V} \in \mathbb{R}^{|V| \times |V|}$ be a $n \times n$ matrix where n is the number of vertices. Let (i, j) and $(k, l) \in V \times V$. Let $\omega_{i,j,kl}$ the odds-ratio of sampling one multi-edge between (i, j) instead of (k, l) . The entries Ω_{ij} and Ω_{kl} of the propensity matrix $\mathbf{\Omega}$ are then defined such that $\Omega_{ij}/\Omega_{kl} = \omega_{i,j,kl}$. This implies that the propensity matrix $\mathbf{\Omega}$ is defined up to a constant, as multiplying $\mathbf{\Omega}$ with any constant preserves the specified odds-ratios.

Now, we define the generalized hypergeometric ensemble (gHyPE), a random graph model that combines degree-related combinatorial effects, i.e., the configuration model, and the newly introduced edge propensities. We do so by using propensities to *bias* the sampling process described above. In the urn model analogy, such biased sampling implies that the probability of drawing a certain number of balls of a given colour (i.e., multi-edges between the corresponding pair of vertices) depends both on their number and their size, as illustrated in Fig. 2. The probability distribution resulting from such a biased sampling process is given by the multivariate Wallenius' non-central hypergeometric distribution^{32,33}.

Theorem 3 (Distribution of gHypEG). Let $\mathcal{G}(V, E)$ be a directed graph with $n = |V|$ vertices and $m = |E|$ edges. Under the assumptions introduced above, the generalized hypergeometric ensemble of graphs (gHypEG) X induced by \mathcal{G} and a given propensity matrix $\mathbf{\Omega}$ follows the multivariate Wallenius' non-central hypergeometric distribution given in Eq. (15).

Let $\mathbf{A} \in \mathbb{N}^n \times \mathbb{N}^n$ be the adjacency matrix associated with \mathcal{G} and $\mathbf{\Xi} \in \mathbb{N}^n \times \mathbb{N}^n$ be its combinatorial matrix defined in Eq. (1). Then the gHypEG defined by $\mathbf{\Xi}$ and $\mathbf{\Omega}$, \mathcal{G} is distributed as follows:

$$\Pr(X = \mathbf{A}) = \left[\prod_{i,j \in V} \binom{\Xi_{ij}}{A_{ij}} \right] \int_0^1 \prod_{i,j \in V} \left(1 - z^{\frac{\Omega_{ij}}{S_{\mathbf{\Omega}}}} \right)^{A_{ij}} dz \quad (15)$$

with

$$S_{\mathbf{\Omega}} = \sum_{i,j \in V} \Omega_{ij} (\Xi_{ij} - A_{ij}). \quad (16)$$

The distribution describing the biased sampling from an urn is a generalisation of the multivariate hypergeometric distribution. The proof of Theorem 3 follows from the fact that, when the sampling is performed without replacement with given relative odds, this sampling process corresponds to the multivariate Wallenius' non-central hypergeometric distribution. Details of this derivation can be found in³² for the univariate case, and in^{34,35} for the multivariate case. A thorough review of non-central hypergeometric distributions has been done in³⁶.

The next two corollaries directly follow from properties of Wallenius' non-central hypergeometric distribution.

Corollary 3.1. For each pair of vertices $i, j \in V$, the probability to draw exactly A_{ij} edges between i and j is given by the marginal distributions of the multivariate Wallenius' non-central hypergeometric distribution, i.e.,

$$\Pr(X_{ij} = A_{ij}) = \binom{\Xi_{ij}}{A_{ij}} \binom{M - \Xi_{ij}}{m - A_{ij}} \int_0^1 \left[\left(1 - z^{\frac{\Omega_{ij}}{S_{\mathbf{\Omega}}}} \right)^{A_{ij}} \left(1 - z^{\frac{\bar{\Omega}_{ij}}{S_{\mathbf{\Omega}}}} \right)^{m - A_{ij}} \right] dz \quad (17)$$

where

$$\bar{\Omega}_{ij} = \frac{\sum_{(l,m) \in (V \times V) \setminus \{(i,j)\}} \Xi_{lm} \Omega_{lm}}{(M - \Xi_{ij})}. \quad (18)$$

Corollary 3.2. The entries of the expected adjacency matrix $\mathbb{E}[X_{ij}]$ can be obtained by solving the following system of equations:

$$\left(1 - \frac{\mathbb{E}[X_{11}]}{\Xi_{11}} \right)^{\frac{1}{\Omega_{11}}} = \left(1 - \frac{\mathbb{E}[X_{12}]}{\Xi_{12}} \right)^{\frac{1}{\Omega_{12}}} = \dots \quad (19)$$

with the constraint $\sum_{i,j \in V} \mathbb{E}[X_{ij}] = m$.

The undirected formulation of the gHypEG follows the same reasoning of Theorem 2, with the addition of a symmetric propensity matrix. That is, the distribution of the undirected generalized hypergeometric ensemble is hence given by a Wallenius' distribution similar to Eq. (15), but corresponding to the upper triangular part of the matrices (i.e., for $i \leq j$) in accordance with Eq. (11).

Estimation of the propensity matrix

In this final section, we show how to define the propensity matrix Ω such that the expected graph $\mathbb{E}[X]$ from the model X defined by Ω coincides with an arbitrary graph G . By doing so, we create a random model *centered* around the inducing graph. The result is described in the following corollary, which follows from the properties of 'Wallenius' non-central hypergeometric distribution.

Corollary 3.3. *Let $\mathcal{G}(V, E)$ be a graph with $n = |V|$ vertices and A its adjacency matrix. Let Ω be a $n \times n$ propensity matrix, characterized by elements Ω_{ij} with $i, j \in V$. Then, \mathcal{G} coincides with the expectation $\mathbb{E}[X]$ of the gHypEG X induced by \mathcal{G} and Ω if and only if the following relation holds.*

$$\forall c \in \mathbb{R}^+ \quad \Omega_{ij} = \frac{1}{c} \log(1 - A_{ij}/\Xi_{ij}) \quad \forall i, j \in V, \quad (20)$$

Where Ξ is the combinatorial matrix associated with X and Ξ_{ij} its elements.

Proof. Equation (20) follows directly from Corollary 3.1. In particular, when solving Eq. (19) for Ω with the assumption $\mathbb{E}[X] = A$ we obtain the following system of $|V|^2$ equations (in the case of a directed graph with self-loops) for $|V|^2 + 1$ variables.

$$\begin{cases} \left(1 - \frac{\mathbb{E}[X_{11}]}{\Xi_{11}}\right)^{\frac{1}{\Omega_{11}}} = C \\ \left(1 - \frac{\mathbb{E}[X_{12}]}{\Xi_{12}}\right)^{\frac{1}{\Omega_{12}}} = C \\ \vdots \end{cases} \quad (21)$$

The solution of this system is Eq. (20). □

A wide range of statistical patterns that go beyond degree effects can be encoded in the graph model by specifying the matrix Ω of edge propensities. The encoding and fitting techniques of such arbitrary propensity matrices are beyond the scope of this article, and will not be discussed here.

Finally, we highlight that the hypergeometric configuration model is a special case of the generalized hypergeometric ensemble. The hypergeometric configuration model described in Theorem 1 can be recovered from the generalized model by setting all entries in the propensity matrix to the same value. By doing so, the odds-ratio between the propensities for any pair of vertices is 1, and the edge sampling process is not biased. Thus, the probability distribution of the model reduces to a function of the degree sequences and the number of multi-edges sampled. Theorem 7, in the SI, provides a formal proof for this result.

Discussion

We have proposed a novel framework for the study of random multi-edge graphs in terms of an urn problem. By doing so, we have arrived at an analytically tractable formulation of the widely used configuration model of random graphs. It preserves degree sequences in expectation and the number of edges exactly. Our model relies on representing edges as balls in an urn containing the appropriate number of them, as explained in Definition 1. Sampling without replacement from such an urn generates random realizations from the model. Thanks to this parallel, we can identify the underlying probability distribution as the *multivariate hypergeometric distribution*. As we show, this formulation accurately reproduces the properties of standard configuration models, with the added benefit of a closed-form formulation.

We have further expanded such a configuration model to the *generalized hypergeometric ensemble*, a new class of models that can incorporate arbitrary dyadic biases into the probabilities of sampling edges. These biases, which we call edge propensities, control the relative likelihood of drawing an individual edge between the respective pair of vertices instead of any other pair. The proposed formulation allows incorporating into the model arbitrary graph patterns that can be reduced to dyadic relations. Again, in Theorem 3 we exploit the parallel to an urn problem to formalize the probability distribution underlying the generalized model: the *multivariate Wallenius' non-central hypergeometric distribution*.

In³⁷, we provide an open-source software implementation of the generalized hypergeometric ensemble within the R library GHYPERNET. The library is freely available for download on CRAN at <https://cloud.r-project.org/package=ghypernet>. The full library documentation can be found at <https://ghyper.net>.

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Author contributions statement

G.C. and V.N. developed the methods and wrote the manuscript. All authors reviewed the manuscript.

Additional information

Competing interests.

The authors declare no competing interests.

Figures

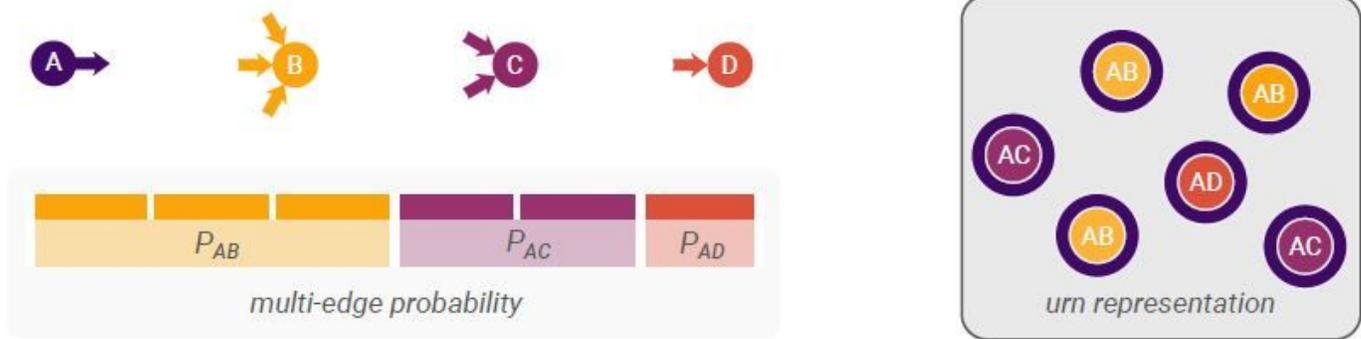


Figure 1

The configuration model represented (upper left) as a conventional edge rewiring process and (right) as an urn problem. In the former case, once the out-stub (A; .) has been sampled for rewiring, then one in-stub is sampled uniformly at random from those available, to draw a new multi-edge. If we represent each possible combination of an out-stub and an in-stub as a ball, we arrive at the urn problem without replacement. For the shown vertices, the odds of observing a multi-edge (A;B) are three times higher than of observing a multi-edge between (A;D) and 1:5 times higher than of observing a multi-edge between (A;C) in both model representations.

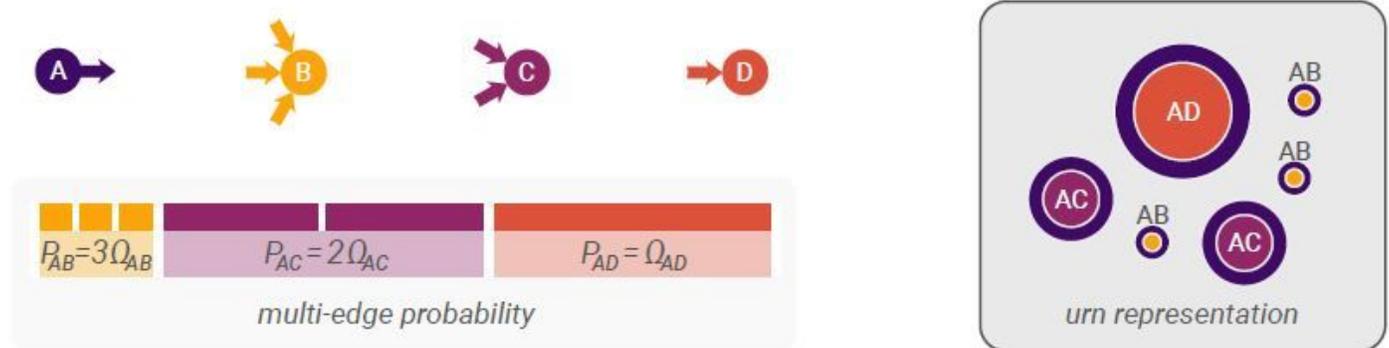


Figure 2

The effect of edge propensities on the configuration model. Differently from the standard configuration model, here the stubs are not sampled uniformly at random (cf. Fig. 1). Given an out-stub, each in-stub is characterized by a propensity W_{ij} of being chosen. As a result, the probability of wiring the out-stub (A; .) to the vertex D is larger than that of B due to a very large edge propensity W_{AD} , even though vertex B has three times more in-stubs than vertex D.